Semantics in Generative Grammar
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Semantics in Generative Grammar

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Preface

This book is an introduction to the craft of doing formal semantics for linguists. It is not an overview of the field and its current developments. The many recent handbooks provide just that. We want to help students develop the ability for semantic analysis, and, in view of this goal, we think that exploring a few topics in detail is more effective than offering a bird's-eye view of everything. We also believe that foundational and philosophical matters can be better discussed once students have been initiated into the practice of semantic argumentation. This is why our first chapter is so short. We dive right in.

The students for whom we created the lectures on which this book is based were graduate or advanced undergraduate students of linguistics who had already had a basic introduction to some formal theory of syntax and had a first understanding of the division of labor between semantics and pragmatics. Not all of them had had an introduction to logic or set theory. If necessary, we filled in gaps in formal background with the help of other books.

We learned our craft from our teachers, students, and colleagues. While working on this book, we were nurtured by our families and friends. Paul Hirschbühler, Molly Diesing, Kai von Fintel, Jim Higginbotham, the students in our classes, and reviewers whose names we don’t know gave us generous comments we could use. The staff of Blackwell Publishers helped us to turn lecture notes we passed back and forth between the two of us into a book. We thank them all.

Irene Heim and Angelika Kratzer
Cambridge, Mass., and Amherst, April 1997
1 Truth-conditional Semantics and the Fregean Program

1.1 Truth-conditional semantics

To know the meaning of a sentence is to know its truth-conditions. If I say to you

(1) There is a bag of potatoes in my pantry

you may not know whether what I said is true. What you do know, however, is what the world would have to be like for it to be true. There has to be a bag of potatoes in my pantry. The truth of (1) can come about in ever so many ways. The bag may be paper or plastic, big or small. It may be sitting on the floor or hiding behind a basket of onions on the shelf. The potatoes may come from Idaho or northern Maine. There may even be more than a single bag. Change the situation as you please. As long as there is a bag of potatoes in my pantry, sentence (1) is true.

A theory of meaning, then, pairs sentences with their truth-conditions. The results are statements of the following form:

**Truth-conditions**
The sentence “There is a bag of potatoes in my pantry” is true if and only if there is a bag of potatoes in my pantry.

The apparent banality of such statements has puzzled generations of students since they first appeared in Alfred Tarski’s 1935 paper “The Concept of Truth in Formalized Languages.”¹ Pairing English sentences with their truth-conditions seems to be an easy task that can be accomplished with the help of a single schema:

**Schema for truth-conditions**
The sentence “______” is true if and only if ______.
A theory that produces such schemata would indeed be trivial if there wasn't another property of natural language that it has to capture: namely, that we understand sentences we have never heard before. We are able to compute the meaning of sentences from the meanings of their parts. Every meaningful part of a sentence contributes to its truth-conditions in a systematic way. As Donald Davidson put it:

The theory reveals nothing new about the conditions under which an individual sentence is true; it does not make those conditions any clearer than the sentence itself does. The work of the theory is in relating the known truth conditions of each sentence to those aspects ("words") of the sentence that recur in other sentences, and can be assigned identical roles in other sentences. Empirical power in such a theory depends on success in recovering the structure of a very complicated ability — the ability to speak and understand a language.\(^2\)

In the chapters that follow, we will develop a theory of meaning composition. We will look at sentences and break them down into their parts. And we will think about the contribution of each part to the truth-conditions of the whole.

### 1.2 Frege on compositionality

The semantic insights we rely on in this book are essentially those of Gottlob Frege, whose work in the late nineteenth century marked the beginning of both symbolic logic and the formal semantics of natural language. The first worked-out versions of a Fregean semantics for fragments of English were by Lewis, Montague, and Cresswell.\(^3\)

It is astonishing what language accomplishes. With a few syllables it expresses a countless number of thoughts, and even for a thought grasped for the first time by a human it provides a clothing in which it can be recognized by another to whom it is entirely new. This would not be possible if we could not distinguish parts in the thought that correspond to parts of the sentence, so that the construction of the sentence can be taken to mirror the construction of the thought. . . . If we thus view thoughts as composed of simple parts and take these, in turn, to correspond to simple sentence-parts, we can understand how a few sentence-parts can go to make up a great multitude of sentences to which, in turn, there correspond a great multitude of thoughts. The question now arises how the construction of the thought proceeds, and by what means the parts are put together
so that the whole is something more than the isolated parts. In my essay "Negation," I considered the case of a thought that appears to be composed of one part which is in need of completion or, as one might say, unsaturated, and whose linguistic correlate is the negative particle, and another part which is a thought. We cannot negate without negating something, and this something is a thought. Because this thought saturates the unsaturated part or, as one might say, completes what is in need of completion, the whole hangs together. And it is a natural conjecture that logical combination of parts into a whole is always a matter of saturating something unsaturated.4

Frege, like Aristotle and his successors before him, was interested in the semantic composition of sentences. In the above passage, he conjectured that semantic composition may always consist in the saturation of an unsaturated meaning component. But what are saturated and unsaturated meanings, and what is saturation? Here is what Frege had to say in another one of his papers.

Statements in general, just like equations or inequalities or expressions in Analysis, can be imagined to be split up into two parts; one complete in itself, and the other in need of supplementation, or "unsaturated." Thus, e.g., we split up the sentence

"Caesar conquered Gaul"

into "Caesar" and "conquered Gaul." The second part is "unsaturated" — it contains an empty place; only when this place is filled up with a proper name, or with an expression that replaces a proper name, does a complete sense appear. Here too I give the name "function" to what this "unsaturated" part stands for. In this case the argument is Caesar.5

Frege construed unsaturated meanings as functions. Unsaturated meanings, then, take arguments, and saturation consists in the application of a function to its arguments. Technically, functions are sets of a certain kind. We will therefore conclude this chapter with a very informal introduction to set theory. The same material can be found in the textbook by Partee et al.6 and countless other sources. If you are already familiar with it, you can skip this section and go straight to the next chapter.

1.3 Tutorial on sets and functions

If Frege is right, functions play a crucial role in a theory of semantic composition. "Function" is a mathematical term, and formal semanticists nowadays use it in
Truth-conditional Semantics

exactly the way in which it is understood in modern mathematics. Since functions are sets, we will begin with the most important definitions and notational conventions of set theory.

1.3.1 Sets

A set is a collection of objects which are called the “members” or “elements” of that set. The symbol for the element relation is “\( \in \)”. \( x \in A \) reads “\( x \) is an element of \( A \)”. Sets may have any number of elements, finite or infinite. A special case is the empty set (symbol “\( \emptyset \)”), which is the (unique) set with zero elements.

Two sets are equal iff they have exactly the same members. Sets that are not equal may have some overlap in their membership, or they may be disjoint (have no members in common). If all the members of one set are also members of another, the former is a subset of the latter. The subset relation is symbolized by “\( \subseteq \)”. \( A \subseteq B \) reads “\( A \) is a subset of \( B \)”.

There are a few standard operations by which new sets may be constructed from given ones. Let \( A \) and \( B \) be two arbitrary sets. Then the intersection of \( A \) and \( B \) (in symbols: \( A \cap B \)) is that set which has as elements exactly the members that \( A \) and \( B \) share with each other. The union of \( A \) and \( B \) (in symbols: \( A \cup B \)) is the set which contains all the members of \( A \) and all the members of \( B \) and nothing else. The complement of \( A \) in \( B \) (in symbols: \( B - A \)) is the set which contains precisely those members of \( B \) which are not in \( A \).

Specific sets may be defined in various ways. A simple possibility is to define a set by listing its members, as in (1).

(1) Let \( A \) be that set whose elements are \( a \), \( b \), and \( c \), and nothing else.

A more concise rendition of (1) is (1’).\(^9\)

(1’) \( A := \{a, b, c\} \).

Another option is to define a set by abstraction. This means that one specifies a condition which is to be satisfied by all and only the elements of the set to be defined.

(2) Let \( A \) be the set of all cats.

(2’) Let \( A \) be that set which contains exactly those \( x \) such that \( x \) is a cat.
(2'), of course, defines the same set as (2); it just uses a more convoluted formulation. There is also a symbolic rendition:

\( A := \{ x : x \text{ is a cat} \}. \)

Read "\( \{ x : x \text{ is a cat} \} \)" as "the set of all \( x \) such that \( x \) is a cat". The letter "\( x \)" here isn't meant to stand for some particular object. Rather, it functions as a kind of place-holder or variable. To determine the membership of the set \( A \) defined in (2''), one has to plug in the names of different objects for the "\( x \)" in the condition "\( x \) is a cat". For instance, if you want to know whether \( \text{Kaline} \in A \), you must consider the statement "Kaline is a cat". If this statement is true, then \( \text{Kaline} \in A \); if it is false, then \( \text{Kaline} \notin A \) ("\( x \notin A \)" means that \( x \) is not an element of \( A \)).

### 1.3.2 Questions and answers about the abstraction notation for sets

**Q1:** If the "\( x \)" in "\( \{ x : x \text{ is a positive integer less than 7} \} \)" is just a place-holder, why do we need it at all? Why don't we just put a blank as in "\( \{ _: _ \text{ is a positive integer less than 7} \} \)"?

**A1:** That may work in simple cases like this one, but it would lead to a lot of confusion and ambiguity in more complicated cases. For example, which set would be meant by "\( \{ _: \{ _: _ \text{ likes } _ \} = \emptyset \} \)"? Would it be, for instance, the set of objects which don't like anything, or the set of objects which nothing likes? We certainly need to distinguish these two possibilities (and also to distinguish them from a number of additional ones). If we mean the first set, we write "\( \{ x : \{ y : x \text{ likes } y \} = \emptyset \} \)". If we mean the second set, we write "\( \{ x : \{ y : y \text{ likes } x \} = \emptyset \} \)".

**Q2:** Why did you just write "\( \{ x : \{ y : y \text{ likes } x \} = \emptyset \} \)" rather than "\( \{ y : \{ x : x \text{ likes } y \} = \emptyset \} \)"?

**A2:** No reason. The second formulation would be just as good as the first, and they specify exactly the same set. It doesn't matter which letters you choose; it only matters in which places you use the same letter, and in which places you use different ones.

**Q3:** Why do I have to write something to the left of the colon? Isn't the condition on the right side all we need to specify the set? For example, instead of "\( \{ x : x \text{ is a positive integer less than 7} \} \)", wouldn't it be good enough to write simply "\( \{ x \text{ is a positive integer less than 7} \} \)"?
A3: You might be able to get away with it in the simplest cases, but not in more complicated ones. For example, what we said in A1 and A2 implies that the following two are different sets:

\[
\{x : \{y : x \text{ likes } y\} = \emptyset\}
\]
\[
\{y : \{x : x \text{ likes } y\} = \emptyset\}
\]

Therefore, if we just wrote “\(\{x \text{ likes } y\} = \emptyset\)”, it would be ambiguous. A mere statement enclosed in set braces doesn’t mean anything at all, and we will never use the notation in this way.

Q4: What does it mean if I write “\{California : California is a western state\}”?

A4: Nothing, it doesn’t make any sense. If you want to give a list specification of the set whose only element is California, write “\{California\}”. If you want to give a specification by abstraction of the set that contains all the western states and nothing else but those, the way to write it is “\{x : x \text{ is a western state}\}”. The problem with what you wrote is that you were using the name of a particular individual in a place where only place-holders make sense. The position to the left of the colon in a set-specification must always be occupied by a place-holder, never by a name.

Q5: How do I know whether something is a name or a place-holder? I am familiar with “California” as a name, and you have told me that “x” and “y” are place-holders. But how can I tell the difference in other cases? For example, if I see the letter “a” or “d” or “s”, how do I know if it’s a name or a place-holder?

A5: There is no general answer to this question. You have to determine from case to case how a letter or other expression is used. Sometimes you will be told in so many words that the letters “b”, “c”, “t”, and “u” are made-up names for certain individuals. Other times, you have to guess from the context. One very reliable clue is whether the letter shows up to the left of the colon in a set-specification. If it does, it had better be meant as a place-holder rather than a name; otherwise it doesn’t make any sense. Even though there is no general way of telling names apart from place-holders, we will try to minimize sources of confusion and stick to certain notational conventions (at least most of the time). We will normally use letters from the end of the alphabet as place-holders, and letters from the beginning of the alphabet as names. Also we will never employ words that are actually used as names in English (like “California” or “John”) as place-holders. (Of course, we could so use them if we wanted to, and then
we could also write things like “[California : California is a western state]”, and it would be just another way of describing the set \{x : x is a western state\}. We could, but we won’t.)

Q6: In all the examples we have had so far, the place-holder to the left of the colon had at least one occurrence in the condition on the right. Is this necessary for the notation to be used properly? Can I describe a set by means of a condition in which the letter to the left of the colon doesn’t show up at all? What about “[x : California is a western state]”?

A6: This is a strange way to describe a set, but it does pick one out. Which one? Well, let’s see whether, for instance, Massachusetts qualifies for membership in it. To determine this, we take the condition “California is a western state” and plug in “Massachusetts” for all the “x”s in it. But there are no “x”s, so the result of this “plug-in” operation is simply “California is a western state” again. Now this happens to be true, so Massachusetts has passed the test of membership. That was trivial, of course, and it is evident now that any other object will qualify as a member just as easily. So \{x : California is a western state\} is the set containing everything there is. (Of course, if that’s the set we mean to refer to, there is no imaginable good reason why we’d choose this of all descriptions.) If you think about it, there are only two sets that can be described at all by means of conditions that don’t contain the letter to the left of the colon. One, as we just saw, is the set of everything; the other is the empty set. The reason for this is that when a condition doesn’t contain any “x” in it, then it will either be true regardless of what value we assign to “x”, or it will be false regardless of what value we assign to “x”.

Q7: When a set is given with a complicated specification, I am not always sure how to figure out which individuals are in it and which ones aren’t. I know how to do it in simpler cases. For example, when the set is specified as “\{x : x + 2 = x^2\}”, and I want to know whether, say, the number 29 is in it, I know what I have to do: I have to replace all occurrences of “x” in the condition that follows the colon by occurrences of “29”, and then decide whether the resulting statement about 29 is true or false. In this case, I get the statement “29 + 2 = 29^2”; and since this is false, 29 is not in the set. But there are cases where it’s not so easy. For example, suppose a set is specified as “\{x : x ∈ \{x : x ≠ 0\}\}”, and I want to figure out whether 29 is in this one. So I try replacing “x” with “29” on the right side of the colon. What I get is “29 ∈ \{29 : 29 ≠ 0\}”. But I don’t understand this. We just learned that names can’t occur to the left of the colon; only place-holders make sense there. This looks just like the example “[California : California is a western state]” that I brought up in Q5. So I am stuck. Where did I go wrong?
A7: You went wrong when you replaced all the "x" by "29" and thereby went from "\{x : x \in \{x : x \neq 0\}\}" to "\{29 : 29 \neq 0\}". The former makes sense, the latter doesn't (as you just noted yourself); so this cannot have been an equivalent reformulation.

Q8: Wait a minute, how was I actually supposed to know that "\{x : x \in \{x : x \neq 0\}\}" made sense? For all I knew, this could have been an incoherent definition in the first place, and my reformulation just made it more transparent what was wrong with it.

A8: Here is one way to see that the original description was coherent, and this will also show you how to answer your original question: namely, whether \(29 \in \{x : x \in \{x : x \neq 0\}\}\). First, look only at the most embedded set description, namely "\{x : x \neq 0\}". This transparently describes the set of all objects distinct from 0. We can refer to this set in various other ways: for instance, in the way I just did (as "the set of all objects distinct from 0"), or by a new name that we especially define for it, say as "\(S := \{x : x \neq 0\}\)". Given that the set \(\{x : x \neq 0\}\) can be referred to in all these different ways, we can also express the condition "\(x \in \{x : x \neq 0\}\)" in many different, but equivalent, forms — for example, these three:

"\(x \in \text{the set of all objects distinct from 0}\)"
"\(x \in S \text{ (where } S \text{ is as defined above)}\)"
"\(x \in \{y : y \neq 0\}\)"

Each of these is fulfilled by exactly the same values for "x" as the original condition "\(x \in \{x : x \neq 0\}\)". This, in turn, means that each can be substituted for "\(x \in \{x : x \neq 0\}\)" in "\(\{x : x \in \{x : x \neq 0\}\}\)". So we have:

\[
\{x : x \in \{x : x \neq 0\}\} = \{x : x \in \text{the set of all objects distinct from 0}\} = \{x : x \in S \text{ (where } S \text{ is as defined above)}\} = \{x : x \in \{y : y \neq 0\}\}.
\]

Now if we want to determine whether 29 is a member of \(\{x : x \in \{x : x \neq 0\}\}\), we can do this by using any of the alternative descriptions of this set. Suppose we take the third one above. So we ask whether \(29 \in \{x : x \in S\}\). We know that it is iff \(29 \in S\). By the definition of \(S\), the latter holds iff \(29 \in \{x : x \neq 0\}\). And this in turn is the case iff \(29 \neq 0\). Now we have arrived at an obviously true statement, and we can work our way back and conclude, first, that \(29 \in S\), second, that \(29 \in \{x : x \in S\}\), and third, that \(29 \in \{x : x \in \{x : x \neq 0\}\}\).
Q9: I see for this particular case now that it was a mistake to replace all occurrences of "x" in the condition "x ∈ {x : x ≠ 0}" by "29". But I am still not confident that I wouldn't make a similar mistake in another case. Is there a general rule or fool-proof strategy that I can follow so that I'll be sure to avoid such illegal substitutions?

A9: A very good policy is to write (or rewrite) your conditions in such a way that there is no temptation for illegal substitutions in the first place. This means that you should never reuse the same letter unless this is strictly necessary in order to express what you want to say. Otherwise, use new letters wherever possible. If you follow this strategy, you won't ever write something like "(x : x ∈ {x : x ≠ 0})" to begin with, and if you happen to read it, you will quickly rewrite it before doing anything else with it. What you would write instead would be something like "(x : x ∈ {y : y ≠ 0})". This (as we already noted) describes exactly the same set, but uses distinct letters "x" and "y" instead of only "x"s. It still uses each letter twice, but this, of course, is crucial to what it is meant to express. If we insisted on replacing the second "x" by a "z", for instance, we would wind up with one of those strange descriptions in which the "x" doesn't occur to the right of the colon at all, that is, "(x : z ∈ {y : y ≠ 0})". As we saw earlier, sets described in this way contain either everything or nothing. Besides, what is "z" supposed to stand for? It doesn't seem to be a place-holder, because it's not introduced anywhere to the left of a colon. So it ought to be a name. But whatever it is a name of, that thing was not referred to anywhere in the condition that we had before changing "x" to "z", so we have clearly altered its meaning.

Exercise

The same set can be described in many different ways, often quite different superficially. Here you are supposed to figure out which of the following equalities hold and which ones don't. Sometimes the right answer is not just plain "yes" or "no", but something like "yes, but only if . . . " For example, the two sets in (i) are equal only in the special case where a = b. In case of doubt, the best way to check whether two sets are equal is to consider an arbitrary individual, say John, and to ask if John could be in one of the sets without being in the other as well.

(i) {a} = {b}
(ii) {x : x = a} = {a}
(iii) \( \{x : x \text{ is green}\} = \{y : y \text{ is green}\} \)
(iv) \( \{x : x \text{ likes a}\} = \{y : y \text{ likes b}\} \)
(v) \( \{x : x \in A\} = A \)
(vi) \( \{x : x \in \{y : y \in B\}\} = B \)
(vii) \( \{x : \{y : y \text{ likes } x\}\} = \emptyset \) = \( \{x : \{x : x \text{ likes } x\}\} = \emptyset \)

### 1.3.3 Functions

If we have two objects \( x \) and \( y \) (not necessarily distinct), we can construct from them the ordered pair \( \langle x, y \rangle \). \( \langle x, y \rangle \) must not be confused with \( \{x, y\} \). Since sets with the same members are identical, we always have \( \{x, y\} = \{y, x\} \). But in an ordered pair, the order matters: except in the special case of \( x = y \), \( \langle x, y \rangle \neq \langle y, x \rangle \). ¹⁰

A (2-place) relation is a set of ordered pairs. Functions are a special kind of relation. Roughly speaking, in a function (as opposed to a non-functional relation), the second member of each pair is uniquely determined by the first. Here is the definition:

(3) A relation \( f \) is a function iff it satisfies the following condition:
   
   For any \( x \): if there are \( y \) and \( z \) such that \( \langle x, y \rangle \in f \) and \( \langle x, z \rangle \in f \), then \( y = z \).

Each function has a domain and a range, which are the sets defined as follows:

(4) Let \( f \) be a function.

   Then the domain of \( f \) is \( \{x : \text{there is a } y \text{ such that } \langle x, y \rangle \in f\} \), and the range of \( f \) is \( \{x : \text{there is a } y \text{ such that } \langle y, x \rangle \in f\} \).

When \( A \) is the domain and \( B \) the range of \( f \), we also say that \( f \) is from \( A \) and onto \( B \). If \( C \) is a superset of \( f \)'s range, we say that \( f \) is into (or to) \( C \). For “\( f \) is from \( A \) (in)to \( B \)”, we write “\( f : A \rightarrow B \)”. The uniqueness condition built into the definition of functionhood ensures that whenever \( f \) is a function and \( x \) an element of its domain, the following definition makes sense:

(5) \( f(x) := \text{the unique } y \text{ such that } \langle x, y \rangle \in f \).

For “\( f(x) \)”, read “\( f \) applied to \( x \)” or “\( f \) of \( x \)”. \( f(x) \) is also called the “value of \( f \) for the argument \( x \)”, and we say that \( f \) maps \( x \) to \( y \). “\( f(x) = y \)” (provided that it is well-defined at all) means the same thing as “\( \langle x, y \rangle \in f \)” and is normally the preferred notation.
Functions, like sets, can be defined in various ways, and the most straightforward one is again to simply list the function’s elements. Since functions are sets of ordered pairs, this can be done with the notational devices we have already introduced, as in (6), or else in the form of a table like the one in (7), or in words such as (8).

(6) \[ F := \{<a, b>, <c, b>, <d, e>\} \]

(7) \[
F :=
\begin{bmatrix}
a \rightarrow b \\
c \rightarrow b \\
d \rightarrow e
\end{bmatrix}
\]

(8) Let \( F \) be that function \( f \) with domain \( \{a, c, d\} \) such that \( f(a) = f(c) = b \) and \( f(d) = e \).

Each of these definitions determines the same function \( F \). The convention for reading tables like the one in (7) is transparent: the left column lists the domain and the right column the range, and an arrow points from each argument to the value it is mapped to.

Functions with large or infinite domains are often defined by specifying a condition that is to be met by each argument-value pair. Here is an example.

(9) Let \( F_{+1} \) be that function \( f \) such that
\[
f : \mathbb{N} \rightarrow \mathbb{N}, \text{ and for every } x \in \mathbb{N}, f(x) = x + 1.\]
\((\mathbb{N} \text{ is the set of all natural numbers.})\)

The following is a slightly more concise format for this sort of definition:

(10) \[ F_{+1} := f : \mathbb{N} \rightarrow \mathbb{N} \]
\[
\text{For every } x \in \mathbb{N}, f(x) = x + 1.
\]

Read (10) as: “\( F_{+1} \) is to be that function \( f \) from \( \mathbb{N} \) into \( \mathbb{N} \) such that, for every \( x \in \mathbb{N}, f(x) = x + 1. \)” An even more concise notation (using the \( \lambda \)-operator) will be introduced at the end of the next chapter.

Notes


7 This was not true of Frege. He distinguished between the function itself and its extension (German: *Wertverlauf*). The latter, however, is precisely what mathematicians today call a “function”, and they have no use for another concept that would correspond to Frege’s notion of a function. Some of Frege’s commentators have actually questioned whether that notion was coherent. To him, though, the distinction was very important, and he maintained that while a function is unsaturated, its extension is something saturated. So we are clearly going against his stated intentions here.

8 “Iff” is the customary abbreviation for “if and only if”.

9 We use the colon in front of the equality sign to indicate that an equality holds by definition. More specifically, we use it when we are defining the term to the left of “:=” in terms of the one to the right. In such cases, we should always have a previously unused symbol on the left, and only familiar and previously defined material on the right. In practice, of course, we will reuse the same letters over and over, but whenever a letter appears to the left of “:=”, we thereby cancel any meaning that we may have assigned it before.

10 It is possible to define ordered pairs in terms of sets, for instance as follows: \( <x, y> := \{\{x\}, \{x, y\} \} \). For most applications of the concept (the ones in this book included), however, you don’t need to know this definition.

11 The superset relation is the inverse of the subset relation: \( A \) is a superset of \( B \) iff \( B \subseteq A \).
2 Executing the Fregean Program

In the pages to follow, we will execute the Fregean program for a fragment of English. Although we will stay very close to Frege's proposals at the beginning, we are not interested in an exegesis of Frege, but in the systematic development of a semantic theory for natural language. Once we get beyond the most basic cases, there will be many small and some not-so-small departures from the semantic analyses that Frege actually defended. But his treatment of semantic composition as functional application (Frege's Conjecture), will remain a leading idea throughout.

Modern syntactic theory has taught us how to think about sentences and their parts. Sentences are represented as phrase structure trees. The parts of a sentence are subtrees of phrase structure trees. In this chapter, we begin to explore ways of interpreting phrase structure trees of the kind familiar in linguistics. We will proceed slowly. Our first fragment of English will be limited to simple intransitive and transitive sentences (with only proper names as subjects and objects), and extremely naive assumptions will be made about their structures. Our main concern will be with the process of meaning composition. We will see how a precise characterization of this process depends on, and in turn constrains, what we say about the interpretation of individual words.

This chapter, too, has sections which are not devoted to semantics proper, but to the mathematical tools on which this discipline relies. Depending on the reader's prior mathematical experience, these may be supplemented by exercises from other sources or skimmed for a quick review.

2.1 First example of a Fregean interpretation

We begin by limiting our attention to sentences that consist of a proper name plus an intransitive verb. Let us assume that the syntax of English associates these with phrase structures like that in (1).
We want to formulate a set of semantic rules which will provide denotations for all trees and subtrees in this kind of structure. How shall we go about this? What sorts of entities shall we employ as denotations? Let us be guided by Frege.

Frege took the denotations of sentences to be truth-values, and we will follow him in this respect. But wait. Can this be right? The previous chapter began with the statement "To know the meaning of a sentence is to know its truth-conditions". We emphasized that the meaning of a sentence is not its actual truth-value, and concluded that a theory of meaning for natural language should pair sentences with their truth-conditions and explain how this can be done in a compositional way. Why, then, are we proposing truth-values as the denotations for sentences? Bear with us. Once we spell out the complete proposal, you'll see that we will end up with truth-conditions after all.

The Fregean denotations that we are in the midst of introducing are also called "extensions", a term of art which is often safer to use because it has no potentially interfering non-technical usage. The extension of a sentence, then, is its actual truth-value. What are truth-values? Let us identify them with the numbers 1 (True) and 0 (False). Since the extensions of sentences are not functions, they are saturated in Frege's sense. The extensions of proper names like "Ann" and "Jan" don't seem to be functions either. "Ann" denotes Ann, and "Jan" denotes Jan.

We are now ready to think about suitable extensions for intransitive verbs like "smokes". Look at the above tree. We saw that the extension for the lexical item "Ann" is the individual Ann. The node dominating "Ann" is a non-branching N-node. This means that it should inherit the denotation of its daughter node. The N-node is again dominated by a non-branching node. This NP-node, then, will inherit its denotation from the N-node. So the denotation of the NP-node in the above tree is the individual Ann. The NP-node is dominated by a branching S-node. The denotation of the S-node, then, is calculated from the denotation of the NP-node and the denotation of the VP-node. We know that the denotation of the NP-node is Ann, hence saturated. Recall now that Frege conjectured that all semantic composition amounts to functional application. If that is so, we
must conclude that the denotation of the VP-node must be unsaturated, hence a function. What kind of function? Well, we know what kinds of things its arguments and its values are. Its arguments are individuals like Ann, and its values are truth-values. The extension of an intransitive verb like "smokes", then, should be a function from individuals to truth-values.

Let's put this all together in an explicit formulation. Our semantics for the fragment of English under consideration consists of three components. First, we define our inventory of denotations. Second, we provide a lexicon which specifies the denotation of each item that may occupy a terminal node. Third, we give a semantic rule for each possible type of non-terminal node. When we want to talk about the denotation of a lexical item or tree, we enclose it in double brackets. For any expression $\alpha$, then, $[[\alpha]]$ is the denotation of $\alpha$. We can think of $[ ]$ as a function (the interpretation function) that assigns appropriate denotations to linguistic expressions. In this and most of the following chapters, the denotations of expressions are extensions. The resulting semantic system is an extensional semantics. Towards the end of this book, we will encounter phenomena that cannot be handled within an extensional semantics. We will then revise our system of denotations and introduce intensions.

A. Inventory of denotations
Let $D$ be the set of all individuals that exist in the real world. Possible denotations are:

- Elements of $D$, the set of actual individuals.
- Elements of $\{0, 1\}$, the set of truth-values.
- Functions from $D$ to $\{0, 1\}$.

B. Lexicon

$[[\text{Ann}]] = \text{Ann}$
$[[\text{Jan}]] = \text{Jan}$
etc. for other proper names.

$[[\text{works}]] = f : D \to \{0, 1\}$
For all $x \in D$, $f(x) = 1$ iff $x$ works.

$[[\text{smokes}]] = f : D \to \{0, 1\}$
For all $x \in D$, $f(x) = 1$ iff $x$ smokes.
etc. for other intransitive verbs.

C. Rules for non-terminal nodes
In what follows, Greek letters are used as variables for trees and subtrees.
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Executing the Fregean Program

(S1) If $\alpha$ has the form $S$, then $[[\alpha]] = [[\gamma]]([[\beta]])$.

(S2) If $\alpha$ has the form $NP$, then $[[\alpha]] = [[\beta]]$.

(S3) If $\alpha$ has the form $VP$, then $[[\alpha]] = [[\beta]]$.

(S4) If $\alpha$ has the form $N$, then $[[\alpha]] = [[\beta]]$.

(S5) If $\alpha$ has the form $V$, then $[[\alpha]] = [[\beta]]$.

2.1.1 Applying the semantics to an example

Does this set of semantic rules predict the correct truth-conditions for "Ann smokes"? That is, is "Ann smokes" predicted to be true if and only if Ann smokes? "Of course", you will say, "that's obvious". It's pretty obvious indeed, but we are still going to take the trouble to give an explicit proof of it. As matters get more complex in the chapters to come, it will be less and less obvious whether a given set of proposed rules predict the judgments it is supposed to predict. But you can always find out for sure if you draw your trees and work through them node by node, applying one rule at a time. It is best to get used to this while the calculations are still simple. If you have some experience with computations of this kind, you may skip this subsection.

We begin with a precise statement of the claim we want to prove:

Claim:

$$[[S \quad NP \quad VP \quad N \quad V \quad Ann \quad smokes]] = 1 \text{ iff Ann smokes.}$$
We want to deduce this claim from our lexical entries and semantic rules (S1)–(S5). Each of these rules refers to trees of a certain form. The tree

\[
S \rightarrow \text{NP} \quad \text{VP} \\
\quad \text{N} \quad \text{V} \\
\text{Ann} \quad \text{smokes}
\]

is of the form specified by rule (S1), repeated here, so let's see what (S1) says about it.

(S1) If \( \alpha \) has the form \( S \), then \( [\alpha] = [\gamma]\([\beta]\)\).

When we apply a general rule to a concrete tree, we must first match up the variables in the rule with the particular constituents that correspond to them in the application. In this instance, \( \alpha \) is

\[
S \rightarrow \text{NP} \quad \text{VP} \\
\quad \text{N} \quad \text{V} \\
\text{Ann} \quad \text{smokes}
\]

so \( \beta \) must be \( \text{NP} \) and \( \gamma \) must be \( \text{VP} \).
The rule says that \([\alpha] = [\gamma]\([\beta]\)]\), so this means in the present application that

\[
(2) \quad S \quad \begin{array}{c}
\text{NP} \\
\text{V}
\end{array} \quad \begin{array}{c}
\text{VP} \\
\text{N}
\end{array} \quad \begin{array}{c}
\text{smokes} \\
\text{Ann}
\end{array} = \begin{array}{c}
\text{VP} \\
\text{V}
\end{array} \quad \begin{array}{c}
\text{NP} \\
\text{N}
\end{array} \quad \begin{array}{c}
\text{smokes} \\
\text{Ann}
\end{array}
\]

Now we apply rule (S3) to the tree

\[
\text{VP} \quad \begin{array}{c}
\text{V}
\end{array} \quad \text{smokes}
\]

(This time, we skip the detailed justification of why and how this rule fits this tree.) What we obtain from this is

\[
(3) \quad \begin{array}{c}
\text{VP} \\
\text{V}
\end{array} = \begin{array}{c}
\text{V}
\end{array} \quad \text{smokes}
\]

From (2) and (3), by substituting equals for equals, we infer (4).

\[
(4) \quad S \quad \begin{array}{c}
\text{NP} \\
\text{V}
\end{array} \quad \begin{array}{c}
\text{VP} \\
\text{N}
\end{array} \quad \begin{array}{c}
\text{smokes} \\
\text{Ann}
\end{array} = \begin{array}{c}
\text{VP} \\
\text{V}
\end{array} \quad \begin{array}{c}
\text{NP} \\
\text{N}
\end{array} \quad \begin{array}{c}
\text{smokes} \\
\text{Ann}
\end{array}
\]

Now we apply rule (S5) to the appropriate subtree and use the resulting equation for another substitution in (4):
(5) \[
\begin{array}{c}
S \\
\hline
\begin{array}{c}
NP \\
N \\
\hline
Ann \\
\end{array} \\
\begin{array}{c}
VP \\
V \\
\hline
\text{smokes} \\
\end{array}
\end{array}
= \begin{array}{c}
\text{[smokes]} \\
\hline
\begin{array}{c}
NP \\
N \\
\hline
\text{Ann} \\
\end{array}
\end{array}
\]

Now we use rule (S2) and then (S4), and after substituting the results thereof in (5), we have (6).

(6) \[
\begin{array}{c}
S \\
\hline
\begin{array}{c}
NP \\
N \\
\hline
Ann \\
\end{array} \\
\begin{array}{c}
VP \\
V \\
\hline
\text{smokes} \\
\end{array}
\end{array}
= \begin{array}{c}
\text{[smokes]}([\text{Ann}]) \\
\end{array}
\]

At this point, we look up the lexical entries for Ann and smokes. If we just use these to substitute equals for equals in (6), we get (7).

(7) \[
\begin{array}{c}
S \\
\hline
\begin{array}{c}
NP \\
N \\
\hline
\text{Ann} \\
\end{array} \\
\begin{array}{c}
VP \\
V_i \\
\hline
\text{smokes} \\
\end{array}
\end{array}
= \begin{array}{c}
\text{[f : D} ightarrow \{0, 1\} \\
\text{For all x ∈ D, f(x) = 1 iff x smokes]} \\
\end{array}(\text{Ann})
\]

Let's take a close look at the right-hand side of this equation. It has the gross form "function (argument)", so it denotes the value that a certain function yields for a certain argument. The argument is Ann, and the function is the one which maps those who smoke to 1 and all others to 0. If we apply this function to Ann, we will get 1 if Ann smokes and 0 if she doesn't. To summarize what we have just determined:
executing the Fregean Program

\[ f : D \rightarrow \{0, 1\} \]

For all \( x \in D, f(x) = 1 \text{ iff } x \text{ smokes} \]

\( (Ann) = 1 \text{ iff Ann smokes.} \)

And now we have reached the goal of our proof: (7) and (8) together imply exactly the claim which we stated at the beginning. QED.

This was not the only way in which we could have constructed the proof of this claim. What matters is (a) that we use each applicable rule or lexical entry to obtain an equation regarding the denotation of a certain subtree; (b) that we keep using some of these equations to substitute equals for equals in others, thereby getting closer and closer to the target equation in our claim; and (c) that we employ the definitions of functions that we find in the lexicon to calculate their values for specified arguments. There is no unique specified order in which we must perform these steps. We can apply rules to the smallest subtrees first, or start at the top of the tree, or anywhere in the middle. We can collect a long list of separate equations before we begin to draw conclusions from any two of them, or else we can keep alternating applications of semantic rules with substitutions in equations derived previously. The soundness of the proof is not affected by these choices (although, of course, some strategies may be easier than others to follow through without getting confused).

We have used the word “proof” a number of times in this section. What exactly do we mean by this term? The notion of “proof” has been made precise in various ways in the history of logic. The most rigorous notion equates a proof with a syntactic derivation in an axiomatic or natural deduction system. Above, we relied on a less regimented notion of “proof” that is common in mathematics. Mathematical proofs are rarely algorithmic derivations. They are usually written in plain English (or some other natural language), supplemented by technical vocabulary that has been introduced through definitions. Conclusions are licensed by inference patterns that are known to be valid but are not spelled out formally. The proofs in this book are all “semi-formal” in this way. The standards of rigor followed in mathematics should be good enough for what we want to accomplish here.

2.1.2 Deriving truth-conditions in an extensional semantics

The proof we just gave illustrates how a semantic system based on extensions allows us to compute the truth-conditions, and hence the meaning, of a sentence. If you check the proof again, you will see that we end up with the truth-conditions of “Ann smokes” because the lexicon defines the extensions of predicates by specifying a condition. Had we defined the function denoted by “smokes” by displaying it in a table, for example, we would have obtained a
mere truth-value. We didn’t really have a choice, though, because displaying the function in a table would have required more world knowledge than we happen to have. We do not know of every existing individual whether or not (s)he smokes. And that’s certainly not what we have to know in order to know the meaning of “smoke”. We could look at a fictitious example, though.

Suppose Ann, Jan, and Maria are the only individuals in the actual world, and Ann and Jan are the only smokers. The extension of the verb “smokes” can now be displayed in a table:

\[
\text{[[smokes]]} = \begin{bmatrix}
\text{Ann} & \rightarrow & 1 \\
\text{Jan} & \rightarrow & 1 \\
\text{Maria} & \rightarrow & 0
\end{bmatrix}
\]

Using this way of defining the extension of “smokes”, our computation would have ended as follows:

\[
\begin{array}{c}
\text{NP} \\
\text{N} \\
\text{Ann}
\end{array}
\quad \begin{array}{c}
\text{VP} \\
\text{V} \\
\text{smokes}
\end{array}
\quad =
\begin{bmatrix}
\text{Ann} & \rightarrow & 1 \\
\text{Ann} & \rightarrow & 1 \\
\text{Maria} & \rightarrow & 0
\end{bmatrix}
\]

Here, the sentence “Ann smokes” would not be paired with its truth-conditions, but with the value 1.

The issue of how an extensional system can yield a theory of meaning concerns the relationship between what Frege called “Sinn” and “Bedeutung”. Frege’s “Bedeutung” corresponds to our term “extension”, and is sometimes translated as “reference”. Frege’s “Sinn” is usually translated as “sense”, and corresponds to what we have called “meaning”. How does Frege get us from Bedeutung to Sinn? In his book on Frege’s philosophy of language, Michael Dummett answers this question as follows:

It has become a standard complaint that Frege talks a great deal about the senses of expressions, but nowhere gives an account of what constitutes such a sense. This complaint is partly unfair: for Frege the sense of an expression is the manner in which we determine its reference, and he tells us a great deal about the kind of reference possessed by expressions of different types, thereby specifying the form that the senses of such expressions must
The sense of an expression is the mode of presentation of the referent: in saying what the referent is, we have to choose a particular way of saying this, a particular means of determining something as a referent.\textsuperscript{3}

What Dummett says in this passage is that when specifying the extension (reference, \textit{Bedeutung}) of an expression, we have to choose a particular way of presenting it, and it is this manner of presentation that might be considered the meaning (sense, \textit{Sinn}) of the expression. The function that is the extension of a predicate can be presented by providing a condition or by displaying it in a table, for example. Only if we provide a condition do we choose a mode of presentation that "shows"\textsuperscript{4} the meaning of the predicates and the sentences they occur in. Different ways of defining the same extensions, then, can make a theoretical difference. Not all choices yield a theory that pairs sentences with their truth-conditions. Hence not all choices lead to a theory of meaning.

\subsection{Object language and metalanguage}

Before we conclude this section, let us briefly reflect on a typographical convention that we have already been using. When we referred to words and phrases of English (represented as strings or trees), we replaced the customary quotes by bold-face. So we had, for example:


```
NP    VP    NP    VP
\text{Ann} \quad \text{smokes} = \quad \text{Ann} \quad \text{smokes}
```

The expressions that are bold-faced or enclosed in quotes are expressions of our \textit{object language}, the language we are investigating. In this book, the object language is English, since we are developing a semantic theory for English. The language we use for theoretical statements is the \textit{metalanguage}. Given that this book is written in English, our metalanguage is English as well. Since we are looking at the English object language in a fairly technical way, our English metalanguage includes a fair amount of technical vocabulary and notational
conventions. The abstraction notation for sets that we introduced earlier is an example. Quotes or typographical distinctions help us mark the distinction between object language and metalanguage. Above, we always used the bold-faced forms when we placed object language expressions between denotation brackets. For example, instead of writing the lexical entry for the name “Ann”:

\[
\text{["Ann"]} = \text{Ann}
\]

we wrote:

\[
[\text{Ann}] = \text{Ann}.
\]

This lexical entry determines that the denotation of the English name “Ann” is the person Ann. The distinction between expressions and their denotations is important, so we will usually use some notational device to indicate the difference. We will never write things like “[Ann]”. This would have to be read as “the denotation of (the person) Ann”, and thus is nonsense. And we will also avoid using bold-face for purposes other than to replace quotes (such as emphasis, for which we use italics).

---

**Exercise on sentence connectives**

Suppose we extend our fragment to include phrase structures of the forms below (where the embedded S-constituents may either belong to the initial fragment or have one of these three forms themselves):

```
S
---
S

S
S S S
And
\Delta
\Delta
\Delta

S
S S S
Or
\Delta
\Delta
\Delta
```

How do we have to revise and extend the semantic component in order to provide all the phrase structures in this expanded fragment with interpretations? Your task in this exercise is to define an appropriate semantic value for each new lexical item (treat “it-is-not-the-case-that” as a single lexical item here) and to write appropriate semantic rules for the new types of non-terminal nodes. To do this, you will also have to expand the inventory of possible semantic values. Make sure that you stick to our working hypothesis that all semantic composition is functional application (Frege’s Conjecture).
2.2 Sets and their characteristic functions

We have construed the denotations of intransitive verbs as functions from individuals to truth-values. Alternatively, they are often regarded as sets of individuals. This is the standard choice for the extensions of 1-place predicates in logic. The intuition here is that each verb denotes the set of those things that it is true of. For example: \( \text{[sleep]} = \{ x \in D : x \text{ sleeps} \} \). This type of denotation would require a different semantic rule for composing subject and predicate, one that isn’t simply functional application.

---

**Exercise**

Write the rule it would require.

Here we have chosen to take Frege’s Conjecture quite literally, and have avoided sets of individuals as denotations for intransitive verbs. But for some purposes, sets are easier to manipulate intuitively, and it is therefore useful to be able to pretend in informal talk that intransitive verbs denote sets. Fortunately, this make-believe is harmless, because there exists a one-to-one correspondence between sets and certain functions.

1. Let \( A \) be a set. Then \( \text{char}_A \), the characteristic function of \( A \), is that function \( f \) such that, for any \( x \in A \), \( f(x) = 1 \), and for any \( x \notin A \), \( f(x) = 0 \).

2. Let \( f \) be a function with range \( \{0, 1\} \). Then \( \text{char}_f \), the set characterized by \( f \), is \( \{ x \in D : f(x) = 1 \} \).

Exploiting the correspondence between sets and their characteristic functions, we will often switch back and forth between function talk and set talk in the discussion below, sometimes saying things that are literally false, but become true when the references to sets are replaced by references to their characteristic functions (or vice versa).

Here is an illustration: Suppose our universe consists of three individuals, \( D = \{ \text{Ann, Jan, Maria} \} \). Suppose further that Ann and Jan are the ones who sleep, and Ann is the only one who snores. If we treat intransitive verbs as denoting sets, we may then assign the following denotations to sleep and snore:

3. \( \text{[sleep]} = \{ \text{Ann, Jan} \} \).

4. \( \text{[snore]} = \{ \text{Ann} \} \).
We can now write things like the following:

(5)  \( \text{Ann} \in [\text{sleep}] \).

(6)  \([\text{snore}] \subseteq [\text{sleep}] \).

(7)  \(|[\text{snore}] \cap [\text{sleep}]| = 1 \).

\(|A| \) (the \textit{cardinality} of \( A \)) is the number of elements in the set \( A \).

(5) means that Ann is among the sleepers, (6) means that the snorers are a subset of the sleepers, and (7) means that the intersection of the snorers and the sleepers has exactly one element. All these are true, given (3) and (4). Now suppose we want to switch to a treatment under which intransitive verbs denote characteristic functions instead of the corresponding sets.

(3')  \([\text{sleep}] = \begin{bmatrix} \text{Ann} & \rightarrow & 1 \\ \text{Jan} & \rightarrow & 1 \\ \text{Maria} & \rightarrow & 0 \end{bmatrix} \)

(4')  \([\text{snore}] = \begin{bmatrix} \text{Ann} & \rightarrow & 1 \\ \text{Jan} & \rightarrow & 0 \\ \text{Maria} & \rightarrow & 0 \end{bmatrix} \)

If we want to make statements with the same import as (5)–(7) above, we can no longer use the same formulations. For instance, the statement

\( \text{Ann} \in [\text{sleep}] \)

if we read it literally, is now false. According to (3'), \([\text{sleep}] \) is a function. Functions are sets of ordered pairs, in particular,

\([\text{sleep}] = \{<\text{Ann}, 1>, <\text{Jan}, 1>, <\text{Maria}, 0>\} \).

Ann, who is not an ordered pair, is clearly not among the elements of this set. Likewise,

\([\text{snore}] \subseteq [\text{sleep}] \)

is now false, because there is one element of \([\text{snore}] \), namely the pair \(<\text{Jan}, 0>\), which is not an element of \([\text{sleep}] \). And
\[|\text{\texttt{snore}} \cap \text{\texttt{sleep}}| = 1\]

is false as well, because the intersection of the two functions described in (3') and (4') contains not just one element, but two, namely \(<\text{Ann}, 1>\) and \(<\text{Maria, 0}>\).

The upshot of all this is that, once we adopt (3') and (4') instead of (3) and (4), we have to express ourselves differently if we want to make statements that preserve the intuitive meaning of our original (5)-(7). Here is what we have to write instead:

\[(5')\] \[\text{\texttt{sleep}}(\text{Ann}) = 1\]

\[(6')\] For all \(x \in D : \text{if} \ [\text{\texttt{snore}}](x) = 1, \text{then} \ [\text{\texttt{sleep}}](x) = 1\]

Or, equivalently:
\[\{x : [\text{\texttt{snore}}](x) = 1\} \subseteq \{x : [\text{\texttt{sleep}}](x) = 1\}\]

\[(7')\] \[|\{x : [\text{\texttt{snore}}](x) = 1\} \cap \{x : [\text{\texttt{sleep}}](x) = 1\}| = 1\]

As you can see from this, using characteristic functions instead of sets makes certain things a little more cumbersome.

### 2.3 Adding transitive verbs: semantic types and denotation domains

Our next goal is to extend our semantics to simple transitive clauses like “Ann likes Jan”. We take it that the structures that the syntax assigns to these look like this:

```
S
   NP  VP
      N  V  NP
           Ann  likes  N
                             Jan
```
The transitive verb combines with its direct object to form a VP, and the VP combines with the subject to form a sentence. Given structures of this kind, what is a suitable denotation for transitive verbs? Look at the above tree. The lexical item “likes” is dominated by a non-branching V-node. The denotation of “likes”, then, is passed up to this node. The next node up is a branching VP-node. Assuming that semantic interpretation is local, the denotation of this VP-node must be composed from the denotations of its two daughter nodes. If Frege’s Conjecture is right, this composition process amounts to functional application. We have seen before that the denotations of NP-nodes dominating proper names are individuals. And that the denotation of VP-nodes are functions from individuals to truth-values. This means that the denotation of a transitive verb like “likes” is a \textit{function from individuals to functions from individuals to truth-values}.\footnote{When we define such a function-valued function in full explicitness, we have to nest one definition of a function inside another. Here is the proposed meaning of “likes”.

\[
[\text{like}] = f : D \to \{ g : g \text{ is a function from } D \to \{0, 1\} \}
\]
\[
\text{For all } x \in D, f(x) = g_x : D \to \{0, 1\}
\]
\[
\text{For all } y \in D, g_x(y) = 1 \text{ iff } y \text{ likes } x.
\]

This reads: \([\text{like}]\) is that function \(f\) from \(D\) into the set of functions from \(D\) to \((0, 1)\) such that, for all \(x \in D\), \(f(x)\) is that function \(g_x\) from \(D\) into \((0, 1)\) such that, for all \(y \in D\), \(g_x(y) = 1 \) iff \(y\) likes \(x\). Fortunately, this definition can be compressed a bit. There is no information lost in the following reformulation: \footnote{\([\text{like}]\) is a 1-place function, that is, a function with just one argument. The arguments of \([\text{like}]\) are interpreted as the individuals which are liked; that is, they correspond to the grammatical \textit{object} of the verb “like”. This is so because we have assumed that transitive verbs form a VP with their direct object. The direct object, then, is the argument that is closest to a transitive verb, and is therefore semantically processed before the subject. What is left to spell out are the interpretations for the new kinds of non-terminal nodes:

\textbf{(S6)} \text{ If } \alpha \text{ has the form } \wedge \beta \gamma, \text{ then } [\alpha] = [\beta][\gamma].
}
What we just proposed implies an addition to our inventory of possible denotations: aside from individuals, truth-values, and functions from individuals to truth-values, we now also employ functions from individuals to functions from individuals to truth-values. It is convenient at this point to introduce a way of systematizing and labeling the types of denotations in this growing inventory. Following a common practice in the tradition of Montague, we employ the labels “e” and “t” for the two basic types.9

1. e is the type of individuals.
   \[ D_e := D. \]

2. t is the type of truth-values.
   \[ D_t := \{0, 1\}. \]

Generally, \( D_t \) is the set of possible denotations of type \( \tau \). Besides the basic types \( e \) and \( t \), which correspond to Frege’s saturated denotations, there are derived types for various sorts of functions, Frege’s unsaturated denotations. These are labeled by ordered pairs of simpler types: \( <\sigma, \tau> \) is defined as the type of functions whose arguments are of type \( \sigma \) and whose values are of type \( \tau \). The particular derived types of denotations that we have employed so far are \( <e, t> \) and \( <e, <e, t>> \):

3. \[ D_{<e, t>} := \{ f : f \text{ is a function from } D_e \text{ to } D_t \} \]

4. \[ D_{<e, <e, t>>} := \{ f : f \text{ is a function from } D_e \text{ to } D_{<e, t>} \} \]

Further additions to our type inventory will soon become necessary. Here is a general definition:

5. Semantic types
   (a) \( e \) and \( t \) are semantic types.
   (b) If \( \sigma \) and \( \tau \) are semantic types, then \( <\sigma, \tau> \) is a semantic type.
   (c) Nothing else is a semantic type.

Semantic denotation domains
   (a) \( D_e := D \) (the set of individuals).
   (b) \( D_t := \{0, 1\} \) (the set of truth-values).
   (c) For any semantic types \( \sigma \) and \( \tau \), \( D_{<\sigma, \tau>} \) is the set of all functions from \( D_\sigma \) to \( D_\tau \).

(5) presents a recursive definition of an infinite set of semantic types and a parallel definition of a typed system of denotation domains. Which semantic
types are actually used by natural languages is still a matter of debate. The issue of "type economy" will pop up at various places throughout this book, most notably in connection with adjectives in chapter 4 and quantifier phrases in chapter 7. So far, we have encountered denotations of types $e$, $<e,t>$, and $<e,<e,t>\rangle$ as possible denotations for lexical items: $D_e$ contains the denotations of proper names, $D_{<e,t>\rangle}$ the denotations of intransitive verbs, and $D_{<e,<e,t>\rangle}$ the denotations of transitive verbs. Among the denotations of non-terminal constituents, we have seen examples of four types: $D$ contains the denotations of all Ss, $D_e$ (= $D$) the denotations of all Ns and NPs, $D_{<e,t>\rangle}$ the denotations of all VPs and certain Vs, and $D_{<e,<e,t>\rangle}$ the denotations of the remaining Vs.

2.4 Schönfinkelization

Once more we interrupt the construction of our semantic component in order to clarify some of the underlying mathematics. Our current framework implies that the denotations of transitive verbs are 1-place functions. This follows from three assumptions about the syntax–semantics interface that we have been making:

Binary Branching
In the syntax, transitive verbs combine with the direct object to form a VP, and VPs combine with the subject to form a sentence.

Locality
Semantic interpretation rules are local: the denotation of any non-terminal node is computed from the denotations of its daughter nodes.

Frege's Conjecture
Semantic composition is functional application.

If transitive verbs denote 1-place function-valued functions, then our semantics contrasts with the standard semantics for 2-place predicates in logic, and it is instructive to reflect somewhat systematically on how the two approaches relate to each other.

In logic texts, the extension of a 2-place predicate is usually a set of ordered pairs: that is, a relation in the mathematical sense. Suppose our domain $D$ contains just the three goats Sebastian, Dimitri, and Leopold, and among these, Sebastian is the biggest and Leopold the smallest. The relation "is-bigger-than" is then the following set of ordered pairs:
R_{bigger} = \{<Sebastian, Dimitri>, <Sebastian, Leopold>, <Dimitri, Leopold>\}.

We have seen above that there is a one-to-one correspondence between sets and their characteristic functions. The "functional version" of $R_{bigger}$ is the following function from $D \times D$ to $\{0, 1\}$.

\[
\begin{array}{c}
\langle L, S \rangle \rightarrow 0 \\
\langle L, D \rangle \rightarrow 0 \\
\langle L, L \rangle \rightarrow 0 \\
\langle S, L \rangle \rightarrow 1 \\
\langle S, D \rangle \rightarrow 1 \\
\langle S, S \rangle \rightarrow 0 \\
\langle D, L \rangle \rightarrow 1 \\
\langle D, S \rangle \rightarrow 0 \\
\langle D, D \rangle \rightarrow 0
\end{array}
\]

$f_{bigger}$ is a 2-place function. In his paper "Über die Bausteine der mathematischen Logik," Moses Schönfinkel showed how n-place functions can quite generally be reduced to 1-place functions. Let us apply his method to the 2-place function above. That is, let us Schönfinkel\footnote{\textsuperscript{10}} the function $f_{bigger}$. There are two possibilities. $f'_\text{bigger}$ is the left-to-right Schönfinkelization of $f_{\text{bigger}}$. $f''_\text{bigger}$ is the right-to-left Schönfinkelization of $f_{\text{bigger}}$.

\[
\begin{array}{c}
L \rightarrow \begin{bmatrix} L \rightarrow 0 \\ S \rightarrow 0 \\ D \rightarrow 0 \end{bmatrix} \\
S \rightarrow \begin{bmatrix} L \rightarrow 1 \\ S \rightarrow 0 \\ D \rightarrow 1 \end{bmatrix} \\
D \rightarrow \begin{bmatrix} L \rightarrow 1 \\ S \rightarrow 0 \\ D \rightarrow 0 \end{bmatrix}
\end{array}
\]

$f'_\text{bigger}$ is a function that applies to the first argument of the "bigger" relation to yield a function that applies to the second argument. When applied to Leopold, it yields a function that maps any goat into 1 if it is smaller than Leopold. There is no such goat. Hence we get a constant function that assigns 0 to all the goats. When applied to Sebastian, $f'_\text{bigger}$ yields a function that maps any goat into 1 if it is smaller than Sebastian. There are two such goats, Leopold and Dimitri. And when applied to Dimitri, $f'_\text{bigger}$ yields a function that maps any goat into 1 if it is smaller than Dimitri. There is only one such goat, Leopold.
Executing the Fregean Program

Let \( f''_{\text{bigger}} \) be the function that applies to the second argument of the "bigger" relation to yield a function that applies to the first argument. When applied to Leopold, it yields a function that maps any goat into 1 if it is bigger than Leopold. These are all the goats except Leopold. When applied to Sebastian, \( f''_{\text{bigger}} \) yields a function that maps any goat into 1 if it is bigger than Sebastian. There is no such goat. And when applied to Dimitri, \( f''_{\text{bigger}} \) maps any goat into 1 if it is bigger than Dimitri. There is only one such goat, Sebastian.

On both methods, we end up with nothing but 1-place functions, and this is as desired. This procedure can be generalized to any n-place function. You will get a taste for this by doing the exercises below.

Now we can say how the denotations of 2-place predicates construed as relations are related to the Fregean denotations introduced above. The Fregean denotation of a 2-place predicate is the right-to-left Schönfinkeled version of the characteristic function of the corresponding relation. Why the right-to-left Schönfinkelandization? Because the corresponding relations are customarily specified in such a way that the grammatical object argument of a predicate corresponds to the right component of each pair in the relation, and the subject to the left one. (For instance, by the "love"-relation one ordinarily means the set of lover–loved pairs, in this order, and not the set of loved–lover pairs.) That’s an arbitrary convention, in a way, though suggested by the linear order in which English realizes subjects and objects. As for the Fregean denotations of 2-place predicates, remember that it is not arbitrary that their (unique) argument corresponds to the grammatical object of the predicate. Since the object is closest to the predicate in hierarchical terms, it must provide the argument for the function denoted by the predicate.

Exercise 1

Suppose that our universe \( D \) contains just two elements, Jacob and Maria. Consider now the following binary and ternary relations:

\[
\begin{bmatrix}
L & S & D \\
0 & 1 & 1 \\
L & 0 & 0 \\
S & 0 & 0 \\
D & 0 & 0 \\
\end{bmatrix}
\]

\( f''_{\text{bigger}} \) is a function that applies to the second argument of the "bigger" relation to yield a function that applies to the first argument. When applied to Leopold, it yields a function that maps any goat into 1 if it is bigger than Leopold. These are all the goats except Leopold. When applied to Sebastian, \( f''_{\text{bigger}} \) yields a function that maps any goat into 1 if it is bigger than Sebastian. There is no such goat. And when applied to Dimitri, \( f''_{\text{bigger}} \) maps any goat into 1 if it is bigger than Dimitri. There is only one such goat, Sebastian.

On both methods, we end up with nothing but 1-place functions, and this is as desired. This procedure can be generalized to any n-place function. You will get a taste for this by doing the exercises below.

Now we can say how the denotations of 2-place predicates construed as relations are related to the Fregean denotations introduced above. The Fregean denotation of a 2-place predicate is the right-to-left Schönfinkeled version of the characteristic function of the corresponding relation. Why the right-to-left Schönfinkelandization? Because the corresponding relations are customarily specified in such a way that the grammatical object argument of a predicate corresponds to the right component of each pair in the relation, and the subject to the left one. (For instance, by the “love”-relation one ordinarily means the set of lover–loved pairs, in this order, and not the set of loved–lover pairs.) That’s an arbitrary convention, in a way, though suggested by the linear order in which English realizes subjects and objects. As for the Fregean denotations of 2-place predicates, remember that it is not arbitrary that their (unique) argument corresponds to the grammatical object of the predicate. Since the object is closest to the predicate in hierarchical terms, it must provide the argument for the function denoted by the predicate.

Exercise 1

Suppose that our universe \( D \) contains just two elements, Jacob and Maria. Consider now the following binary and ternary relations:
R_{adores} = \{<\text{Jacob, Maria}>, <\text{Maria, Maria}>\}
R_{assigns to} = \{<\text{Jacob, Jacob, Maria}>, <\text{Maria, Jacob, Maria}>\}

In standard predicate logic, these would be suitable extensions for the 2-place and 3-place predicate letters "F^2" and "G^3" as used in the following scheme of abbreviation:

"F^2": "a adores b"
"G^3": "a assigns b to c"

Find the characteristic functions for both of these relations, and then Schönfinkel them from right to left. Could the two Schönfinkeled functions be suitable denotations for the English verbs "adore" and "assign (to)" respectively? If yes, why? If not, why not?

---

**Exercise 2**

In the exercise on sentence connectives in section 2.1, we stipulated ternary branching structures for sentences with "and" and "or". Now assume that all English phrase structures are at most binary branching, and assign accordingly revised syntactic analyses to these sentences. (Whether you choose right-branching or left-branching structures does not matter here, but stick to one option.) Then revise the semantics accordingly. As always, be sure to provide every subtree with a semantic value, as well as to adhere to our current assumptions about the semantic interpretation component (Locality and Frege's Conjecture).

Using the labeling system introduced at the end of section 2.3, specify the type of denotation for each node in your binary branching structure for the sentence "Jan works, and it is not the case that Jan smokes".

---

**Exercise 3**

(a) Extend the fragment in such a way that phrase structure trees of the following kind are included.
Add the necessary semantic rules and lexical entries, sticking to Locality and Frege’s Conjecture. Assume that the preposition “to” is a semantically vacuous element; that is, assume the ad hoc rule below:

\[
\text{If } \alpha \text{ has the form } P \beta, \text{ then } [\alpha] = [\beta].
\]

(b) Suppose now that the actual world contains just three individuals, Ann, Maria, and Jacob. And suppose further that Ann introduces Maria to Jacob, and Maria introduces Jacob to Ann, and no further introductions take place. Which particular function is \([\text{introduce}]\) in this case? Display it in a table.

(c) Using the table specification of \([\text{introduce}]\) from (b) and the lexical entries for the names, calculate the denotations of each non-terminal node in the tree under (a).

(d) Under standard assumptions, a predicate logic formalization of the English sentence “Ann introduces Maria to Jacob” might look like this:

\[ I^3 (A M J) \]

Scheme of abbreviation:

“\(A\)”: “Ann”

“\(M\)”: “Maria”

“\(J\)”: “Jacob”

“\(I^3\)”: “a introduces b to c”
Executing the Fregean Program

The extension of "lsIt un der this scheme of abbreviation is the following set $X$ of ordered triples:

$$X := \{ <x, y, z> \in D \times D \times D : x \text{ introduces } y \text{ to } z \}.$$  

How is this extension related to the extension — let's call it $f$ — for introduce that you defined in (a)? Give your answer by completing the following statement:

For any $x, y, z \in D$, $f(x)(y)(z) = 1 \text{ iff } \ldots \in X.$

2.5 Defining functions in the $\lambda$-notation

The final section of this chapter is devoted to yet another technical matter. You have already had a taste of the ubiquity of functions among the denotations that our Fregean semantics assigns to the words and phrases of natural languages. We will now introduce another notation for describing functions, which will save us some ink in future chapters.

The format in which we have defined most of our functions so far was introduced in section 1.3 with the following example:

(1) $F_{+1} = f : \mathbb{N} \to \mathbb{N}$  
For every $x \in \mathbb{N}$, $f(x) = x + 1$.

The same definition may henceforth be expressed as follows:

(2) $F_{+1} := \lambda x : x \in \mathbb{N} . x + 1$

The $\lambda$-term, "$\lambda x : x \in \mathbb{N} . x + 1$", is to be read as "the (smallest) function which maps every $x$ such that $x$ is in $\mathbb{N}$ to $x + 1$".

Generally, $\lambda$-terms are constructed according to the following schema:

(3) $[\lambda \alpha : \phi . \gamma]$

We say that $\alpha$ is the argument variable, $\phi$ the domain condition, and $\gamma$ the value description. The domain condition is introduced by a colon, and the value description by a period. $\alpha$ will always be a letter standing for an arbitrary argument of the function we are trying to define. In (2), this is the letter "$x$", which we generally use to stand for arbitrary individuals. The domain condition $\phi$ defines the domain of our function, and it does this by placing a condition on
possible values for $\alpha$. In our example, $\phi$ corresponds to \("x \in \mathbb{N}\), which encodes the information that the domain of $F_{+1}$ contains all and only the natural numbers. The value description $\gamma$, finally, specifies the value that our function assigns to the arbitrary argument represented by $\alpha$. In (2), this reads \("x + 1\), which tells us that the value that $F_{+1}$ assigns to each argument is that argument’s successor. The general convention for reading $\lambda$-terms in (semi-mathematical) English is such that (3) reads as (3\)':

\[(3') \text{ the smallest function which maps every } \alpha \text{ such that } \phi \text{ to } \gamma\]

We will typically omit “smallest”, but it is always understood, and it is strictly speaking necessary to pick out the intended function uniquely. Notice, for instance, that besides $F_{+1}$, there are lots of other functions which also map every natural number to its successor: namely, all those functions which are supersets of $F_{+1}$, but have larger domains than $\mathbb{N}$. By adding “smallest”, we make explicit that the domain condition $\phi$ delimits the domain exactly; that is, that in (2), for instance, \("x \in \mathbb{N}\)” is not only a sufficient, but also a necessary, condition for \("x \in \text{dom}(F_{+1})\).\(^{15}\)

Like other function terms, $\lambda$-terms can be followed by argument terms. So we have:

\[(4) \ [\lambda x : x \in \mathbb{N} . x + 1](5) = 5 + 1 = 6.\]

The $\lambda$-notation as we have just introduced it is not as versatile as the format in (1). It cannot be used to abbreviate descriptions of functions which involve a distinction between two or more cases. To illustrate this limitation, let’s look at the function $G$ defined in (5).

\[(5) \quad G = f : \mathbb{N} \rightarrow \mathbb{N} \]

For every $x \in \mathbb{N}$, $f(x) = 2$, if $x$ is even, and $f(x) = 1$ otherwise.

The problem we encounter if we attempt to press this into the shape \(\[\lambda \alpha : \phi \cdot \gamma\]\) is that there is no suitable value description $\gamma$. Obviously, neither “1” nor “2” is the right choice. \(\[\lambda x : x \in \mathbb{N} . 1\] would be that function which maps every natural number to 1, clearly a different function from the one described in (5). And similarly, of course, \(\[\lambda x : x \in \mathbb{N} . 2\]” would be inappropriate.\(^{16}\)

This implies that, as it stands, the new notation is unsuitable for precisely the kinds of functions that figure most prominently in our semantics. Take the extension of an intransitive verb.

\[(6) \quad \text{[smoke]} = f : \mathbb{D} \rightarrow \{0, 1\} \]

For every $x \in \mathbb{D}$, $f(x) = 1$ iff $x$ smokes.
(6) stipulates that \( f(x) \) is to be 1 if \( x \) smokes and 0 otherwise.\(^{17}\) So we face the same difficulty as with (5) above when it comes to deciding on the value description in a suitable \( \lambda \)-term.

We will get rid of this difficulty by defining an extended use of the \( \lambda \)-notation. So far, you have been instructed to read \( \text{"}[\lambda \alpha : \phi \cdot \gamma]\text{"} \) as \( \text{"the function which maps every } \alpha \text{ such that } \phi \text{ to } \gamma\text{"} \). This paraphrase makes sense only when the value description \( \gamma \) is a noun phrase. If \( \gamma \) had the form of a sentence, for instance, we would get little more than word salad. Try reading out a \( \lambda \)-term like \( \text{"}[\lambda x : x \in \mathbb{N} . x \text{ is even}]\text{"} \). What comes out is: \( \text{"the function which maps every } x \text{ in } \mathbb{N} \text{ to } x \text{ is even}\text{"} \). This is neither colloquial English nor any kind of technical jargon; it just doesn't make any sense.

We could, of course, change the instructions and stipulate a different way to read the notation. Suppose we decided that \( \text{"}[\lambda \alpha : \phi \cdot \gamma]\text{"} \) was to be read as follows: \( \text{"the function which maps every } \alpha \text{ such that } \phi \text{ to } 1 \text{ if } \gamma \text{ and to } 0 \text{ otherwise}\text{"} \). Then, of course, \( \text{"}[\lambda x : x \in \mathbb{N} . x \text{ is even}]\text{"} \), would make perfect sense: \( \text{"the function which maps every } x \text{ in } \mathbb{N} \text{ to } 1, \text{ if } x \text{ is even, and to } 0 \text{ otherwise}\text{"} \). With this new convention in force, we could also use the \( \lambda \)-notation to abbreviate our lexical entry for \( \text{"smoke\text{"}} \).

\[
\text{(7) } [\text{smoke}] := [\lambda x : x \in D . x \text{ smokes}]
\]

If this reads: \( \text{"let } [\text{smoke}] \text{ be the function which maps every } x \text{ in } D \text{ to } 1, \text{ if } x \text{ smokes, and to } 0 \text{ otherwise}\text{"} \), it is easily seen as an equivalent reformulation of (6). If we add an argument term, we have:

\[
\text{(8) } [\text{smoke}](\text{Ann}) = [\lambda x : x \in D . x \text{ smokes}](\text{Ann}) = 1 \text{ if Ann smokes} \\
= 0 \text{ otherwise}.
\]

Instead of (8), we'll write:

\[
\text{(8') } [\text{smoke}](\text{Ann}) = [\lambda x : x \in D . x \text{ smokes}](\text{Ann}) = 1 \text{ iff Ann smokes}.
\]

The down side of substituting this new convention for reading \( \lambda \)-terms for the previous one would be that it makes garbage of those cases which we considered at first. For instance, we could no longer write things like \( \text{"}[\lambda x : x \in \mathbb{N} . x + 1]\text{"} \) if this had to be read: \( \text{"the function which maps every } x \text{ in } \mathbb{N} \text{ to } 1, \text{ if } x + 1 \text{ [sic], and to } 0 \text{ otherwise}\text{"} \).

We will have our cake and eat it too, by stipulating that \( \lambda \)-terms may be read in either one of the two ways which we have considered, whichever happens to make sense in the case at hand.
Executing the Fregean Program

Read “[λα : φ . γ]” as either (i) or (ii), whichever makes sense.

(i) “the function which maps every α such that φ to γ”

(ii) “the function which maps every α such that φ to 1, if γ, and to 0 otherwise”

Luckily, this convention doesn’t create any ambiguity, because only one clause will apply in each given case. If γ is a sentence, that’s clause (ii), otherwise (i).

You may wonder why the same notation has come to be used in two distinct senses. Wouldn’t it have been wiser to have two different notations? This is a legitimate criticism. There does exist a precise and uniform interpretation of the λ-notation, where λ-terms can be shown to have the same meaning in the two cases, after all, despite the superficial disparity of the English paraphrases. This would require a formalization of our current informal metalanguage, however. We would map phrase structure trees into expressions of a λ-calculus, which would in turn be submitted to semantic interpretation. Here, we use λ-operators and variables informally in the metalanguage, and rely on a purely intuitive grasp of the technical locutions and notations involving them (as is the practice, by the way, in most mathematical and scientific texts). Our use of the λ-notation in the metalanguage, then, has the same status as our informal use of other technical notation, the abstraction notation for sets, for example.

The relative conciseness of the λ-notation makes it especially handy for the description of function-valued functions. Here is how we can express the lexical entry of a transitive verb.

In (10), we have a big λ-term, whose value description is a smaller λ-term. Which clause of the reading convention in (9) applies to each of these? In the smaller one (“[λy : y ∈ D . y loves x]”), the value description is evidently sentential, so we must use clause (ii) and read this as “the function which maps every y in D to 1, if y loves x, and to 0 otherwise”. And since this phrase is a noun phrase, clause (i) must apply for the bigger λ-term, and thus that one reads: “the function which maps every x in D to the function which maps every y in D to 1, if y loves x, and to 0 otherwise”.

Functions can have functions as arguments. Here, too, the λ-notation is handy. Take:

The function in (11) maps functions from $D_{<c,t>}$ into truth-values. Its arguments, then, are functions from individuals to truth-values. The function in (12) is a possible argument:
(12) \[\lambda y : y \in D_e . y \text{ stinks}\]

If we apply the function in (11) to the function in (12), we get the value 1 if there is some \(x \in D_e\) such that \([\lambda y : y \in D_e . y \text{ stinks}] (x) = 1\). Otherwise, the value is 0. That is, we have:

(13) \[\lambda f : f \in D_{\leq,t} . there\ is\ some\ x \in D_e\ such\ that\ f(x) = 1] ( [\lambda y : y \in D_e . y \text{ stinks}] ) = 1\]

iff there is some \(x \in D_e\) such that \(\lambda y : y \in D_e . y \text{ stinks}] (x) = 1\)

iff there is some \(x \in D_e\) such that \(x \text{ stinks}\).

Let us introduce a couple of abbreviatory conventions which will allow us to describe the most common types of functions we will be using in this book even more concisely. First, we will sometimes omit the outermost brackets around a \(\lambda\)-term which is not embedded in a larger formal expression. Second, we will contract the domain condition when it happens to be of the form "\(\alpha \in \beta\)". Instead of "\(\lambda \alpha : \alpha \in \beta \cdot \gamma\)", we will then write "\(\lambda \alpha \in \beta \cdot \gamma\". (This corresponds to shortening the paraphrase "every \(\alpha\) such that \(\alpha\ is\ in\ \beta\) to "every \(\alpha\ in\ \beta\".) And sometimes we will leave out the domain condition altogether, notably when it happens to be "\(x \in D\)". So the lexical entry for "love" may appear, for example, in either of the following shorter versions:

(14) \[[\text{love}] := \lambda x \in D . (\lambda y \in D . y \text{ loves } x)\]

\[[\text{love}] := \lambda x . (\lambda y . y \text{ loves } x)\]

You have to be careful when \(\lambda\)-terms are followed by argument terms. Without further conventions, 15(a) is not legitimate, since it is ambiguous between 15(b) and 15(c), which are not equivalent:

(15) (a) \(\lambda x \in D . (\lambda y \in D . y \text{ loves } x)(\text{Sue})\)

(b) \(\lambda x \in D . (\lambda y \in D . y \text{ loves } x)(\text{Sue}) = \lambda x \in D . \text{Sue loves } x\)

(c) \((\lambda x \in D . (\lambda y \in D . y \text{ loves } x))(\text{Sue}) = \lambda y \in D . y \text{ loves Sue}\)

In cases of this kind, use either the formulation in (b) or the one in (c), whichever corresponds to the intended meaning.

There is a close connection between the abstraction notation for sets and the \(\lambda\)-notation for functions. The characteristic function of the set \(\{x \in \mathbb{N} : x \neq 0\}\) is \(\lambda x \in \mathbb{N} . x \neq 0\), for example. Much of what you have learned about the abstraction notation for sets in chapter 1 can now be carried over to the \(\lambda\)-notation for functions. Set talk can be easily translated into function talk. Here are the correspondences for some of the cases we looked at earlier:
Set talk

29 ∈ \{x ∈ \mathbb{N} : x ≠ 0\} iff 29 ≠ 0
Massachusetts ∈ \{x ∈ D : California is a western state\} iff California is a western state.
(x ∈ D : California is a western state) = D if California is a western state.
(x ∈ D : California is a western state) = \emptyset if California is not a western state.

{x ∈ \mathbb{N} : x ≠ 0} = \{y ∈ \mathbb{N} : y ≠ 0\}
{x ∈ \mathbb{N} : x ∈ \{x ∈ \mathbb{N} : x ≠ 0\}} =
{x ∈ \mathbb{N} : x ≠ 0}
{x ∈ \mathbb{N} : x ∈ \{y ∈ \mathbb{N} : y ≠ 0\}} =
{x ∈ \mathbb{N} : x ≠ 0}

Function talk

[λx ∈ \mathbb{N} . x ≠ 0](29) = 1 iff 29 ≠ 0
[λx ∈ D . California is a western state](Massachusetts) = 1 iff California is a western state.
[λx ∈ D . California is a western state](x) = 1 for all x ∈ D if California is a western state.
[λx ∈ D . California is a western state](x) = 0 for all x ∈ D if California is not a western state.

[λx ∈ \mathbb{N} : x ≠ 0] = [λy ∈ \mathbb{N} : y ≠ 0]
[λx ∈ \mathbb{N} . [λx ∈ \mathbb{N} . x ≠ 0](x)] =
[λx ∈ \mathbb{N} ∈ . x ≠ 0]
[λx ∈ \mathbb{N} . [λy ∈ \mathbb{N} . y ≠ 0](x)] =
[λx ∈ \mathbb{N} . x ≠ 0]

If you are still unclear about some of the statements in the left column, go back to chapter 1 and consult the questions and answers about the abstraction notation for sets. Once you understand the set talk in the left column, the transition to the function talk in the right column should be straightforward.

---

Exercise 1

Describe the following functions in words:

(a) \( \lambda x ∈ \mathbb{N} . x > 3 \) and \( x < 7 \)
(b) \( \lambda x : x \) is a person . \( x \)'s father
(c) \( \lambda X ∈ \text{Pow}(D) . \{y ∈ D : y \not\in X\}^{20} \)
(d) \( \lambda X : X \subseteq D . [\lambda y ∈ D . y \not\in X] \)

---

Exercise 2

In this exercise, simple functions are described in a rather complicated way. Simplify the descriptions as much as possible.
(a) \[\{\lambda x \in D . [\lambda y \in D . [\lambda z \in D . z \text{ introduced } x \text{ to } y]]\}(Ann)(Sue)\]
(b) \[\{\lambda x \in D . [\lambda y \in D . [\lambda z \in D . z \text{ introduced } x \text{ to } y]]\}(Ann)(Sue)\]
(c) \[\{\lambda x \in D . [\lambda y \in D . [\lambda z \in D . z \text{ introduced } x \text{ to } y]]\}(Ann)(Sue)\]
(d) \[\{\lambda x \in D . [\lambda y \in D . [\lambda z \in D . z \text{ introduced } x \text{ to } y]]\}(Ann)(Sue)\]
(e) \[\{\lambda f \in D_{<e,l>} . [\lambda x \in D_e . f(x) = 1 \text{ and } x \text{ is gray}]\}(\{\lambda y \in D_e . y \text{ is a cat}\})\]
(f) \[\{\lambda f \in D_{<e,<e,l>} . [\lambda x \in D_e . f(x)(Ann) = 1]\}(\{\lambda y \in D_e . [\lambda z \in D_e . z \text{ saw } y]\})\]
(g) \[\{\lambda x \in |N . [\lambda y \in |N . y > 3 \text{ and } y < 7](x)\}\]
(h) \[\{\lambda z \in |N . [\lambda y \in |N . [\lambda x \in |N . x > 3 \text{ and } x < 7](y)](z)\}\]

Exercise 3

Suppose “and” and “or” have Schönfinkeled denotations, that is \[[\text{and}]\] and \[[\text{or}]\] are both members of \(D_{<1,<1,1>}\). They are functions that map truth-values into functions from truth-values to truth-values. Specify the two functions using the \(\lambda\)-notation.

Exercise 4

Replace the “?” in each of the following statements (you may want to review definition (5) of section 2.3 before tackling this exercise):

(a) \[\{\lambda f \in D_{<e,l>} . [\lambda x \in D_e . f(x) = 1 \text{ and } x \text{ is gray}]\} \in D?\]
(b) \[\{\lambda f \in D_{<e,<e,l>} . [\lambda x \in D_e . f(x)(Ann) = 1]\} \in D?\]
(c) \[\{\lambda y \in D_e . [\lambda f \in D_{<e,l>} . [\lambda x \in D_e . f(x) = 1 \text{ and } x \text{ is in } y]]\} \in D?\]
(d) \[\{\lambda f \in D_{<e,l>} . \text{ there is some } x \in D_e \text{ such that } f(x) = 1\} \in D?\]
(e) \[\{\lambda f \in D_{<e,l>} . \text{ Mary} \} \in D?\]
(f) \[\{\lambda f \in D_{<e,l>} . [\lambda g \in D_{<e,l>} . \text{ there is no } x \in D_e \text{ such that } f(x) = 1 \text{ and } g(x) = 1]\} \in D?\]

Notes

1 Here and below, when we speak of the denotation of a node in a tree, we really mean the denotation of the subtree dominated by that node.
2 M. Black and P. Geach, Translations from the Philosophical Writings of Gottlob Frege (Oxford, Basil Blackwell, 1960), use “reference” to translate Frege’s “Bedeutung”.


4 The use of Wittgenstein's term "show" in this connection is due to Dummett, ibid.

5 This is another mathematical-background section, which the mathematically sophisticated need only skim.

6 The notation "\(\{x \in D : x \text{ sleeps}\}\)" is a standard abbreviation for "\(\{x : x \in D \text{ and } x \text{ sleeps}\}\)". (Recall that when we first introduced the set-abstraction notation, we allowed only a variable to the left of the colon.)

7 This conclusion, even though we are motivating it by using his general proposal about semantic composition, Frege himself would not have endorsed. As discussed by Dummett (*Frege*, pp. 40ff.), he did not allow for function-valued functions.

8 Notice that the implicit bracketing in "\(f(x)(y)\)" is "\([f(x)](y)\)" not "\(f[(x)(y)]\)". The latter wouldn't make any sense. (What could we possibly mean by "\((x)(y)\)"?) So it is not necessary to make the correct parse explicit in the notation.

9 R. Montague, *Formal Philosophy* (New Haven, Yale University Press, 1974); "e" is for "entity", "t" for "truth-value".

10 \(D \times D\) is the *Cartesian product* of \(D\) with \(D\), which is defined as the set of ordered pairs of elements of \(D\).

11 A 2-place function on a domain \(A\) is a function with domain \(A \times A\), which is defined as \(\{<x, y> : x \in A \text{ and } y \in A\}\).


13 This procedure is also called "Currying" after the logician H. B. Curry, who built on Schönfinkel's work. Molly Diesing informs us that by Stephen Jay Gould's "Brontosaurus principle", one could argue that we should use "Currying": generality of use takes priority over temporal precedence. We are not sure how general the use of "Currying" is at this time, however, hence we don't know whether the Brontosaurus Principle applies to this case. We'll stick to temporal priority, then. W. Kneale and M. Kneale, *The Development of Logic* (Oxford, Clarendon Press, 1962), pp. 522f., credit Schönfinkel for Schönfinkelizeation.

14 Most versions of the \(\lambda\)-notation in the literature look a little different. What we are calling the "domain condition" is typically absent, and the intended domain is indicated instead by using argument variables that are assigned to fixed semantic types. The value description is frequently enclosed in brackets rather than introduced by a period. The terms "argument variable", "domain condition", and "value description" are also our own invention.

15 The attentive reader may have noticed that another piece of information seems to get lost in the reformulation from (1) to (2): viz. information about which set \(F_{+}\) is into. Nothing in (2) corresponds to the part "\(\rightarrow \text{IN}\)" in (1). Is this a problem? No. If you go back to our initial definition of "function" in section 1.3, you can see that the information supplied in (2) is entirely sufficient to define a unique function. It already follows from (2) that all values of \(F_{+}\) are in \(\text{IN}\). In other words, the format employed in (1) is actually redundant in this respect.

16 The best we can do, if we insist on using the \(\lambda\)-notation to define \(G\), is to describe \(G\) as the union of two separate functions with smaller domains:
G := \lambda x : x \in \mathbb{N} \& x \text{ is even} . \ 2 \cup \lambda x : x \in \mathbb{N} \& x \text{ is odd} .

But this is not much of an abbreviation of (5). Remember, we were planning to save ink.

Incidentally, the indication that \( f \) is into \{0, 1\} is not redundant in (6). Without this information, we could not conclude that \( f(x) = 0 \) when \( x \) doesn't smoke. From the condition "\( f(x) = 1 \) iff \( x \) smokes" by itself, we can infer only that \( f(x) \neq 1 \) in this case. If we want to render "\( \to \{0, 1\} \)" as redundant here as "\( \to \mathbb{N} \)" was in (1) and (5), we have to replace "\( f(x) = 1 \) iff \( x \) smokes" by something like the formulation in the text below: "\( f(x) = 1 \) if \( x \) smokes and 0 otherwise'.

Even with this disjunctive convention, by the way, the \( \lambda \)-notation remains unsuitable for abbreviating definitions like the one for \( G \) in (5). But we can live with this, because functions like \( G \) are not common in semantic applications.


"\( \text{Pow}(D) \)" reads "the power set of \( D \)". The power set of a set is the set of all its subsets. (More formally: for any set \( X \), \( \text{Pow}(X) := \{Y : Y \subseteq X\} \).)
In defining the semantic component for the small fragment of English we considered in the previous chapter, we followed the traditional method of writing a special semantic rule for each syntactic configuration that we encountered in our trees. There was a rule for Ss dominating an NP and a VP, a rule for VPs dominating a V and an NP, a rule for VPs dominating just a V, and so on. But if Frege's Conjecture is right, and all non-trivial semantic composition is functional application, then we shouldn't need all these construction-specific interpretation rules. It should be sufficient to specify the denotations of the lexical items, and the rest should follow automatically.

We begin the present chapter by implementing this attractive suggestion. Then we proceed to some reflections about the place of the semantic component within the grammar, especially its relation to the syntactic component. Our Fregean approach gives us an interesting perspective on the relation between a verb and its arguments, which has been discussed from more syntactic points of view throughout the history of generative grammar.

3.1 Type-driven interpretation

The term "type-driven interpretation" was coined by Ewan Klein and Ivan Sag, whose paper "Type-Driven Translation" criticized the construction-specific interpretation method of classical Montague Grammar and proposed essentially the same revision that we present in this section. We continue to assume that the input for the semantic component is a set of phrase structure trees. But we no longer allow semantic rules for specific types of subtrees like the ones we wrote in chapter 2. Instead, we posit three very general principles:

(1) *Terminal Nodes* (TN)

If $\alpha$ is a terminal node, $[\alpha]$ is specified in the lexicon.
(2) **Non-Branching Nodes (NN)**
   If $\alpha$ is a non-branching node, and $\beta$ is its daughter node, then $[\alpha] = [\beta]$.

(3) **Functional Application (FA)**
   If $\alpha$ is a branching node, $[\beta, \gamma]$ is the set of $\alpha$'s daughters, and $[\beta]$ is a function whose domain contains $[\gamma]$, then $[\alpha] = [\beta]([\gamma])$.

Notice that (3) does not specify the linear order of $\beta$ and $\gamma$. Nevertheless, it applies in a unique way to each given binary branching tree. If $\alpha$ is of the form $[s \text{ NP VP}]$, we have to apply it in such a way that the right-hand daughter corresponds to $\beta$ and the left-hand daughter to $\gamma$. How do we know? Because this is the only way to satisfy the condition that $[\beta]$ must be a function whose domain contains $[\gamma]$. If $\alpha$ is of the form $[v \text{ V NP}]$, it's the other way round: $\beta$ must be the left node, and $\gamma$ the right one. Here you can see, by the way, what is behind the name “type-driven interpretation”: it's the semantic types of the daughter nodes that determine the procedure for calculating the meaning of the mother node. The semantic interpretation component, then, can ignore certain features that syntactic phrase structure trees are usually assumed to have. All it has to see are the lexical items and the hierarchical structure in which they are arranged. Syntactic category labels and linear order are irrelevant.

During the course of the book, we will add one or two additional principles to the above list, but we will strive to keep it as parsimonious as possible. When we look at a new construction for which we don't yet have a semantic analysis, we always try first to accommodate it by adding only to the lexicon.

Our current set of lexical entries is just the beginning of a long list that we will extend as we go:

(4) **Sample of lexical entries:**
   
   (i) $[\text{Ann}] = \text{Ann}$
   
   (ii) $[\text{smokes}] = \lambda x \in D_e \cdot x \text{ smokes}$
   
   (iii) $[\text{loves}] = \lambda x \in D_e \cdot [\lambda y \in D_e \cdot y \text{ loves } x]$
   
   (iv) $[\text{and}] = \begin{bmatrix}
   1 & \rightarrow & 1 \\
   0 & \rightarrow & 0 \\
   1 & \rightarrow & 0 \\
   0 & \rightarrow & 0
   \end{bmatrix}$

   or, using the $\lambda$-notation:

   $[\text{and}] = \lambda p \in D_t \cdot [\lambda q \in D_t \cdot p = q = 1]^3$

   etc.

Given the lexicon, the three interpretive principles TN, NN, and FA should suffice to derive all the predictions that were made by our old semantic component.
from chapter 2. For the most part, it is transparent how this works out. The semantic rules (S2)–(S5) in chapter 2 are evidently special cases of NN, and (S1) and (S6) are special cases of FA. What is not covered by the new theory is the interpretation of ternary trees, like those we initially assumed for sentences with “and” and “or”. The semantic rules from the exercise on connectives in section 2.1 cannot be seen as special cases of any current principle. (Our formulation of FA above stipulates that $\alpha$ has no more than two daughters.) But rather than worry about this limitation, we will henceforth assume (mainly for pedagogical reasons) that phrase structures in natural language are at most binary branching. If necessary, it would not be too difficult to adjust our system of composition rules so that it could interpret a more realistic range of input structures.

3.2 The structure of the input to semantic interpretation

We have been assuming that the entities which are assigned denotations by the semantic component are phrase structure trees, and that these are somehow generated by the syntactic component of the grammar. We have not committed ourselves to any concrete assumptions about the kind of syntax that does this. Indeed, a variety of views on this matter are compatible with our approach to semantics and have been entertained. According to the so-called “Standard Theory” of early generative grammar, the input to semantic interpretation consisted of the Deep Structures generated by the base component of the syntax, essentially a context-free grammar. In Generative Semantics, Deep Structures were representations that resembled formulas of the Predicate Calculus, and could involve decomposition of predicates. From a more modern Chomskyan perspective, the inputs to semantic interpretation are Logical Forms, which are the output of transformational derivations. Many other syntactic theories can be, and have been, combined with the same kind of semantic theory – for instance, categorial grammars and monostratal phrase structure grammars. The only requirement for the syntax is that it provides us with phrase structure trees.

When we say “(phrase structure) trees”, what exactly do we mean? According to standard definitions, a phrase structure tree consists of a (finite) set of labeled nodes which are ordered by a dominance relation and a linear precedence relation. Many structural concepts can be defined in terms of dominance and precedence, among them the ones that happen to be referred to by our interpretation principles TN, NN, and FA: namely, “(non-) terminal node”, “(non-) branching”, and “daughter”. This being so, phrase structure trees in the standard sense are suitable structures for our semantic rules to apply to them.

We mentioned already that certain somewhat more impoverished structures would be suitable as well. Nothing in our semantic component – in particular, none of our principles TN, NN, and FA – makes any direct or indirect mention of non-terminal node labels or linear precedence. "(Non-) terminal node", "(non-) branching", and "daughter" are all definable in terms of the dominance relation alone, and these are the only syntactic notions that have been mentioned. We therefore conclude that our semantics can also be combined with a syntax that provides unlinearized structures without syntactic category labels instead of standard phrase structure trees. If the syntactic theory you are working with is of that kind, you may tacitly ignore the syntactic category labeling and linearization information encoded by the tree diagrams which we will be drawing in this book. It won’t make any difference.

Even though our semantics does not rely on syntactic category labels, we will normally employ at least some commonly used labels when we present phrase structure trees. Using syntactic category labels makes it easier for us to talk about particular nodes or types of nodes in a tree. We want to emphasize, however, that our choice of labels does not have any theoretical significance; the semantic component doesn’t have to see them. Here is an overview of the lexical category labels that we will be using most often:

<table>
<thead>
<tr>
<th>syntactic category</th>
<th>label</th>
</tr>
</thead>
<tbody>
<tr>
<td>verb</td>
<td>V</td>
</tr>
<tr>
<td>noun</td>
<td>N</td>
</tr>
<tr>
<td>adjective</td>
<td>A</td>
</tr>
<tr>
<td>preposition</td>
<td>P</td>
</tr>
<tr>
<td>determiner</td>
<td>D</td>
</tr>
<tr>
<td>inflectional elements</td>
<td>I</td>
</tr>
<tr>
<td>(&quot;do&quot;, tense, modal auxiliaries, etc.)</td>
<td></td>
</tr>
<tr>
<td>complementizer</td>
<td>C</td>
</tr>
</tbody>
</table>

Labels for phrasal categories are coined in the usual way. A maximal verbal projection is a VP (verb phrase), for example. An intermediate verbal projection is a V (V-bar). We will be using "S" or "IP" (inflection phrase) as labels for sentences. Other labels may be introduced as we go along.

The kind of semantic theory we are developing here is compatible with a wide range of approaches to syntax. It may be worth pointing out, however, that it is nevertheless fundamentally incompatible not only with many conceivable proposals regarding the structure of the input to semantic interpretation, but even with some that have actually been made. For instance, Jackendoff and other representatives of the Extended Standard Theory argued that meaning depended on both Deep Structure and Surface Structure. Roughly, Deep Structure was to
determine predicate–argument relations, and Surface Structure scope and binding relations. In our terms, this would mean that the domain of the interpretation function \[ D \] should consist of (or include) something like pairs of phrase structure trees. What this would amount to in concrete detail is far from clear, and our current conception of semantic composition would have to be substantially altered and/or enriched to make it work.\(^\text{10}\)

### 3.3 Well-formedness and interpretability

Montague's view of the role of syntax was that the syntactic component served to generate *exactly* the trees to which the semantic component assigned denotations. In such a framework, there is no such thing as a tree generated by the syntax but not in the domain of the interpretation function \[ D \], or a tree in the domain of \[ D \] but not generated by the syntax. Syntactic well-formedness and semantic interpretability coincide completely. This was just the situation that philosophical logicians were used to from their work on formal languages.

But linguists have questioned this set-up for a long time. Even before the word "modularity" became common, it was widely agreed that the grammars acquired by human beings do not just divide the set of all possible expressions into the well-formed, meaningful ones versus all the rest. Sentences or near-sentences may be judged deviant for a variety of independent reasons. It is easy to find examples that speakers reject as ungrammatical despite perfectly clear and definite intuitions about what they mean. Considering how general and universal our principles of interpretation are, of course, it is only to be expected that they will apply not only to those structures that happen to represent grammatical English sentences, but to many others besides.

The existence of the reverse case – that is, of uninterpretable but otherwise completely well-formed examples – is harder to establish without already presupposing a certain amount of theory, but there is certainly no reason to rule it out a priori either. The following examples might be cases in point.

(1) *Ann laughed Jan.

(2) *It is not the case that greeted Ann.

Suppose the syntax of English provides derivations for these which are not ruled out by any syntactic (or phonological) principle and which assign them the following phrase structures.
What happens if we try to calculate denotations for these trees by means of our current composition rules and lexicon?

In (1') we obtain a truth-value for the constituent “laughed Jan” (by applying the function \([\text{laughed}]\), of type \(\langle e,t\rangle\), to Jan, of type \(e\)). But at the next higher node, we can’t apply any of our rules. FA demands that either \([\text{Ann}]\) or \([\text{laughed Jan}]\) be a function. Since one is a person and the other a truth-value, it cannot apply. In (2'), we obtain a function of type \(\langle e,t\rangle\) for the node dominating \(\text{greeted Ann}\) (by applying \([\text{greeted}]\), of type \(\langle e,\langle e,t\rangle\rangle\), to Ann, of type \(e\)). There again we are stuck at the next node up: \([\text{it is not the case that}]\) is a function of type \(\langle t,t\rangle\) and \([\text{greeted Ann}]\) a function of type \(\langle e,t\rangle\). Neither has the other in its domain, so FA cannot apply.

Thus, in both cases we are looking at trees that don’t receive any denotation by our semantic component. That is, they are not in the domain of the \(\left[\right]\) function as defined by our current lexicon and composition rules. It is reasonable to hypothesize that this – and this alone – is what accounts for the deviance judgments that we represented by the asterisks in (1) and (2). We have not shown, of course, that there isn’t something else wrong with these structures in addition to their failure to receive denotations. But let’s suppose for the sake of the argument here that there isn’t.

Structures which, like (1') and (2'), fail to receive denotations will be called uninterpretable. We take it that uninterpretability is one among other sources of ungrammaticality. Uninterpretable structures are those filtered out by the semantic component of the grammar. Here is a more precise formulation of our set of semantic rules, which explicitly takes into account the possible existence of uninterpretable (sub)trees. We define the interpretation function \(\left[\right]\) as the smallest function that fulfills the following conditions:

\[ (\text{3}) \quad \text{Terminal Nodes (TN)} \]

If \(\alpha\) is a terminal node, then \(\alpha\) is in the domain of \(\left[\right]\) if \([\alpha]\) is specified in the lexicon.
(4) **Non-Branching Nodes (NN)**

If $\alpha$ is a non-branching node, and $\beta$ is its daughter node, then $\alpha$ is in the domain of $[\ ]$ if $\beta$ is. In this case, $[\alpha] = [\beta]$.

(5) **Functional Application (FA)**

If $\alpha$ is a branching node and $\{\beta, \gamma\}$ is the set of $\alpha$'s daughters, then $\alpha$ is in the domain of $[ ]$ if both $\beta$ and $\gamma$ are and $[\beta]$ is a function whose domain contains $[\gamma]$. In this case, $[\alpha] = [\beta](\gamma)$.

Notice that each of (3)–(5) gives sufficient (but not necessary) conditions for $\alpha$ being in the domain of the interpretation function $[\ ]$. But by defining $[\ ]$ as the *smallest* function that meets all these conditions, we say, in effect, that a tree is not in the domain of $[\ ]$ unless one of (3)–(5) implies that it is. We could have achieved the same result, of course, by using "iff" rather than "if" in (3)–(5).

Finally, let us make explicit the filtering function of the semantic component:

(6) **Principle of Interpretability**

All nodes in a phrase structure tree must be in the domain of the interpretation function $[\ ]$.

In sum, we are adopting a view of the grammar as a whole on which syntax and semantics are independent modules. Each imposes its own constraints on the grammatical structures of the language, and we expect there to be structures that are interpretable though syntactically illegitimate, as well as structures that are syntactically correct but uninterpretable.

### 3.4 The $\Theta$-Criterion

In the syntactic literature, the ungrammaticality of examples like (1) and (2) has sometimes been attributed to the so-called $\Theta$-Criterion (Theta-Criterion).

(1) *Ann laughed Jan.

(2) *It is not the case that greeted Ann.

Chomsky gives the following formulation.\(^{11}\)

(3) **$\Theta$-Criterion**

Each argument bears one and only one $\Theta$-role, and each $\Theta$-role is assigned to one and only one argument.
What are Θ-roles? The term “Θ-role” (“theta-role”) is an abbreviation for the more commonly used term “thematic role”. If an event of greeting takes place somewhere, then there is a greeter and someone who is greeted. Greeter and greeter are two specific Θ-roles. If a kicking takes place, then there is a kicker and someone who is kicked. Kicker and kickee are again two specific Θ-roles. A verb like “greet” requires one argument that is interpreted as the greeter, and another argument that is interpreted as the one that is greeted. We may say, then, that it assigns these two specific Θ-roles to its arguments. A verb like “kick” requires one argument that is interpreted as the kicker, and another argument that is interpreted as the one that is kicked. Again, we may say that it assigns those two specific Θ-roles. Kicking and greeting are both actions, kickers and greeters are both agents, and whatever is greeted or kicked is a patient or theme according to a common terminology. Usually, when you hear about Θ-roles, it’s general Θ-roles like “agent”, “theme”, “patient”, or “experiencer” that are being talked about. General Θ-roles are potentially interesting, since there are generalizations linking them to syntactic positions in a systematic way. As far as the Θ-criterion is concerned, however, only specific Θ-roles are relevant. Now, what does the Θ-criterion do?

One part of the Θ-Criterion says that whenever a lexical element requires an argument with a certain Θ-role, then there must be such an argument somewhere in the syntactic representation. And there can be only one. There couldn’t be two NPs, for example, that are both interpreted as the greeter argument for the verb “greet”. The other part of the Θ-Criterion says that whenever there is a “candidate” for an argument, like an NP or whatever you, then this element must be in fact an argument of a lexical element that assigns it a Θ-role. That it can’t be assigned more than one Θ-role means that it can’t fill more than one argument position at once.

The Θ-Criterion is meant to exclude the following kinds of cases:

(4) *Ann laughed Jan. (= (1) above)

(5) *Greeted Ann.

(4) is ruled out by it since one of the NPs cannot be assigned a Θ-role. “Laugh” has only one Θ-role to assign, and since this one Θ-role cannot be assigned to two NPs, (4) cannot possibly mean that Ann and Jan both laughed. (5) is ruled out by the Θ-Criterion since “greeted” assigns two Θ-roles, but there is only one NP present. Since no NP can receive two Θ-roles, (5) cannot mean that Ann greeted herself.

We saw above how similar examples are treated by our current semantic theory: (4) came out as simply uninterpretable; that is, it receives no denotation.
In other words, it is predicted to have no truth-conditions at all, neither those of "Ann and Jan both laughed" nor any others. As far as (4) is concerned, then, the predictions of our semantics coincide fully with those of the $\Theta$-Criterion.

May we conclude that the $\Theta$-Criterion is – apart from differences in terminology – just a corollary of our current theory? In other words, does every $\Theta$-Criterion violation reduce to a case of uninterpretability in the sense we specified above? Or is there merely a certain overlap in predictions? Let's take a closer look at this question.

What about (5)? The Principle of Interpretability succeeds (just like the $\Theta$-Criterion) in ruling out examples like (2) above, in which (5) appears as an embedded constituent. But (5) all by itself isn't actually uninterpretable according to our semantics. It gets, in effect, the denotation of the VP "greeted Ann" (which is the function $\lambda x \in D_e \cdot x$ greeted Ann). We therefore predict correctly that "greeted Ann" cannot be used to make a statement – that is, is not capable of truth or falsity.

Summarizing the discussion so far, we still haven't seen any predictions of the $\Theta$-Criterion that were not covered by the Principle of Interpretability in our current theory. But it would be premature to stop looking. In fact, the following abstract consideration shows that the $\Theta$-Criterion is a genuinely stronger constraint than Interpretability.

Suppose we have a predicate $\alpha$ with one $\Theta$-role to assign. In our terms, suppose that $[\alpha]$ is of type $<e,t>$. According to the $\Theta$-Criterion, $\alpha$ must appear in the vicinity of something which receives its $\Theta$-role. That means $\alpha$ has to have a sister node with a meaning of type e. According to our Interpretability principle, on the other hand, a sister node of type e is not strictly required. It would provide one suitable environment for $\alpha$, but not the only kind. Imagine instead that $\alpha$ has a sister node whose meaning is a function with domain $D_{<e,t>}$ (For instance, it might be of type $<<e,t>,e>$. In that case, the next higher node could be interpreted by applying this sister's meaning to $[\alpha]$. So we could have an interpretable structure which does not contain an argument for $\alpha$! $\alpha$ would not be assigning its $\Theta$-role to any phrase, in violation of the $\Theta$-Criterion. Yet Interpretability would be fulfilled, in virtue of $\alpha$ being a suitable argument for its sister node.

The question now is: Can we find concrete instances of this sort of situation? And if yes, are they in fact grammatical (as predicted by Interpretability) or ungrammatical (as predicted by the stronger $\Theta$-Criterion)?

In the following chapter, we will propose that common nouns like "barn" are 1-place predicates (type $<e,t>$). In other words, they have a $\Theta$-role to assign, and thus the $\Theta$-Criterion requires the presence of an argument. In certain examples (predicative uses), this is unproblematic:

(6) This is a barn.
The required argument here is the subject NP “this”. (6) is true if and only if the object referred to by “this” has the property of being a barn. But consider the following sentence:

(7) The barn burned down.

(7) contains no phrase that receives the Θ-role of “barn”. It thus seems to violate the Θ-Criterion. Yet it is perfectly fine, and we will see below how it can be interpreted by assigning “the” a meaning of type \(\langle \langle t,t,\rangle,\rangle\), suitable to take \([\text{barn}]\) as an argument. So this is the sort of case we have been looking for. The Interpretability Principle and the Θ-Criterion make different predictions here, and if the analysis we will give for (7) is on the right track, then the empirical facts favor the former.\(^\text{13}\)

Another case which prima facie points in the same direction arises with coordinated predicates. So far, all our examples involving “and” and “or” had these connectives combining with sentences, as in propositional logic. But English apparently also allows these connectives to conjoin subsentential constituents, such as two VPs, as in (8).\(^\text{14}\)

(8)\[ S \]
\[ NP \quad VP \]
\[ Ann \quad VP \quad and \quad VP \]
\[ \text{sings} \quad \text{dances} \]

To interpret this, we need a new lexical entry for the homonym of “and” that we encounter here. (The familiar meaning of type \(\langle t, t, t, \rangle\) evidently doesn’t work here.)

(9) \[ [\text{and}] = \lambda f \in D_{\langle t, t, \rangle} . [\lambda g \in D_{\langle t, t, \rangle} . [\lambda x \in D_{\langle t, t, \rangle} . f(x) = g(x) = 1]] \]

You can verify that, given (9), the tree in (8) is interpretable by our general principles and receives intuitively correct truth-conditions. What interests us here is that this is another interpretable structure which seems to violate the Θ-Criterion, in that there are not enough arguments to go around for all the Θ-roles that need to be assigned in (8). “Sing” and “dance” each have a Θ-role to assign, but only one potential argument (the NP “Ann”) is present. Once more, we tentatively conclude that the weaker requirements imposed by our Interpretability Principle make the better empirical predictions.
These two arguments against the $\Theta$-Criterion are not beyond question, of course. They are only as good as the syntactic and semantic analyses we have sketched. It might very well turn out upon closer inspection that there is more than meets the eye to the structures of examples (7) and (8). Specifically, we might find evidence for non-overt constituents which provide just the superficially missing arguments that are demanded by the $\Theta$-Criterion. In that event, these examples might not only cease to be counterexamples, but might ultimately turn out even to support the $\Theta$-Criterion. We will set this possibility aside for now, but in principle it remains open.

3.5 Argument structure and linking

Some syntactic theories posit a syntactic representation of a verb's argument structure that is distinct from the representation of the verb's denotation. Argument structure representations are meant to encode "the syntactically relevant argument-taking properties of a verb" (and any lexical item that takes arguments). Argument structure representations play a role in theories of linking—that is, theories about how a verb's arguments are linked to syntactic positions in a tree. In this section, we will look at some proposals that have actually been made for argument structure representations and linking. We will see that some of the information that has been attributed to argument structure representations and linking principles turns out to be redundant, given the semantic theory we have been developing.

Minimally, the argument structure representations found in the syntactic literature list the arguments a verb takes. This information has been thought to be relevant to the syntax, because of the deviance of sentences like (1) or (2):

(1) *Ann laughed Jan.

(2) *Greeted Ann.

In the previous section, we discussed the possibility that the deviance of (1) and (2) might actually be accounted for in the semantics. (1) fails Interpretability, and (2) receives a VP-denotation, hence cannot be used to make a statement capable of truth or falsity. If this view of the syntax–semantics interface should turn out to be correct, there is no need for separate syntactic representations of argument structure, unless they provide more information than a mere list of the verb's arguments.
In her book *Argument Structure*, Jane Grimshaw explores the hypothesis that argument structure representations also reflect prominence relations among arguments. The verb "introduce" would be given the following argument structure representation, for example:

\[ \text{introduce (agent (goal (theme)))} \]

(3) says that "introduce" has three arguments that are hierarchically ordered: the agent argument is the highest, the goal argument comes next, and the theme argument is at the bottom. Grimshaw emphasizes that the general thematic role labels "agent", "goal", or "theme" have no theoretical status; they merely serve to identify the verb's arguments. She could have used labels for specific thematic roles instead. A prominence relation among arguments is part and parcel of our Fregean verb denotations. The Fregean denotation of "introduce" is

\[ \lambda x \in D_e . [\lambda y \in D_e . \lambda z \in D_e . z \text{ introduces } x \text{ to } y] \]

Like (3), (4) encodes the information that the agent argument of "introduce" is most prominent, the goal argument is next, and the theme argument fills the lowest position. If we apply the function in (4) to an individual, Sue, for example, we get (5).

\[ \lambda y \in D_e . [\lambda z \in D_e . z \text{ introduces Sue to } y] \]

The first argument to be processed, then, is given the role of the person who is introduced (the theme). Next, if we apply the function in (5) to Ann, the result is (6): Ann is assigned the role of the person who Sue is introduced to (the goal):

\[ \lambda z \in D_e . z \text{ introduces Sue to Ann} \]

If we finally apply the function in (6) to, say, Pat, we end up with a truth-value: "True" if Pat introduces Sue to Ann, and "False" otherwise. That is, the last argument to be processed is understood as the agent of the introduction event. We can conclude, then, that we do not need separate argument structure representations to display prominence relations among arguments. This information is already provided by the representation of the verb's denotation.

Our system of type-driven interpretation principles implies a rather strong claim about the linking of a verb's arguments to syntactic positions. The lexically determined prominence relations must be preserved in the syntax. This means that there couldn't be a natural language that has structures like (7) as well as structures like (8), with the truth-conditions (7') and (8') respectively.
Since the V-denotation must combine with the denotation of its direct object by Functional Application, the lexical meaning of “loves” determines that “Jan” is interpreted as the one who is loved in both (7) and (8). Consequently, (7) and (8) cannot have different truth-conditions. If we allowed construction-specific interpretation rules, we could in principle have rules like (9) and (10):

(9) If $\alpha$ has the form $\lambda x \in D_e \quad [\beta](\{\gamma\})$, then $[\alpha] = [\beta](\{\beta\}(\{\gamma\}))$.

(10) If $\alpha$ has the form $\lambda x \in D_e \quad [\beta](\{\gamma\})$, then $[\alpha] = [\beta](\{\beta\}(\{\gamma\}))$.

Keeping our earlier rules for the interpretation of lexical, non-branching, and S-nodes, both (7) and (8) would now be assigned the intended truth-conditions.
Active structure (7)
\[[\text{VP}_{\text{active}}] = [\lambda y \in D_e \cdot [\lambda z \in D_e \cdot z \text{ loves } y]](\text{Jan}) = [\lambda z \in D_e \cdot z \text{ loves } \text{Jan}].
\]
\[[7] = [\lambda z \in D_e \cdot z \text{ loves } \text{Jan}](\text{Ann}) = 1 \text{ iff Ann loves Jan}.
\]

Passive structure (8)
\[[\text{VP}_{\text{passive}}] = [\lambda x \in D_e \cdot [\lambda y \in D_e \cdot [\lambda z \in D_e \cdot z \text{ loves } y]](x)(\text{Jan})]
\]
\[[8] = [\lambda x \in D_e \cdot \text{Jan loves } x](\text{Ann}) = 1 \text{ iff Jan loves Ann}.
\]

Many syntactic theories consider it necessary to stipulate principles that prevent a language from having both structures like (7) and structures like (8). Take Mark Baker's Uniformity of Theta Assignment Hypothesis (UTAH):\(^19\)

\[(11) \text{ The Uniformity of Theta Assignment Hypothesis (UTAH)}
\]
Identical thematic relationships between items are represented by identical structural relationships.

The UTAH can be given a weaker or stronger interpretation. On the weaker interpretation, the UTAH says that NPs that bear the same specific thematic role bear the same syntactic relationship to their verb. The stronger interpretation requires that all NPs that bear the same general thematic role (given some inventory of general thematic roles) bear the same syntactic relationship to their verb. The weaker version of the UTAH requires that all arguments that bear the lover role must be syntactically realized in a uniform way, and the same is true of all arguments that bear the role of the one who is loved. The stronger version of the UTAH might require, for example, that all arguments that bear the experiencer (of an emotion) role appear in the same syntactic position, and likewise all arguments bearing the theme, patient, or agent role. Both versions exclude the hypothetical situation described above. If (7') states the correct truth-conditions for (7), then the truth-conditions (8') are ruled out for (8), and if (8') states the correct truth-conditions for (8), then the truth-conditions (7') are ruled out for (7).

We have just seen that some of the work done by principles like the UTAH is automatically taken care of by our semantic component. In fact, the weak version of the UTAH is at least close to becoming superfluous. What about the more interesting strong version? What about general thematic roles in our framework? It is fairly obvious that the prominence relations among a verb's arguments do not have to be learned separately for each verb. There are generalizations involving general thematic roles. Agent arguments are generally higher than theme or patient arguments, for example. The exact nature of such generalizations is still a matter of debate. The most worked-out proposal is that of Dowty,\(^20\)
whose argument selection principles are stated on the basis of thematic proto-roles. Dowty assumes that thematic roles like agent or patient are cluster concepts like the prototypes of Rosch and Mervis. According to him, the argument with the greatest number of proto-agent properties is selected as the lexically most prominent argument, for example. With ditransitive verbs, the argument with the greatest number of proto-patient properties would be the lowest argument. The middle argument would have fewer proto-patient properties than the lowest argument and fewer proto-agent properties than the highest argument. Be this as it may, whatever the correct generalizations about lexical prominence relations are, our semantic interpretation system automatically imposes them on the hierarchical line-up of arguments in the syntax. The syntactic component, then, does not have to worry about thematic roles, be they specific or general. What appear to be generalizations about the syntactic realization of arguments might in fact be rooted in uniformities of prominence relations across lexical items.

Does this mean that we can dispense with syntactic linking principles altogether? Not quite yet. There is an important syntactically relevant distinction between a verb's arguments that does not yet fall out from our semantics. Syntactic work that was most vigorously pursued within Relational Grammar has established that not all arguments that are lexically most prominent show the same syntactic behavior.

(12) Unaccusative verb: \([\text{die}] = \lambda x \in D_e . x \text{ dies}\]
Unergative verb: \([\text{work}] = \lambda x \in D_e . x \text{ works}\]

(13) Transitive agentive verb: \([\text{greet}] = \lambda x \in D_e . [\lambda y \in D_e . y \text{ greets } x]\]
Transitive experiencer verb: \([\text{worry}] = \lambda x \in D_e . [\lambda y \in D_e . y \text{ worries } x]\]

As far as lexical prominence relations are concerned, there is no difference between "die" and "work", and the same is true for "greet" and "worry". Yet syntactically, unaccusative verbs behave differently from unergative verbs, and agentive verbs do not pattern with (object) experiencer verbs in important respects. Using the terminology of Williams, the most prominent argument of unergative and agentive transitive verbs is an external argument. The most prominent argument of unaccusative and (object) experiencer verbs is an internal argument. According to Williams, the external argument is located external to the maximal projection of the verb, whereas internal arguments appear within the maximal projection of the verb (at some level of representation). In one way or other, the difference between external and internal arguments has been held responsible for the fact that subjects of unaccusative and (object) experiencer verbs show certain properties of objects, unlike the subjects of unergative and agentive transitive verbs.
The syntactic impact of the distinction between external and internal arguments cannot be directly derived from our semantics as is. To be sure, there are plenty of proposals which try to correlate the distinction with some semantic property. But that's not enough. This all by itself will not spare us the stipulation of a special linking principle. Our semantic composition rules only impose hierarchical relationships on the syntactic representation of arguments. There is no way for them to influence whether a most prominent argument is realized within or outside a VP, for example. One possible conclusion, then, is that the distinction between external and internal arguments might be the only piece of information about a verb's argument structure that has to be taken care of in the syntax. This is not a necessary conclusion, however. Marantz\(^4\) has argued that external arguments are not arguments of their verbs at all. Within our current semantic framework, it is hard to see how to even make sense of such an idea. It is possible, however, to implement it within a theory that construes verb denotations along the lines of Donald Davidson.\(^5\) Kratzer shows how this move eliminates the need for a syntactic representation of the distinction between external and internal arguments.\(^6\) All information about a verb's argument structure, then, would be directly derivable from its denotation. There would be no syntactic theory of argument structure or linking. While this line of research might have promise, it goes way beyond what we can and should pursue in an introductory text.

Notes

1. This is similar to the method of "rule-by-rule interpretation" (so dubbed by Emmon Bach) which you find in Montague's papers, for example. Suppose we described the set of phrase structures that are interpretable by our current semantics by means of a context-free grammar. Then each of our semantic rules would correspond to exactly one of the phrase structure rules — hence the name "rule-by-rule". See E. Bach, "An Extension of Classical Transformational Grammar," in Problems of Linguistic Metathtory, Proceedings of the 1976 Conference at Michigan State University, 1976, pp. 183–224.


3. Note that we can't write "\(\lambda p \in D_1. [\lambda q \in D_1. p \land q]\)". "p" and "q" are variables for truth-values, so the "and" in "p and q" would have the same status as the "and" in "1 and 0", for example. Without further conventions, it is not clear what this would mean.

4. This has been argued for, e.g., by R. Kayne, Connectedness and Binary Branching (Dordrecht, Foris, 1984).


For a precise definition, see e.g., Partee et al., *Mathematical Methods*, pp. 443–4.


We might stress that this objection is directed at the specific formulation of the Θ-Criterion that we have quoted. It does not apply to all versions of the Θ-Criterion that you find in the literature. James Higginbotham, in particular, has proposed to replace the second part of the Θ-Criterion by the more flexible requirement that every thematic role must somehow be discharged. See J. Higginbotham, “On Semantics,” *Linguistic Inquiry*, 16 (1985), pp. 547–93; idem, “Elucidations of Meaning,” *Linguistics and Philosophy*, 12 (1989), pp. 465–517. In the latter he spells out a list of different discharge mechanisms (pp. 475ff.).

Recall our pedagogical decision to avoid ternary branching – hence the somewhat arbitrary right-branching structure.


This is actually not quite true. Our semantics doesn't force us to assume that there is no way to change the argument structure of a verb. It does imply, however, that if there are such changes at all, they have to be brought about by specific morphemes. We could have a morpheme of the following kind, for example:

$$\text{[passive]} = \lambda f \in D_{\text{ee,et,te}}. [\lambda x \in D_\varepsilon, \text{there is some } a \in D_\varepsilon \text{ such that } f(x)(a) = 1]$$

[passive] operates on transitive V- or V-denotations, and eliminates the highest argument. An analysis of this kind for the participle affix “-en” was proposed by D. Dowty, “Governed Transformations as Lexical Rules in a Montague Grammar,” *Linguistic Inquiry*, 9 (1978), pp. 393–426. We will not take a stand here as to
whether natural languages have items like passive. All we are interested in right now is to point out that our framework requires any change in argument structure to be brought about by some morpheme.

19 M. C. Baker, *Incorporation. A Theory of Grammatical Function Changing* (Chicago, University of Chicago Press, 1988), p. 46. Baker's condition includes the requirement that UTAH is a principle for Deep Structure representations. At the present stage of our investigation, we are not presupposing a multistratal syntax, so we left out the level-specification in Baker's original formulation of the UTAH. Baker's UTAH is a stronger version of Perlmutter and Postal's Universal Alignment Hypothesis (UAH), which states that "there exist principles of universal grammar which predict the initial relation borne by each nominal in a given clause from the meaning of the clause" (D. M. Perlmutter and P. M. Postal, "The 1-Advancement Exclusiveness Law," in D. M. Perlmutter and C. G. Rosen (eds), *Studies in Relational Grammar*, vol. 2 (Chicago, University of Chicago Press, 1984), pp. 81-125). Principles like UAH or UTAH are defended in a very sophisticated way in D. Pesetsky, *Zero Syntax, Experiencers and Cascades* (Cambridge, Mass., MIT Press, 1995), but have been argued to be untenable by Carol Rosen, "The Interface between Semantic Roles and Initial Grammatical Relations," in Perlmutter and Rosen (eds), *Relational Grammar*, pp. 38-77.


4 More of English: Nonverbal Predicates, Modifiers, Definite Descriptions

We ultimately want to determine appropriate types of meanings for lexical items of all syntactic categories and to predict how these meanings are composed in all sorts of syntactic constructions. In this chapter, we take some further small steps towards this goal. We continue with the Fregean working hypothesis that there are two types of basic, or saturated, meanings — namely, individuals and truth-values — and that all other meanings are functions that are somehow constructed out of these.

We will run into important general issues like vagueness, context-dependency, and presupposition, but will not stop to consider them seriously and systematically. We will also make quite a few concrete decisions without defending them against alternatives. The primary purpose of this chapter is to build a reasonably broad base of simple constructions for which we have at least a preliminary treatment, so that we are not too constrained in our choice of examples in the subsequent chapters. A secondary purpose is to give an overview of some of the basic issues regarding modification and presupposition.

4.1 Semantically vacuous words

Some lexical items are widely held to make no semantic contribution to the structures in which they occur. The standard example is of certain occurrences of prepositions, such as “of” in “proud of John” or “father of John”. Another plausible candidate is the copula “be” in predicative sentences like “John is rich”. We will also assume, at least for the time being, that the indefinite article “a” is vacuous when it occurs in predicate nominals such as “a cat” in “Kaline is a cat”.¹ We would want the following equalities, for example:
(1) \([\text{of John}] = [\text{John}]\)
\([\text{be rich}] = [\text{rich}]\)
\([\text{a cat}] = [\text{cat}]\)

There are various ways of making this come out. One way is to list semantically vacuous items in the lexicon as denoting the identity function of the appropriate type, for instance:

(2) \([\text{of}] = \lambda x \in D_e . x\)

(3) \([\text{be}] = \lambda f \in D_{<e,t} . f\)

(4) \([\text{a}] = \lambda f \in D_{<e,t} . f\)

In words: \([\text{of}]\) is that function which maps every individual in \(D_e\) to itself, and \([\text{be}]\) (\(= [\text{a}]\)) is that function which maps every function in \(D_{<e,t}\) to itself.

An even easier possibility is to assume that the semantic component simply “doesn’t see” such items. In other words, a structure that is really binary-branched may be treated as non-branching in the semantics: a branch occupied only by a vacuous item doesn’t count. The principle for nonbranching nodes then applies and passes up the meaning unchanged. Either way, we ensure the equalities in (1).

### 4.2 Nonverbal predicates

What we have assumed for verbs can be extended straightforwardly to adjectives, nouns, and prepositions. Just as intransitive verbs denote functions from individuals to truth-values, so do many nouns and adjectives, for instance:

(1) \([\text{cat}] = \lambda x \in D_e . x \text{ is a cat}\)

(2) \([\text{gray}] = \lambda x \in D_e . x \text{ is gray}\)

Among prepositions, intransitive (1-place, monadic) ones are the exception rather than the rule, but there are some candidates:

(3) \([\text{out}] = \lambda x \in D_e . x \text{ is not in } x\text{'s home}\)
Each of these categories also has transitive (2-place, dyadic) members, whose extensions are just like those of transitive verbs – for example, “part” in “part of Europe” and “proud” in “proud of John” (note the vacuity of “of”), and all run-of-the-mill prepositions:

(4) \[[\text{part}] = \lambda x \in D_e . [\lambda y \in D_e . y \text{ is part of } x]\]

(5) \[[\text{fond}] = \lambda x \in D_e . [\lambda y \in D_e . y \text{ is fond of } x]\]

(6) \[[\text{in}] = \lambda x \in D_e . [\lambda y \in D_e . y \text{ is in } x]\]

These lexical entries allow us to calculate appropriate denotations for the phrases “part of Europe”, “fond of Joe”, “in Texas”, by means of our composition principle Functional Application (FA), for instance:

By FA: \[[\text{in Texas}] = [[\text{in}}][[\text{Texas}]]\]
By lexical entry for Texas: \[[\text{Texas}] = \text{Texas}\]
Hence: \[[\text{in Texas}] = [[\text{in}}](\text{Texas})\]
By lexical entry for in: \[[\text{in}}](\text{Texas}) = \lambda x \in D_e . [\lambda y \in D_e . y \text{ is in } x]\](\text{Texas}) = \lambda y \in D_e . y \text{ is in Texas.}
Hence: \[[\text{in Texas}] = \lambda y \in D_e . y \text{ is in Texas.}\]

We will disregard the case of ditransitive (3-place, triadic) predicates, though there are presumably some analogs to verbs like “give” and “introduce” in other syntactic categories.

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**Exercise**

Calculate the truth-conditions for at least one of the sentences “Joe is in Texas”, “Joe is fond of Kaline”, and “Kaline is a cat”.

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**4.3 Predicates as restrictive modifiers**

It has often been observed that prepositional phrases (PPs) may appear inside NPs in three distinct semantic roles: as arguments, as restrictive modifiers, or as nonrestrictive modifiers. 2 Typical examples of each type are the following:
We have already dealt with the first case: PPs that are arguments are headed by vacuous prepositions, thus they have the same denotations as the NPs they contain (individuals), and they are arguments of 2-place (relational) nouns such as “part”.

About the third case, we have next to nothing to say, and we mentioned it only to guard against confusing nonrestrictive modifiers with restrictive ones. The basic intuition that most authors have expressed about the semantics of nonrestrictive modification is that nonrestrictive modifiers are not semantically composed at all with the phrases they modify. Rather, they have the status of separate sentences which serve to make side-remarks of some kind. For example, the meaning of (4) is not unlike that of (5).

(4) It is surprising that Susan, from Nebraska, finds it cold in here.

(5) It is surprising that Susan finds it cold in here. Note that she is from Nebraska.

This makes it reasonable to assume that at the level at which our semantic rules apply, the nonrestrictive modifier isn’t part of the structure at all, so the question of how its denotation should be composed with that of its modifier doesn’t arise in the first place. This said, we will concentrate on restrictive modifiers, which are our primary topic in this section.

While the distinction between arguments and restrictive modifiers is notoriously difficult to make in practice, the basic semantic intuition behind it is simple: Arguments reduce the adicity of the noun they combine with; modifiers leave it unchanged.

What it means to reduce the adicity of the noun is illustrated by the example we have treated: “part” is a 2-place predicate, while the result of combining it with “of Europe” to form “part of Europe” is a 1-place predicate. Restrictive modifiers, by contrast, are characterized by the fact that they leave the semantic type, including the adicity, of the modifier completely unchanged. Intuitively, “city in Texas” has the same kind of extension as “city”: namely, the characteristic function of a set. More specifically, if \([\text{city in Texas}]\) is the characteristic function of a set A and \([\text{city}]\) is the characteristic function of a set B, then A is a subset of B: namely, that subset which results by intersecting B with the set of things in Texas. Consider the truth-conditions of “Lockhart is a city in Texas” to confirm this intuition.
Now we have already proposed an analysis of “in” above (motivated at the
time by occurrences of this preposition in simple copula + PP phrases), under
which “in Texas” has a denotation of type \(<e,t>\). Given this decision and our
current inventory of composition rules, we predict “city in Texas” to be
uninterpretable. As a branching structure with daughters “city” and “in Texas”,
we should be interpreting it by applying either \([city]\) to \([in Texas]\) or vice versa.
But neither is possible, as both are functions of type \(<e,t>\).

What shall we do about this problem? There are two directions we could take:
either revise out lexical semantics for (some of) the ingredients, or else stipulate
a new composition rule. We will entertain both options, beginning with the second.

4.3.1 A new composition rule

Here is a composition principle which is tailored to the situation at hand.

(6) **Predicate Modification (PM)**

If \(\alpha\) is a branching node, \(\{\beta, \gamma\}\) is the set of \(\alpha\)'s daughters, and \([\beta]\) and \([\gamma]\)
are both in \(D_{<e,t>}\), then

\([\alpha] = \lambda x \in D_e . [\beta](x) = [\gamma](x) = 1.5\)

Applied to “city in Texas”, (6) gives the desired result:

(7) \([city in Texas]\)

\[= \lambda x \in D_e . [city](x) = [in Texas](x) = 1\]

\[= \lambda x \in D_e . x \text{ is a city and } x \text{ is in Texas.}\]

(In the last step, we used the lexical entries and the result of a previous calculation
from section 4.2 above.)

PM is general enough to cover not just PPs modifying nouns, but likewise
adjective phrases (APs), whether to the right or the left of a noun, and also
stacked modifiers in unlimited numbers. For instance, we can now predict cor-
rectly the truth-conditions of (8).

(8) Kaline is a gray cat in Texas fond of Joe.

To treat this example, we must, of course, impose some binary-branching hier-
archy among the three modifiers “gray”, “in Texas”, and “fond of Joe”. Given
the nature of the semantic operation performed by PM, it so happens that all
our different choices in this regard yield logically equivalent results. This is as
it should be. The syntax of English (to the best of our knowledge) does not
determine a unique parse for this sentence, but it is nevertheless perceived as
truth-conditionally unambiguous.
Exercise

Calculate the truth-conditions for (8), given one possible syntactic parse.

The operation performed by PM has also been called "intersective modification", because if we look at the sets instead of at their characteristic functions, it amounts to set-theoretic intersection. "Conjunctive composition" would be another natural name, highlighting the connection with the semantics of "and".\(^6\) (Notice that "city in Texas" receives exactly the meaning we derive for "city and in Texas", using the predicate-coordinating "and" we defined in section 3.4.)

Predicate modification is a genuinely new semantic composition principle on our list. It is obviously not functional application. If it is really needed, there is more to semantic composition than Frege's Conjecture. Are we forced to this conclusion?

4.3.2 Modification as functional application

As we have already mentioned, the alternative to PM is to explore revised lexical entries for the words that may head modifiers. Suppose we insist that \([\text{city}]\) and \([\text{in Texas}]\) combine by Functional Application after all. If we keep \([\text{city}]\) as before, with type \(<\text{e},\text{t}>\), then \([\text{in Texas}]\) will have to be of type \(<<\text{e},\text{t}>,<\text{e},\text{t}>>\). This in turn requires a new semantics for at least one of "in" and "Texas". If we keep the assumption that \([\text{Texas}] = \text{Texas} \in \text{D}_\text{e}\), we must reinterpret "in". It must now denote a function of type \(<\text{e},<<\text{e},\text{t}>,<\text{e},\text{t}>>\). Which such function? Well, we know that we want to be able to derive equations like these:

\[(9) \quad [\text{in}](\text{Texas})(\lambda x \in \text{D}_\text{e} . x \text{ is a city}) = \lambda x \in \text{D}_\text{e} . x \text{ is a city and x is in Texas.} \]
\[= [\text{in}](\text{Building 20})(\lambda x \in \text{D}_\text{e} . x \text{ is a room}) = \lambda x \in \text{D}_\text{e} . x \text{ is a room and x is in Building 20.} \]

The generalization appears to be that, for any individual \(y \in \text{D}_\text{e}\) and any function \(f \in \text{D}_{<\text{e},\text{t}>>}\), \([\text{in}](y)(f) = \lambda x \in \text{D}_\text{e} . f(x) = 1 \text{ and x is in y.}\) This determines directly the desired definition for the function \([\text{in}]\):

\[(10) \quad [\text{in}] = \lambda y \in \text{D}_\text{e} . [\lambda f \in \text{D}_{<\text{e},\text{t}>>} . [\lambda x \in \text{D}_\text{e} . f(x) = 1 \text{ and x is in y}]] \]

By similar reasoning, we can determine an entry for the adjective "gray" that permits phrases like "gray cat" to be interpreted by Functional Application:

\[(11) \quad [\text{gray}] = \lambda f \in \text{D}_{<\text{e},\text{t}>>} . [\lambda x \in \text{D}_\text{e} . f(x) = 1 \text{ and x is gray}] \]
Exercise

Calculate the truth-conditions for (8), given one possible syntactic parse. This time, use FA instead of PM.

By systematically revising the entries of all adjectives and prepositions, we are able to interpret all phrases containing a noun with one or more modifiers in them by means of Functional Application alone, and so we can eliminate Predicate Modification from the theory.

But there is a trade-off. What happens now when an AP or PP stands as a (maximal) predicate by itself, as in "Julius is gray" or "Julius is in Amherst"? If APs and PPs denote functions of type $<e,t>,<e,t>$, these sentences are prima facie uninterpretable. We could try to solve the problem by assigning a suitable denotation to the copula "be".

Exercise

Define such a denotation. There are two distinct solutions.

But then this same "be" would not be interpretable in sentences with a nominal predicate, for example, "Julius is a cat", since NPs like "a cat" are still of type $<e,t>$. So the copula would have to be ambiguous between vacuous and nonvacuous occurrences.

Another solution is to assume a systematic lexical ambiguity in all adjectives and prepositions. Each has both the initially assumed type $<e,t>$ (or $<e,<e,t>$) meaning and the new type $<e,t>,<et>$ (or $<e,<e,t>,<et>$) denotation. The syntax may freely generate both homonyms in all the same places, but the Principle of Interpretability will allow only one in any given environment.

Exercise

It would not be adequate to list the two readings of each preposition or adjective separately in the lexicon, as if they had to be learned individually. Evidently, there is a systematic relation between the two readings, which makes one predicatable, given the other. So we would want to list only one in the
lexicon, and derive the other by means of a general "lexical rule". Spell out two versions of this proposal. For the first version, assume that the homonyms with the simpler types \(<e,t>\) or \(<e,<e,t>>\) are basic and listed in the lexicon. For the second version, assume that the more complicated types \(<<e,t>,<et>>\) or \(<e,<<e,t>,<et>>>\) are the ones of the basic, individually listed items. Your task is to formulate the appropriate lexical rules for either version. That is, you have to specify general recipes that map arbitrary denotations of the basic type to secondary denotations of the appropriate nonbasic type. Rules of this kind are called "type-shifting rules".

Yet another solution is to posit a certain amount of non-overt structure in VPs of the surface form "be" + AP or "be" + PP. Perhaps these VPs contain an invisible predicate that the AP or PP modifies, something like a zero equivalent of a bland noun like "thing" or "individual".

We will not explore or try to evaluate these options further. All we wanted to show here is that the elimination of the Predicate Modification rule is not without its price. As matters stand, it does not look entirely unreasonable if we decide to adopt PM after all.

### 4.3.3 Evidence from nonintersective adjectives?

Both analyses of adjectives that we have entertained so far predict that the following pair of sentences are logically equivalent:

\[(12)\] 
\[(a)\] Julius is a gray cat.  
\[(b)\] Julius is gray and Julius is a cat.

In the analysis that uses PM, the equivalence follows directly from the content of the rule. We can prove it without using any specific information about the meanings of the lexical items, except the information about their semantic type. In the analysis that relies on Functional Application alone, the equivalence follows from the lexical meaning of the adjective "gray". On this alternative, mere inspection of the types of the words and the applicable composition principles does not suffice to prove it.

Since the equivalence in (12) is indeed intuitively valid, both analyses make the correct prediction, albeit in different ways. But it has often been noted\(^8\) that analogous equivalences do not obtain for many other adjectives, and it seems that this fact might have some bearing on the choice we have been contemplating.

Consider adjectives like "large" and "small". One might truthfully assert that a small elephant is still a very large animal. So it is intuitively possible for (13) to be true while (14) is false.
(13) Jumbo is a small elephant.

(14) Jumbo is a small animal.

This shows that (13) does not intuitively entail (14). But “Jumbo is an elephant” does entail “Jumbo is an animal”, and given the meaning of “and”, this implies that (15) entails (16).

(15) Jumbo is small and Jumbo is an elephant.

(16) Jumbo is small and Jumbo is an animal.

So (13) and (14) cannot be equivalent to (15) and (16) respectively. At least one, perhaps both, of these equivalences must be denied, or else we falsely predict that (13) entails (14).

Now we have seen that if adjectives are of type <e,t> and combine with their modifiers by PM, the equivalences in question follow regardless of specific lexical meaning. We have to conclude, therefore, that it is not possible to define an adequate type <e,t> meaning for “small”.

But if the meaning of “small” is of type <<e,t>,<e,t>>, it seems that we have a chance of defining it appropriately, so that the inference from (13) to (14) does not go through. We do not want a lexical entry analogous to (11) above, of course (that is, we don’t want to define \([\text{small}] = \lambda f \in D_{e,t} . [\lambda x . f(x) = 1 \text{ and } x \text{ is small}]\), but fortunately there are other functions of type <<e,t>,<e,t>> to consider. Here is a proposal that seems to reflect the relevant intuition about why small elephants are not small animals.

(17) \([\text{small}] = \lambda f \in D_{e,t} . [\lambda x \in D_e . f(x) = 1 \text{ and the size of } x \text{ is below the average size of the elements of } \{ y : f(y) = 1 \}]\)

One might quibble with some of the details, but (17) is definitely on the right track towards an explanation of why (13) fails to imply (14). According to (17), (13) asserts that Jumbo is an elephant and Jumbo’s size is below the average elephant size. The set of all animals contains mostly individuals that are smaller than any elephants, so the average animal size is much lower than the average elephant size, and it is easy for Jumbo’s size to fall between the two. This is how (14) can be false when (13) is true.

So it seems that we should allow at least some adjectives to denote functions of type <<e,t>,<e,t>>. “Small”, we have seen, cannot possibly be interpreted with the lower type <e,t>, and this is just one of many examples for which the same kind of argument can be made. These adjectives are often called
nonintersective. In the adjective inventories of natural languages, they are apparently in the majority. “Intersective” adjectives – that is, those which validate equivalences like (12), like our initial example “gray” – represent the exception rather than the rule.

Are there any implications of this discussion for the analysis of the intersective adjectives? Not as far as we have seen up to now. Adding type $$<e,t>,<e,t>$$ adjectives like “small” to the lexicon does not seem to cost us anything extra, whether we add them to a grammar that also has type $$<e,t>$$ adjectives and a rule of PM, or to one where all adjectives have type $$<e,t>,<e,t>$$ and PM is absent. Notice that our theory does not place any premium per se on having a uniform semantic type for all members of a given syntactic category. So the fact that on the PM option, some adjectives have type $$<e,t>$$ and others type $$<e,t>,<e,t>$$ is not in itself a reason to disprefer it.

The picture changes if we recall some of the trade-offs we thought to be tied to the elimination of PM. For instance, we observed that type $$<e,t>,<e,t>$$ meanings are less straightforwardly interpretable than type $$<e,t>$$ meanings in predicative occurrences of APs, as in “Julius is gray”. But it turns out that adjectives like “small”, which we have seen cannot have lexical entries of type $$<e,t>$$, also occur as seemingly complete predicates:

(18) Jumbo is small.

What do we make of this? We seem to need one of the mechanisms that we considered in section 4.3.2 as a replacement for PM after all, whether we have PM or not. For instance, we may posit a zero modifiee in the syntax of (18), or assume a lexical rule that routinely produces secondary type $$<e,t>$$ meanings from type $$<e,t>,<e,t>$$ inputs. (See exercise above.) But this tips the balance, and a theory without PM begins to look more parsimonious over all.

Or does it? Let’s consider the whole case based on “small” a bit more carefully before we jump to conclusions. Is a type $$<e,t>,<e,t>$$ analysis as in the entry (17) really forced on us by the data we have considered?

An elementary observation about adjectives like “small” that we have so far left aside is that they are vague and heavily context-dependent. In suitable specific discourse settings, people may have truth-value judgments about utterances containing “small” which are firm and uniform across speakers. But if we try to generalize over all felicitous uses of a given “small” sentence, we find that objects of practically any size can count as “small” and also as “not small”. It is important to note that this vagueness and context-dependency are not limited to predicative occurrences like the one in (18). It remains even when there is a noun that “small” modifies. Consider Jumbo again. When you first read sentence (13), which described him as “a small elephant”, in the context in which we presented it above, you spontaneously interpreted it in the manner we have
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described: namely, as true only if Jumbo is smaller than the average elephant. But imagine we had first introduced a scenario populated with an army of monsters like King Kong. We might then have said something like: “Jumbo doesn’t have a chance; he’s only a small elephant”, and this could have been true even if Jumbo were as large as or even larger than most other elephants.

So the meaning for “small” we codified in (17) represents at best a sort of default that applies when phrases of the form “small N” are interpreted more or less out of context. The contribution of the modified noun’s meaning to the meaning of the whole phrase is not in general this mechanical. Perhaps the basic generalization about “small” is that it means “of a size less than the contextually salient standard”. How the contextually salient standard is established for each given utterance of the word “small” is a complex affair. Previous discourse and the nonlinguistic circumstances of the utterance play a role. The mention of the word “elephant” in the immediate vicinity of the adjective draws attention to the elephant stereotype, including the stereotypical elephant size. Perhaps this is the whole reason why the average size of elephants happens to be the most salient standard in most situations where somebody utters the phrase “small elephant”. In other words, the contribution of the modified noun may be rather indirect and mediated by the context.

If this picture is correct, a type <e,t> interpretation for “small” may be viable after all. The lexical entry might say essentially the following:

(19)  \( [\text{small}] = \lambda x \in D_e. x\text{'s size is below }c, \text{ where }c\text{ is the size standard made salient by the utterance context.} \)

“Small elephant”, as it occurs in (13), could then be interpreted by PM after all. It would receive the truth-conditions in (17) just in case the context of utterance does not supply a more salient size standard than the average size of elephants. To explain the intuition that (13) does not entail (14), we would assume that utterance contexts change quickly. We must acknowledge that it is possible, in fact highly natural, for an utterance of (14) which follows right after (13) to change the prevailing size standard from average elephant size to average animal size. We do predict that (13) entails (14) if the context for both is the same. But this prediction may be compatible with the evidence, if we can tell a plausible story about why the context will automatically change whenever an utterance of (14) follows one of (13).

We have evidently scratched only the surface of a complex of important issues here. Vagueness and context-dependency have been studied quite carefully by philosophers and linguists within the general framework of natural language semantics that we are presenting here. But the results of their work, both substantive and technical, are largely beyond the scope of this text. Our present conclusion must therefore remain open for reconsideration. But for the purposes
of this book, we assume that type <e,t> entries for vague adjectives like “small” are viable, and we continue to work with the Predicate Modification rule.

We should briefly mention that there are a few other types of apparently nonintersective adjectives, with a behavior rather different from “small”. One group is represented by “former” and “alleged”. Clearly, if John is a former teacher, it does not follow that John is former and John is a teacher. The second conjunct is clearly false in this case, and the first is not even grammatical. The latter fact suggests that “former” does not have a meaning of type <e,t>, and moreover that it is not, after all, advisable to make a zero modifier available in the syntax of all VPs of the form “be” + AP. At least, these are the obvious conclusions to draw if we want to predict that “John is former” is plainly uninterpretable. In this respect, “former”, unlike “small”, shows just the distribution that we would expect of an adjective of type <<e,t>,<e,t>>. But unfortunately its meaning can be shown not to be of this type. The reasoning goes as follows: if [[former]] were a function of type <<e,t>,<e,t>>>, then for any two nouns α and β, the following would hold:

\[ (20) \quad \text{If } \alpha = \beta, \text{ then } [[\text{former } \alpha]] = [[\text{former } \beta]]. \]

The reason for (20) is that “former α” (and likewise “former β”) could only be interpreted by Functional Application: that is, [[former α]] = [[former]]([[α]]). Since [[α]] is by assumption the same as [[β]], we must have [[former]][[[α]]] = [[former]][[[β]]]. But (20) implies counterintuitive predictions. Suppose Bill’s lovers happen to be exactly the tenants of 13 Green Street. So for any \( x \in D \), [[lover of Bill’s]](x) = 1 iff [[tenant of 13 Green St]](x) = 1. By the mathematical definition of a function, this means that [[lover of Bill’s]] = [[tenant of 13 Green St]]. With (20), it then follows that, if “John is a former lover of Bill’s” is true, then so is “John is a former tenant of 13 Green Street”. But intuitively, the situation described is entirely compatible with the first of these being true and the second false. So we conclude that [[former]] is not of type <<e,t>,<e,t>>.

Apart from these negative conclusions, our present framework doesn’t enable us to say anything precise. Very roughly, a successful analysis of “former” presupposes a general account of the time parameter in predicates. We must first adapt our semantics to the elementary fact that, for example, “John is a teacher” may be true in 1970 and false in 1980. Our current entries for predicates only make sense if we either disregard any change over time, or tacitly agree that we are considering a certain fixed point in time. The proper treatment of implicit time reference and temporal quantification requires an intensional semantics. An intensional semantics will be introduced in chapter 12, but even there, we will neglect temporal dependencies. Accordingly, we cannot fully answer the question of what the existence of adjectives like “former” implies for the treatment of adjectives and modification in general.
Many other loose ends have been left in this brief introduction to modifiers. The astute reader may have noticed, for instance, that we have said nothing about PPs modifying verbs rather than nouns. A few cautiously chosen examples of this type happen to mean what we predict them to mean: for example, “Julius is sleeping on the couch” does seem to be true just in case Julius is sleeping and is on the couch. (Exercise: Show that this is what we predict.) Many other examples that readily come to mind, however, do not have the predicted meanings at all. Consider “John wrote on the blackboard”, “Mary put the book on the table”, or “Max tossed the salad in the bowl”, to name just a few. There are plenty of good grounds here to suspect that our present account of modifiers is very far from how modification really works in natural language. On the other hand, it might turn out that the problems lie elsewhere – for instance, in our simplistic assumptions about verb meaning and VP structure. Indeed, current research overwhelmingly points to the latter conclusion, and many prima facie counterexamples emerge as cases of intersective modification after all.10

4.4 The definite article

We have proposed that common nouns like “cat” denote the characteristic functions of sets of individuals. What does this imply for the semantic analysis of determiners? We will defer the general version of this question to chapter 6. Right here, we will look at only one determiner: namely, the definite article.11

4.4.1 A lexical entry inspired by Frege

The basic intuition about phrases of the form “the NP” is that they denote individuals, just like proper names. Had it not been for Bertrand Russell’s famous claim to the contrary, few people would think otherwise. Frege, for one, thought it obvious: “let us start, e.g., with the expression ‘the capital of the German Empire.’ This obviously takes the place of a proper name, and has as its reference an object.”12 Hence his practice of referring to definite descriptions as “compound proper names”.

If you read on in that particular passage, Frege imposes a rather odd syntactic analysis on his example. Instead of dividing it into the constituents “the” and “capital of the German Empire”, he splits it up into “the capital of” and “the German Empire”. He then proceeds to analyze “the capital of” as denoting a function from objects to objects (our type <e,e>). Elsewhere, however, he treats a similar example as follows:
"the negative square root of 4". We have here a case in which out of a concept-expression a compound proper name is formed with the help of the definite article in the singular, which is at any rate permissible when one and only one object falls under the concept.13

We follow Frege in this second syntactic and semantic analysis. Notice that by "concept-expression", he means an expression whose meaning is of type \(<e,t>\). In his example, that's the NP "negative square root of 4", which indeed receives a meaning of that type if we analyze it along the lines of the previous sections of this chapter. ("Square root", it seems, is a transitive noun like "part", with a meaning of type \(<e,\langle e,t\rangle>\). "Of" is vacuous, \([\text{square root}]\) applies to 4 by Functional Application, and the result of that composes with \([\text{negative}]\) under Predicate Modification.)

The determiner "the", then, denotes a function with arguments in \(D_{<e,t>}\) and values in \(D_e\). For instance, \([\text{the}]\) applied to the function \([\text{negative square root of 4}]\) yields the number \(-2\). \([\text{the}]\) applied to \([\text{president of the USA}]\) yields Clinton at the time of writing. \([\text{the}]\) applied to \([\text{opera by Beethoven}]\) yields Fidelio. The generalization that emerges is (1).

\[(1)\quad \text{For any } f \in D_{<e,t>} \text{ such that there is exactly one } x \text{ for which } f(x) = 1, \quad [\text{the}]_\langle f \rangle = \text{the unique } x \text{ for which } f(x) = 1.\]

What about functions \(f\) which do not map exactly one individual to 1? What is \([\text{the}]_\langle f \rangle\) for one of those?

Let's examine our intuitions on this matter. What are the objects denoted by the following definites?

\[(2)\quad \text{the escalator in South College}\]
\[(3)\quad \text{the stairway in South College}\]

You should know that South College has no escalator and more than one stairway. Once we are aware of this, we are hard pressed to say which objects (2) and (3) denote. The only natural answer is that neither of these phrases denotes any object at all. Let's go ahead and implement precisely this simple-minded intuition in our lexical entry for "the".

What we are aiming to predict is that (2) and (3) have no semantic value. In other words, there is no such thing as \([\text{the escalator in South College}]\) or \([\text{the stairway in South College}]\). The reason has to be that the functions \([\text{escalator in South College}]\) and \([\text{stairway in South College}]\) are not in the domain of \([\text{the}]\). If they are not in the domain of \([\text{the}]\), then \([\text{the}]\) can't apply to them, and this means that we cannot apply FA to calculate a semantic value for the DP-nodes in (2) or (3). The generalization that emerges regarding the domain of \([\text{the}]\) is this:
The domain of \([\text{the}]\) contains just those functions \(f \in D_{<e,t>e}\) which satisfy the condition that there is exactly one \(x\) for which \(f(x) = 1\).

Putting (1) and (4) together, we can now formulate our lexical entry for "the":

\[
[\text{the}] = \\
\lambda f : f \in D_{<e,t>e} \text{ and there is exactly one } x \text{ such that } f(x) = 1. \\
\text{the unique } y \text{ such that } f(y) = 1.
\]

This is a bit of an unwieldy \(\lambda\)-term, but if you apply it to the examples above, you can see that it describes the function we were trying to define.

Before we end this subsection, let's dispose of a technical matter. What is the semantic type of \([\text{the}]\)? In the strict sense of our definitions so far, it actually has none. To say that its type is \(<<e,t>e>\) would mean that it is a function from \(D_{<t,e>}\) to \(D_e\). But "from \(D_{<e,t>}\)" means "with domain \(D_{<e,t>}\)", and we have just seen that the domain of \([\text{the}]\) is not \(D_{<e,t>}\) but only a subset thereof. At this point, we find ourselves with an inconvenient terminology, and we will simply change it. We henceforth define \(D_{<e,t>} \) (for any types \(\sigma, \tau\)) as the set of all partial functions from \(D_\sigma\) to \(D_\tau\). "Partial function from" is defined as follows:

\[
(6) \text{A partial function from } A \text{ to } B \text{ is a function from a subset of } A \text{ to } B.
\]

(When we emphatically mean "function from" rather than "partial function from", we will sometimes say "total function from".) With these new definitions, we can now say that \([\text{the}]\) is in \(D_{<e,t>e}>\), or that its type is \(<<e,t>e>\).

### 4.4.2 Partial denotations and the distinction between presupposition and assertion

When a tree contains a lexical item that denotes a partial function, this may cause the tree to wind up without a semantic value. We have already seen this happen in (2) and (3) above. In larger structures, there are repercussions all the way up the tree. For instance, the following sentences are predicted not to have any semantic values, neither 1 nor 0 nor anything else.

\[
(7) \text{The stairway in South College is dirty.}
\]

\[
(8) \text{John is on the escalator in South College.}
\]

This is a direct consequence of our set of composition principles. The only principle that could potentially provide a semantic value for the branching node
above the definite description (that is, the S-node in (7) and the PP-node in (8)) is Functional Application, repeated here from chapter 3.

(9) **Functional Application (FA)**

If $\alpha$ is a branching node and $\{\beta, \gamma\}$ is the set of $\alpha$'s daughters, then $\alpha$ is in the domain of $\llbracket \cdot \rrbracket$ if both $\beta$ and $\gamma$ are and $\llbracket \gamma \rrbracket$ is in the domain of $\llbracket \beta \rrbracket$. In this case, $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket(\llbracket \gamma \rrbracket)$.

But, as this formulation makes plain, you can't apply the denotation of one daughter to that of the other unless both daughters have denotations. So FA can't apply (and no other composition principle even came close). By the same reasoning, no semantic values can be obtained for any higher nodes that indirectly dominate a denotationless definite description like (2) or (3).

---

**Exercise**

Consider an example with one definite description embedded in another:

(i) The killer of the black cat escaped.

(a) Draw an interpretable syntactic structure for (i).

(b) Describe three possible states of affairs:

- one where (i) is false;
- another one where (i) lacks a truth-value because "the black cat" has no extension;
- a third one where (i) is also without a truth-value, but this time because "the killer of the black cat" lacks a denotation. (Assume for this third scenario that "the black cat" does have a denotation.)

Are the empirical predictions that are implied by our current semantic component correct? Consider what we predict about (8): If you know English, and if you furthermore know that there is no escalator in South College, then you know that:

(a) the sentence "John is on the escalator in South College" is not true, and

(b) the sentence "John is on the escalator in South College" is not false.
Part (a) of this prediction is unobjectionable. But in apparent disagreement with (b), many informants will spontaneously classify the assertion "John is on the escalator in South College" as false.

Does this mean that our Fregean semantics for the definite article has proved empirically inadequate and must be abandoned? Many philosophers and linguists have drawn this conclusion. It is a reasonable conclusion, but it is not inescapable. An alternative response is to reconsider the straightforward identification which we have assumed so far between the semantic values 1 and 0 and the pre-theoretical notions of truth and falsity. Might we perhaps reconcile the present semantic analysis with the empirical evidence if we posit a somewhat more indirect correspondence between the truth-values of our semantic theory and the intuitions that people report in truth-value judgment tasks?

As a first step in this direction, let us propose that the colloquial term "false" covers both truth-value-less sentences and those that are false in the technical sense of denoting 0. In other words, the technical terms of our theory translate into pre-theoretical terms as follows:

\[
\begin{align*}
\phi \text{ is true} & \quad \Leftrightarrow \quad \left\{ \begin{array}{l}
\phi = 1 \\
\phi \text{ has no semantic value}
\end{array} \right. \\
\phi \text{ is false} & \quad \Leftrightarrow \quad \phi = 0
\end{align*}
\]

This stipulation makes the predictions of our semantics consistent with the data we reported above: for example, with the fact that informants who are told that South College contains no escalator and are asked to decide whether (8) is true or false will choose "false".

Mere compatibility with the data is not all we are aiming for, of course. To justify our choice over competing theories, in particular those that make only a 2-way distinction between true and false, we have to show that the additional distinction between two sources of intuitive falsity does some useful work. For instance, we might argue that it helps us to explain certain other manifestations of semantic competence, which can be observed when we move beyond simple truth-value judgment tasks and elicit subtler intuitions.

Indeed, this kind of argument has been offered. Specifically, it has been argued that the technical distinction between lacking a value and denoting 0 can be systematically related to an intuitive distinction: namely, the distinction between what is asserted and what is presupposed. For an illustration, consider the following three sentences.

(10) (a) John is absent again today.
     (b) Today is not the first time that John is absent.
     (c) John is absent today, and that has happened before.
All three of these sentences somehow express the speaker's belief that John is absent today and has been absent at least once before. But they are not simply interchangeable. If you are talking to somebody with whom you already share the information that John has been absent in the past, but who doesn't yet know about today, (10a) is a natural choice, but (10b) is not. If you are talking to somebody who already knows that John is absent today, but knows nothing about his past history, then (10b) is natural, whereas (10a) is not. And if your audience knows nothing at all about John's past or present attendance, the most natural choice is (10c). We accordingly say that (10a) presupposes that John was absent before and asserts that he is absent today. With (10b), it's the other way round: this sentence presupposes that John is absent today and asserts that he was absent before. Finally, (10c) asserts that John both is absent today and was absent before, without presupposing anything.

The ability to discriminate between (10a), (10b), and (10c) and decide which is most appropriate for an audience with a given state of information is clearly part of understanding English. We would therefore like to capture it somehow in our semantics. The hypothesis presently under consideration is that a semantic theory equipped to distinguish two kinds of non-true sentences is better suited to accomplish this than one that isn't. The concrete proposal is that $\phi$ having no value represents the case that $\phi$ has a false presupposition, and $[\phi] = 0$ means that $\phi$ does not presuppose anything false but makes a false assertion. (The third case, $[\phi] = 1$, thus has to mean that both what $\phi$ presupposes and what it asserts are true.)

Let's return to the definite article. The analysis we presented above (following Frege) may be called a "presuppositional" analysis of "the". In light of what we have said in the present subsection, it predicts that a sentence like (8) ("John is on the escalator in South College") would be used most naturally by a speaker who assumes that her audience knows that there is a unique escalator in South College, but doesn't know about John's whereabouts. This seems basically right. Minimal pairs like (11a, b) below point in the same direction.

(11)  
(a) There will be one mid-term, which will be on November 21st.  
(b) The mid-term will be on November 21st.

If the topic of mid-term exams for this course hasn't come up yet at all, (11a) is the natural choice; (11b) is fully appropriate only when the audience is already aware that there will be one mid-term. Neither sentence can be true unless there will be a unique mid-term, but in (11a) this is part of the assertion, while in (11b) it is presupposed. Our semantic analysis predicts this: If there isn't a unique mid-term, then (11b) has no truth-value at all. (In contrast with (11a), which should denote 0 in this case. We say "should", because we have yet to develop an analysis of the words and constructions in (11a) that actually predicts this.)
We have made an attempt here to provide some initial evidence in favor of our presuppositional analysis. This is not meant to be a compelling justification of this analysis, however. There are many well-known objections which we have not even mentioned, much less addressed. For the time being (especially for the next chapter), it is convenient to be able to assume some concrete analysis of the definite article, even if it turns out to be only a crude approximation. This will allow us to get our analysis of other parts of English off the ground.

Exercise 1

Assume that sentence (i) has the structure (i') at the level at which it is interpreted.

(i) John doesn't use the escalator in South College.

(i')

\[
\begin{array}{c}
S \\
\text{not} \\
S \\
\text{John} \\
\text{VP} \\
\text{use} \\
\text{DP} \\
\text{the} \\
\text{NP} \\
\text{escalator in South College}
\end{array}
\]

In (i'), the determiner "the" is assumed to be a D (determiner, not to be confused with the domain of individuals D, also referred to as "D_n") heading a DP (determiner phrase). The structure (i') differs from the surface structure of (i) with respect to the position of the subject "John". We might assume that (i') is the Deep Structure of (i), where the subject appears below the negation before movement to a higher position. Or else we could assume that (i') is a Logical Form representation that results from reconstruction of the raised subject into its original position.

What does our current theory predict about this sentence? Are the predictions empirically correct? Construct suitable examples and scenarios that help
Exercise 2

Look at the following scenario:

(← LEFT)  🍎  🍎  🍎  🍎  🍎  🍎 (RIGHT →)

and consider the following three definite descriptions:

(i) the leftmost apple in the row
(ii) the leftmost dark apple in the row
(iii) the apple that is both leftmost in the row and dark

(a) In your intuitive judgment, which individual, if any, does each of these definite descriptions refer to in this situation?
(b) What predictions concerning the denotations of the definite descriptions (i), (ii), and (iii) would follow from an analysis that treats adjectives like "leftmost" as 1-place predicates?
(c) Specify a more adequate denotation for attributive uses of "leftmost".
(d) In a compositional fashion, compute the denotation of the definite description "the leftmost dark apple in the row", given the above scenario. For the purpose of this computation, take "apple in the row" as an unanalyzed predicate. That is, you don't have to worry about the PP "in the row".

4.4.3 Uniqueness and utterance context

Frege's uniqueness presupposition has often been objected to as an idealization that does not really fit the definite singular article in English. We frequently say things like "the door is locked" or "the cat wants to be fed", yet we don't believe that there is just one door and just one cat in the world, and nobody that hears us speak this way will attribute such beliefs to us either.

There are a number of different responses to this objection, but this is not the occasion to give them serious consideration. Somehow or other, we have to
concede that “the cat” doesn’t denote the unique cat that there is in the whole world, but rather denotes, on each occasion on which it is uttered, the unique cat among those individuals that are under consideration on this utterance occasion. Our best attempt for now at making this explicit is in the following revised lexical entry:

\[(5') \quad \text{[the]} = \\lambda f : f \in D_{cte} \text{ and there is exactly one } x \in C \text{ such that } f(x) = 1.\]

the unique \( y \in C \text{ such that } f(y) = 1, \)

where \( C \) is a contextually salient subset of \( D \).

Once again, our lack of a serious account of context-dependency prevents us from stating this more precisely. Below, we will assume something like \((5')\) in informal discussion, but will abstract away from context-dependency and use \((5)\) in our calculations.

### 4.4.4 Presupposition failure versus uninterpretability

In chapter 3, we talked about cases of a rather different sort in which a linguistic expression fails to have a semantic value. Recall our treatment of so-called \(\Theta\)-Criterion violations like “Ann laughed Jan”. We observed that this sentence is not in the domain of the \([\ ]\) function as defined by our semantic theory. We called such structures “uninterpretable”, and we proposed that the uninterpretability of “Ann laughed Jan” accounted for the ungrammaticality judgment represented by the asterisk.

In the present chapter, however, we have just suggested that sentences which lack a semantic value are intuitively judged as presupposition failures. So it seems that we have not been consistent. Is it a kind of falsity or a kind of ungrammaticality that we want our theory to capture when it provides no denotation for a given structure? The two are obviously quite different intuitively, and by simply conflating them we would be missing a systematic fact about people’s linguistic intuitions. We might try to draw the intended distinction as follows:

\[(12) \quad \text{If } \alpha \text{ is uninterpretable, then it can be proved from the semantics alone that } \alpha \text{ is outside the domain of } [\ ].\]

\[(13) \quad \text{If it is a contingent matter of fact that } \alpha \text{ is outside the domain of } [\ ], \text{ then } \alpha \text{ is a presupposition failure.}\]

\((12)\) and \((13)\) correctly distinguish between, say, “Ann laughed Jan” on the one hand and “The escalator in South College is moving” on the other. In the former
case, we need not assume anything about the world to show it lacks a denotation. In the latter case, we need to invoke physical facts to show this, and we can easily imagine counterfactual states of affairs in which that sentence would have a truth-value.

But if we tried to turn (12) and (13) into biconditionals that could stand as definitions of uninterpretability and presupposition failure, we would face an objection: namely, that there are sentences which intuitively are “necessary presupposition failures” – for example, “John met the man who died and didn’t die”. By the criteria given in (12), (13), this is indistinguishable from “Ann laughed Jan”: we only need the semantics of English in order to infer that it has no semantic value. But its intuitive status is different, and it should be classified as a presupposition failure rather than as uninterpretable.

So the distinction we are after cannot simply be identified with the difference between necessary and contingent lack of denotation. If we want to characterize it in precise terms, we have to be more specific. In the case of an uninterpretable structure, information about the type of each subtree is sufficient to decide that the structure receives no denotation. To detect presupposition failure, by contrast, we must know more about the denotations of certain subtrees than their mere semantic types.

4.5 Modifiers in definite descriptions

We conclude this chapter by highlighting some predictions that our current semantics makes about the interaction of modifiers and the definite article when these co-occur.

As was already implicit in our treatment of examples in section 4.4, we are assuming a syntax according to which restrictive modifiers within DPs (determiner phrases) form a constituent with the head noun to the exclusion of the determiner. That is, the bracketing is as in (1), not as in (2).

(1) 
```
DP 
  the 
  NP 
    book 
    PP 
      on the pillow 
```

(2) 
```
DP 
  DP 
    PP 
      the 
      book 
    on the pillow 
```
Accordingly, our semantics composes the values of the noun and the modifier before composing the result of this with the value of the determiner.

Syntacticians have frequently entertained the opposite hierarchical organization (2). Suppose this were the structure of the object in the sentence “Ann dislikes the book on the pillow”. What would happen if we attempted to interpret this sentence by our principles of semantic composition?

The first answer to this question is that the alternative structure in (2) leads to uninterpretability within our framework – not within (2) itself, but in whatever larger structure (2) is part of. What happens is that the lower DP, “the book”, denotes the unique (contextually relevant) book (if any), and the higher DP then denotes a truth-value: namely, 1 if the unique book is on the pillow, and 0 if it is not. So any attempt to interpret the next higher node (say, the S-node if (2) is a subject, or the VP-node if it is the object of a transitive verb) will fail due to type-mismatch.19

In a historic debate, Barbara Partee20 took the uninterpretability of structures like (2) (given certain plausible assumptions) to be an argument against syntactic analyses that imply the bracketing [[determiner noun]modifier] and in favor of those that assume [determiner [noun modifier]]. Noam Chomsky21 pointed out in response that there is undeniable evidence for such alleged “uninterpretable” surface bracketings in at least some constructions in some languages. (Directly relevant here, for instance, are Scandinavian languages, where the definite article is realized as a suffix on the noun and intervenes in the linear surface order between it and the PP.) The proper conclusion to draw seems to be that surface structure need not always be the appropriate input to semantic interpretation. Rather, the level that is interpreted is a more abstract one. Our hypothesis then must be that even languages that exhibit the bracketing [[determiner noun] modifier] on the surface have [determiner [noun modifier]] at some other level of representation.

Notes

1 For a more systematic treatment of the indefinite article, see ch. 6.
3 Ibid., chs 4 and 7.
4 A more explicit formulation along the lines of section 3.3 would be (4’).

(4’)

If α is a branching node and {β, γ} is the set of α’s daughters, then α is in the domain of [[ ]] if both β and γ are, and [[β]] and [[γ]] are both in $D_{<\alpha,t>}$.

In this case, $[\alpha] = \lambda x \in D_{\alpha} . [[\beta]](x) = [[\gamma]](x) = 1$.

5 Remember that we couldn’t have written “[[α]] = λ x ∈ D_{α} . [[β]](x) and [[γ]](x)” in the last line of (6). [[β]](x) and [[γ]](x) are truth-values. A value description of the
form \[ \lbrack \bar{p} \rbrack(x) \text{ and } \lbrack \bar{y} \rbrack(x) \], then, is as ill-formed as a value description of the form "1 and 0".


11 And we won't even consider the full range of occurrences of this. Since all the common nouns we use in the book are singular count nouns, we say nothing about definite mass and plural terms. We confine our attention to what philosophers call "definite descriptions": viz. singular terms in which "the" combines with an NP headed by a count noun.


14 If we help ourselves to a little bit of mathematical and logical notation, we can make it shorter:
\[ \lambda f : f \in D_{\text{ext}} \quad \& \exists! x[f(x) = 1] \quad \& \forall y[f(y) = 1]. \]

The abbreviatory conventions used here are the following:

"\( \exists! x[\phi] \)" abbreviates "there is exactly one \( x \) such that \( \phi \)."

"\( \iota x[\phi] \)" abbreviates "the unique \( x \) such that \( \phi \)."

(The first symbol in the last line is the Greek letter iota.)

15 Most notably B. Russell, in “On Denoting,” *Mind*, 14 (1905), pp. 479–93, who proposed an alternative analysis which predicts that \( [[\text{John is on the escalator in South College}]] = 0 \) (given that there is no escalator in South College).


18 The issue did not really arise before section 4.4, since in section 4.3 we only used the determiner “a”. Since that was taken to be vacuous, nothing depended on its exact place in the phrase structure hierarchy.

19 We predict, thus, that if structures like (2) are generated by the syntax of English at all, they can only be used unembedded to make statements, or embedded in positions where truth-values are selected (e.g., as complements to truth-functional connectives). This possibility seems not to be realized, a fact for which we must seek some syntactic explanation.


In this chapter, we consider the internal semantics of another kind of NP-modifier, the relative clause. This will give us occasion to introduce variables and variable binding. The standard interpretation techniques for structures with variables work with variable assignments. Since the notion of a "variable assignment" is not an easy one, we will introduce it in a stepwise fashion. We will start out with a simplified notion of a variable assignment that allows us to interpret structures with (possibly multiple occurrences) of just one variable. Once the essential methods of variable interpretation and variable binding are in place, we will introduce the general notion of a variable assignment and look at structures with multiple variables. The final two sections of the chapter are devoted to general issues of variable binding in syntax and semantics.

5.1 Relative clauses as predicates

Our analysis of relative clauses goes back at least to Quine:¹

The use of the word "relative" in "relative clause" has little to do with its use in "relative term".² A relative clause is usually an absolute term. It has the form of a sentence except that a relative pronoun stands in it where a singular term³ would be needed to make a sentence, and often the word order is switched; thus "which I bought". A general term of this sort is true of just those things which, if named in the place of the relative pronoun, would yield a true sentence; thus "which I bought" is true of just those things x such that x I bought, or, better, such that I bought x.

From this last broad rule we see in particular that a relative pronoun is in a way redundant when it occurs as subject. For example, "who loves Mabel" is true of just the persons of whom "loves Mabel" is true, and "which is bigger than Roxbury" is true of just the things of which "bigger than Roxbury" is true. But the redundant pronoun can serve a grammatical
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purpos e: we switch from “loves Mabel” to “who loves Mabel” for attribu­
tive use as in “brother who loves Mabel”, just because relative clauses
are adjectival⁴ and hence suited, unlike the verbal form “loves Mabel”, to
attributive position. There is less purpose in “which is bigger than Roxbury”,
since “bigger than Roxbury” is adjectival already. The main use of a form
like “which is bigger than Roxbury” is after a comma as an unrestrictive
clause; and we may pass over unrestrictive clauses, for they are only
stylistic variants of coordinate sentences.

At any rate the peculiar genius of the relative clause is that it creates
from a sentence “... x ...” a complex adjective summing up what that
sentence says about x. Sometimes the same effect could be got by dropping
“x is”, as in the last example, or by other expedients; thus, in the case of
“I bought x”, “bought by me” (formed by conversion and application⁵) would serve as well as the relative clause “which I bought”. But often, as
in the case of “the bell tolls for x”, the relative clause is the most concise
adjective available for the purpose.

... A fruitful basis for singular descriptions is the general term of the
form of a relative clause; thence “the car [which] I bought from you”. Let
us build this example from its elements. We have a triadic relative term
“bought from”, which, applied predicatively to the singular terms “I”, “x”
(say), and “you”, gives a sentence form “I bought x from you”. Putting
the relative pronoun for the “x” here and permuting, we get the relative
clause “which I bought from you”. This clause is a general term, adjectival
in status. Combining it attributively with the general term “car”, we get
the general term “car which I bought from you”; and then “the” yields the
singular term.

As Quine makes clear, if we abstract away from their internal syntactic and
semantic composition, relative clauses are just like other modifiers in NP — for
example, the PPs and APs we considered earlier. They have the same type of
denotation (namely, characteristic functions of sets), and they contribute in the
same way to the denotation of the surrounding structure. The latter actually
follows from the former in our framework. Since we don’t allow construction­
specific semantic rules but only general principles of composition, we are
committed to the prediction that phrases in the same environment and with the
same type of denotation must make the same contribution. Let’s consider an
example:

(1) The house which is empty is available.

Omitting the internal structure of the relative clause, the determiner phrase (DP)
in (1) has a structure of the following kind:
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House", "empty", and "available" have the obvious lexical entries; for example, \[\text{empty} \equiv \lambda x \in D_e . x \text{ is empty}\]. We also have a denotation for the determiner "the". So once we decide on an interpretation for the complementizer phrase (CP) "which is empty", we can calculate the truth-conditions of (1').

Following Quine, we hypothesize that "which is empty" has exactly the same denotation as "empty". This is suggested by the fact that substituting "empty" for "which is empty" leads to a sentence with equivalent truth-conditions: "The empty house is available". (Never mind the change in order, which is for some reason required by the syntax of English.) The rest now follows: [house] and [which is empty] combine by Predicate Modification (PM), with the result that the NP "house which is empty" denotes the function \(\lambda x \in D_e . x \text{ is a house and } x \text{ is empty}\). When we apply Functional Application (FA) to compute the denotation of the DP above, this function becomes the argument of [the]. According to the entry for "the", then, the DP "the house which is empty" has a denotation iff there is exactly one empty house, and in that case it denotes the unique empty house. Finally, the whole S (again by FA) has a truth-value iff there is a unique empty house, and it has the value 1 iff the unique empty house is available. This prediction conforms to intuitive judgment.

Notice that this analysis correctly distinguishes restrictive relatives from their nonrestrictive counterparts: for example, (1) from (2).

(2) The house, which is empty, is available.

If we assume (as we did above and as does Quine) that nonrestrictive modifiers are like separate sentences, then [the] in (2) applies to the extension of "house" by itself (and not to the extension of any constituent including "which is empty"). This implies that (2) presupposes there to be exactly one house – unlike (1), which is entirely compatible with there being two or more houses, as long as only one of them is empty.\(^6\) In what follows, we will disregard nonrestrictives and talk exclusively about restrictive relatives.

To sum up, restrictive relatives are just another kind of intersective modifier.
5.2 Semantic composition inside the relative clause

What remains to be worked out, then, is the internal semantic composition of relative clauses. What are their internal syntactic structures, and what lexical entries and composition principles are needed to make them denote the appropriate functions of type <e,t>? For instance, how do we derive systematically what we simply assumed above: namely, that \([\text{which is empty}] = [\text{empty}]\)?

We will adopt a more or less standard syntactic analysis of relative clauses, according to which they look like this:

(1) \[
\begin{array}{c}
\text{CP} \\
\text{\quad which} \\
\text{\quad } \\
\text{\quad C} \\
\text{\quad } \\
\text{\quad S} \\
\text{\quad that} \\
\text{\quad DP} \\
\text{\quad VP} \\
\text{\quad t is empty}
\end{array}
\]

Various other structures would serve our purposes equally well, as long as there is a relative pronoun at the top and a trace in subject position. In (1), either the complementizer (C) or the *wh*-word has to be deleted on the surface. We categorized the trace in (1) as a DP (determiner phrase). From now on, we will take all phrases that show the same syntactic behavior as phrases headed by overt determiners to be DPs. This includes proper names, pronouns, and traces.

The semantics for structures like (1) will require some innovative additions to our current theory. What we know at this point is what we want to come out on top: The CP (complementizer phrase) should get a value in \(D_{e,t}\); more particularly, in the case of (1), it should denote the characteristic function of the set of empty objects. We also know what the VP inside this CP denotes; it's actually that very same function. But it would be too easy to assume therefore that the semantic value simply gets passed up from VP to S to C to CP. That would work in this special case, but not (as Quine already remarked) in general. We must also worry about cases where the trace is not in subject position:
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(2) should denote the function \( \lambda x \in D \). John abandoned \( x \). In this case, this is not the value of any of its subtrees.

The basic question we face at this point is: What are the semantic values of traces? For instance, what is the extension of the object DP in (2)?

We face a dilemma here. On the one hand, we would like traces to have the same type of extension as other DPs, because then we could use the same composition principles as before to interpret the structures they appear in. For instance, we would like “\( t \)” in (2) to denote an individual, because that would make a suitable argument for the extension of the verb “abandon”. If we treated the trace as semantically vacuous, for instance, we would be in trouble: the S-node would denote the characteristic function of the set of all individuals who abandoned John – with John being the one abandoned instead of the aban doner!

On the other hand, does it make sense to assign a referent to the trace? Let’s investigate this possibility.

### 5.2.1 Does the trace pick up a referent?

It is sometimes suggested that the relative pronoun is anaphorically related to and inherits the reference of the relative clause’s head. For instance, in “the movie (which) Mary saw \( t \)”, the trace is said to get its denotation from “which”, which in turn gets it from the head. One immediate objection to this idea is that it would not apply to relatives in quantifier phrases like “no movie that Mary saw”, since such phrases do not denote individuals (as we will argue in more detail below). But even if we stick to definite DPs, it is an idea that raises puzzling questions.

First, what do we mean by the “head” of the relative clause? In the syntactic structure we are assuming, this could only be the whole containing DP:
(3)

There just is no smaller DP in this structure. In fact, there is no smaller constituent of any kind that denotes an individual and from which a referent of the desired type e could thus be picked up by the trace. The whole DP does denote an individual. But if we said that the trace inherits its denotation from that, we would get into a vicious circle: The denotation of the whole is supposed to be built up from the denotations of its parts, and this construction can’t get off the ground if we need the denotation of the whole before we can interpret some of the parts.

Should we reconsider the structure, then, and adopt the following bracketing after all?

(4)

No, that just gets us into different problems. Now [the] applies to [movie], giving rise to a presupposition that there is just one (relevant) movie. Besides, it is mysterious what the CP would have to denote. It couldn’t be a function of type <e,t>, because then the whole DP would denote (if anything) a truth-value instead of an individual.
Let us abandon this line of approach. Here, then, is the dilemma: We would like the trace to denote an individual so that we can interpret the nodes above it, but we can't seem to find a suitable individual. There is no easy way out within the confines of our current theoretical apparatus. It is time to explore the utility of a genuinely new theoretical construct, the variable.

5.2.2 Variables

Variables were invented precisely to be like ordinary referring phrases in the respects we want them to be, but sufficiently unlike them to avoid the puzzles we just ran up against. A variable denotes an individual, but only relative to a choice of an assignment of a value. What is a value assignment for a variable? The simplest definition for our present purposes is this:

(5) Preliminary definition: An assignment is an individual (that is, an element of $D (= D_c)$).

A trace under a given assignment denotes the individual that constitutes that assignment; for example:

(6) The denotation of “t” under the assignment Texas is Texas.

An appropriate notation to abbreviate such statements needs to be a little more elaborate than the simple $[\ldots]$ brackets we have used up to now. We will indicate the assignment as a superscript on the brackets; for instance, (7) will abbreviate (6):

(7) $[t]_{\text{Texas}} = \text{Texas}$.

The general convention for reading this notation is as follows: Read “$[\alpha]^a$” as “the denotation of $\alpha$ under $a$” (where $\alpha$ is a tree and $a$ is an assignment).

(7) exemplifies a special case of a general rule for the interpretation of traces, which we can formulate as follows:

(8) If $\alpha$ is a trace, then, for any assignment $a$, $[\alpha]^a = a$.

The decision to relativize the denotations of traces to assignments has repercussions throughout our system of rules. We must allow the denotations of larger phrases that contain traces to be assignment-relative as well. For instance, a VP whose object is a trace will not denote a fixed function in $D_{\text{ext}}$, but may denote different functions under different assignment functions; for instance:
A sentence like John abandoned t, then, does not have truth-conditions per se, but only with respect to an assignment. We have, for example:
Now that we have admitted traces into our syntactic representations, we find ourselves in a rather odd situation. We have sentences like "John abandoned Mary" that have truth-conditions (and hence a meaning) per se, and our theory should acknowledge this fact, as it has always done. But we also have to deal with sentences like "John abandoned t" that may appear as parts of relative clauses, and need to have assignment-dependent denotations. How can we do justice to both types of sentences without complicating our composition principles? Take the Functional Application principle, for instance. It should be written in such a way that the top node gets an assignment-dependent value whenever either of the daughters does. Now it would be rather inelegant to have to distinguish three different cases here, according to whether the function-denoting daughter, or the argument-daughter, or neither of them happens to contain a trace and therefore to have an assignment-relative value. It is simpler, if a bit artificial, to formulate all our composition principles for assignment-dependent denotations and introduce assignment-independent denotations through a definition, as follows:

(9) For any tree \( \alpha \), \( \alpha \) is in the domain of \([ \ ]\) iff for all assignments \( a \) and \( b \),
\[
[\alpha] = [\alpha].
\]
If \( \alpha \) is in the domain of \([ \ ]\), then for all assignments \( a \),
\[
[a] = [\alpha].
\]

As for lexical rules, if we have an item like laugh, we can still assign it the assignment-independent denotation \([\text{laugh}]\). But with definition (9), we automatically also have a semantic value for this item under any arbitrary assignment. From (9) and the lexical entry for laugh, it follows that

(10) For any assignment \( a \),
\[
[\text{laugh}] = [\text{laugh}] = \lambda x \in D_e . x \text{ laughs}.
\]

Now we can write the new versions of our composition principles almost as simply as the old ones, except that we need to distinguish two cases in the interpretation of terminal nodes. Our previous Terminal Nodes rule now divides into (8) above and (11).
(11) **Lexical Terminals**
If α is a terminal node occupied by a lexical item, then [α] is specified in the lexicon.

(12) **Non-Branching Nodes (NN)**
If α is a non-branching node and β its daughter, then, for any assignment a, [α]ₐ = [β]ₐ.

(13) **Functional Application (FA)**
If α is a branching node and {β, γ} the set of its daughters, then, for any assignment a, if [β]ₐ is a function whose domain contains [γ]ₐ, then [α]ₐ = [β]ₐ([γ]ₐ).

(14) **Predicate Modification (PM)**
If α is a branching node and {β, γ} the set of its daughters, then, for any assignment a, if [β]ₐ and [γ]ₐ are both functions of type <e,t>, then [α]ₐ = λx ∈ D . [β]ₐ(x) = [γ]ₐ(x) = 1.

---

**Exercise**

Consider a state of affairs with the following properties:

D (= Dₑ) = {b, j, m, s}.

s invites b and m; b invites b, j, and m; no other invitations take place.

[Bob] = b.

(a) Show that [[(i)]]ₐ = 1.

(i)  
```
S
 /   
|    |
|    |
|    |
|    |
```

(b) Which assignments make (i) true? (List all.)

(c) Show that [[(i)]] is undefined: that is, that the tree under (i) is not in the domain of [[ ]].
5.2.3 Predicate abstraction

Finally, we can start to think about the semantic principle that determines the denotation of a relative clause like (2) (repeated from above) from its two immediate constituents: the relative pronoun which and the sentence John abandoned t.

(2) CP
   \[\text{which}\]
   \[C\]
   \[C\]
   \[that\]
   \[DP\]
   \[V\]
   \[DP\]
   \[abandoned\]
   \[t\]

We treat the complementizer "that" as semantically vacuous, so the \(\bar{C}\) inherits the value of the S below it. The relative pronoun within CP is also not assigned any denotation of its own. But it is not simply vacuous; its presence will be required to meet the structural description of the composition principle applying to the CP above it. That principle is a new one, unrelated to the ones we have employed thus far:

12 (15) Predicate Abstraction (PA)
   If \(\alpha\) is a branching node whose daughters are a relative pronoun and \(\beta\), then \(\llbracket \alpha \rrbracket = \lambda x \in D . \llbracket \beta \rrbracket^x\).

This rule is also known as "Functional Abstraction" or "Lambda Abstraction" in the literature. The reason for these names is transparent. The semantic value of \(\alpha\) is defined as a function. In fact, our formulation of PA uses the \(\lambda\)-notation to define this function. (This was not, of course, essential; we could have defined it in words as well.) To appreciate what the rule is doing, look again at (2). Being a relative clause, (2) should have an assignment-independent interpretation, and it should denote the function \(\lambda x \in D . \text{John abandoned x}\). Here is the proof that it does:
Relative Clauses, Variables, Variable Binding

CP

[which]

C

S

John

abandoned

= (by PA)

λx ∈ D.

[ C S ]

X

John

abandoned

= (by vacuity of C)

λx ∈ D.

[X]

X

John

abandoned

= (by FA)

λx ∈ D.

[ VP ]

X

abandoned

( [John]) = (by definition (9))

λx ∈ D.

[ VP ]

X

abandoned

( [John]) = (by lexical entry of John)

λx ∈ D.

[ VP ]

X

abandoned

(John) = (by FA)

λx ∈ D. [abandoned]x( [t]x) (John) = (by Traces Rule)

λx ∈ D. [abandoned]x(x)(John) = (by definition (9))
\(\lambda x \in D . \text{[abandoned]}(x)(\text{John}) = (\text{by lexical entry of abandoned})\)

\(\lambda x \in D . [\lambda y \in D . [\lambda z \in D . z \text{ abandoned } y]](x)(\text{John})\)

= (by definition of \(\lambda\)-notation)

\(\lambda x \in D . \text{John abandoned } x. \text{ QED.}\)

The Predicate Abstraction Rule gives the moved relative pronoun what is called a *syncategorematic* treatment. Syncategorematic items don’t have semantic values of their own, but their presence affects the calculation of the semantic value for the next higher constituent. A syncategorematic treatment of relative pronouns goes against the concept of type-driven interpretation that we argued for earlier, and we will eventually want to abolish rules of this kind. At the moment, we will not worry about this blemish of our theory, however, and will keep the current version of the Predicate Abstraction Rule (at least for a little while), since it is easier to handle for beginning semanticists than the theoretically more adequate alternatives.

When you work with assignments, do not confuse denotations *under* assignments with denotations *applied to* assignments. There is a big difference between \([a]^x\) and \([a](x)\). For example,

\[
\begin{align*}
[\text{whom John abandoned } t]\text{Charles} &\neq [\text{whom John abandoned } t](\text{Charles}) \\
[\text{sleeps}]\text{Ann} &\neq [\text{sleeps}](\text{Ann})
\end{align*}
\]

What you see on the left side is a function of type \(<e,t>\), but what you have on the right is the result of applying such a function to an individual — in other words, a truth-value. In many other instances, one of the two notations isn’t even well-defined in the first place. For example,

“\([\text{John abandoned } t]^x\)”

makes sense: it stands for a truth-value; that is, it equals either 1 or 0, depending on what individual “\(x\)” stands for.

“\([\text{John abandoned } t](x)\)”

on the other hand is nonsense, for two reasons. First, it falsely presumes that “John abandoned \(t\)” is in the domain of \([\_]\); that is, that it has a semantic value independently of a specified assignment. This is not the case. Second, even if “John abandoned \(t\)” did happen to have an assignment-independent denotation, its denotation (like that of any other \(S\)) would be a truth-value, not a function. So it would be as though we had written “1\((x)\)” or “0\((x)\)”._
5.2.4 A note on proof strategy: bottom up or top down?

When you are asked to calculate the semantic value of a given tree, you have a choice between two strategies: to work from the bottom up or from the top down. To take a very simple example, you are told that, as a matter of actual fact, Ann doesn’t like Al, and you are asked to show that

$$S$$

$$Ann \quad VP$$

$$likes \quad Al$$

= 0

One way to present your proof is as follows:

**Bottom-up proof**

By lexicon: $[[Ann]] = Ann$ (=: (i))

By lexicon: $[[\text{likes}]] = \lambda x \in D . \lambda y \in D . y \text{ likes } x$ (=: (ii))

By lexicon: $[[Al]] = Al$ (=: (iii))

By FA: $[[likes]]([[Al]])$

Therefore, using (ii) and (iii) from above:

$$VP$$

$$likes \quad Al$$

$= [\lambda x \in D . \lambda y \in D . y \text{ likes } x](Al) = \lambda y \in D . y \text{ likes } Al$ (=: (iv))

$$S$$

$$Ann \quad VP$$

$$likes \quad Al$$

$= [[[Ann]]]([[likes]]([[Al]]))$
Therefore, using (i) and (iv) from above:

\[
\begin{aligned}
S \\
\text{Ann} & \quad \text{VP} \\
\text{likes} & \quad \text{AI}
\end{aligned}
\]

= [\lambda y \in D . y \text{ likes AI}](\text{Ann})

This means:

\[
\begin{aligned}
S \\
\text{Ann} & \quad \text{VP} \\
\text{likes} & \quad \text{AI}
\end{aligned}
\]

= 1 iff Ann likes AI.

And given the description of the facts, Ann doesn’t like AI, therefore

\[
\begin{aligned}
S \\
\text{Ann} & \quad \text{VP} \\
\text{likes} & \quad \text{AI}
\end{aligned}
\]

= 0.

QED.

This was the bottom-up strategy, because we calculated values for the lowest nodes first and got to the top node last. Alternatively, you could have presented your reasoning as follows:

Top-down proof

\[
\begin{aligned}
S \\
\text{Ann} & \quad \text{VP} \\
\text{likes} & \quad \text{AI}
\end{aligned}
\]

= 1

iff (by FA)

\[
\begin{aligned}
\text{VP} \\
\text{likes} & \quad \text{AI}
\end{aligned}
\] (\([\text{Ann}]\)) = 1
iff (by FA)

\[[\text{likes}]([\text{Al}])([\text{Ann}]) = 1\]

iff (by lexical entries)

\[\lambda x \in D . \lambda y \in D . y \text{ likes } x](\text{Al})(\text{Ann}) = 1\]

iff (by definition of \(\lambda\)-notation)

Ann likes Al.

Since she doesn’t, therefore,

\[
\begin{array}{c}
\text{S} \\
\text{Ann} \\
\text{VP} \\
\text{likes} \\
\text{Al}
\end{array}
\]

= 0.

QED.

This second strategy was the top-down strategy. Both strategies are equally sound, and there is no major difference in efficiency. Most people are naturally inclined towards the bottom-up strategy at first, but there is really no reason to prefer it.

As we turn to calculations involving Predicate Abstraction, however, the top-down method begins to have an advantage. Suppose your assignment is to show that the following equation holds:

\[
\begin{array}{c}
\text{CP} \\
\text{who} \\
\text{C} \\
\text{S} \\
\text{NP} \\
\text{t} \\
\text{sleeps}
\end{array}
\]

= \lambda x \in D . x \text{ sleeps.}

Let’s first prove this by the top-down method.
Top-down proof

\[\begin{array}{c}
\text{CP} \\
\text{who} \\
\text{C} \\
\text{C} \\
\text{S} \\
\text{t sleeps}
\end{array}\]

\[\begin{array}{c}
\text{Ann} \\
\text{= (by PA and definition (9))}
\end{array}\]

\[\begin{array}{c}
\lambda x \in D . \\
\text{S} \\
\text{t sleeps}
\end{array}\]

\[\begin{array}{c}
\lambda x \in D . [\text{[sleep]}^x(\text{[[t]^x]])] = (by \text{Traces Rule})
\end{array}\]

\[\begin{array}{c}
\lambda x \in D . [\text{[sleep]}^x(\text{x})] = (by \text{lexical entry of sleep and definition (9)})
\end{array}\]

\[\begin{array}{c}
\lambda x \in D . [[[\lambda y \in D . y \text{sleeps}](\text{x})] = (by \text{definition of } \lambda\text{-notation})
\end{array}\]

\[\begin{array}{c}
\lambda x \in D . \text{x sleeps.}
\end{array}\]

QED.

This was straightforward. Now suppose we had tried to do it bottom up. Here is the right way to do that:

Bottom-up proof

Let \( x \in D \) be an arbitrarily chosen assignment.

By the rule for traces: \([\text{t}]^x = x \) (i)

\([\text{sleeps}]^x = \lambda y \in D . y \text{ sleeps.} \)](ii)

By FA:

\[\begin{array}{c}
\text{S} \\
\text{t sleeps}
\end{array}\]

\[\begin{array}{c}
= [[\text{sleeps}]^x(\text{[[t]^x})]
\end{array}\]
herefore, using (i) and (ii) from above:

\[
S \quad \frac{t \text{ sleeps}}{} \quad ]^x = [\lambda y \in D . y \text{ sleeps}](x).
\]

\[
\begin{array}{c}
C \\
S \\
t \text{ sleeps}
\end{array} \quad ]^x = [\lambda y \in D . y \text{ sleeps}](x)
\]

ince \( x \) was arbitrary, we can now summarize:

\[
\begin{array}{c}
C \\
S \\
t \text{ sleeps}
\end{array} \quad ]^x = [\lambda y \in D . y \text{ sleeps}](x) [=: (iii)]
\]

y PA and definition (9):

\[
\begin{array}{c}
CP \\
\frac{\text{who}}{} \\
C \\
S \\
t \text{ sleeps}
\end{array} \quad ]^x = \lambda x \in D .
\]

\[
\begin{array}{c}
C \\
S \\
t \text{ sleeps}
\end{array} \quad ]^x
\]

y (iii) above, this means:

\[
\begin{array}{c}
CP \\
\frac{\text{who}}{} \\
C \\
S \\
t \text{ sleeps}
\end{array} \quad ]^x = \lambda x \in D . [ [\lambda y \in D . y \text{ sleeps}](x) ]
\]
By the definition of $\lambda$-notation, this is equivalent to:

$$\begin{array}{c}
\text{CP} \\
\text{who} \\
\text{C} \\
\text{C} \\
\text{S} \\
\text{t} \\
\text{sleeps} \end{array} \quad \begin{array}{c}
\text{Ann} \\
= \lambda x \in D \cdot x \text{ sleeps.} \end{array}$$

QED.

This is just as good a proof as the previous one. However, it *required a certain amount of foresight to write up*. We had to anticipate right at the beginning that we would need to work out the semantic values of the subtrees up to $\overline{C}$ for an arbitrary assignment. Had we just plunged in without thinking ahead, we would have been tempted to calculate something which subsequently turns out to be irrelevant: namely, the extension of $\overline{C}$ with respect to the particular assignment Ann. In that case, we would have wasted the first few steps of the proof to determine that

$$\begin{array}{c}
\overline{C} \\
\text{C} \\
\text{S} \\
\text{t} \\
\text{sleeps} \end{array} \quad \begin{array}{c}
\text{Ann} \\
= [\lambda y \in D \cdot y \text{ sleeps}](\text{Ann}). \ [=: (i)] \end{array}$$

This is perfectly correct, but it is not helpful for the continuation of the proof, because for the next node up, we derive:

$$\begin{array}{c}
\text{CP} \\
\text{who} \\
\text{C} \\
\text{C} \\
\text{S} \\
\text{t} \\
\text{sleeps} \end{array} \quad \begin{array}{c}
\text{Ann} \\
= \lambda x \in D \cdot x \text{ sleeps.} \end{array}$$

Now the fact established in (i) gives us only a tiny piece of partial information about which function it is that we have at the right-hand side of this equation:
(i) merely tells us that it is some function or other which to Ann assigns 1 iff she sleeps. It doesn’t tell us how it behaves with arguments other than Ann. So that’s not enough information to finish the proof. We are supposed to prove the equality of two functions, and for this it does not suffice to make sure that they assign the same value to Ann. Rather, we must show that they agree on all their arguments.

The moral of this is the following: If you already see clearly where you are headed, the bottom-up strategy can be as efficient and elegant as the top-down strategy. But if you are still groping in the dark or don’t want to take any chances, top-down is the way to go. The advantage is that by the time you get to the lower nodes, you will already know which assignments it is relevant to consider.

---

**Exercise**

Some interesting issues arise when we bring together our new analysis of relative clauses and the previous chapter’s Fregean treatment of definite descriptions. Suppose we embed a non-denoting definite in a relative clause:

(i) John is a man who attended the 1993 Olympics.

The intuitive status of (i) is the same as that of (ii).

(ii) John attended the 1993 Olympics.

Both are perceived as presupposition failures, since there were no Olympic Games in 1993. Our semantics from chapter 4 predicts that (ii) receives no truth-value. The same should come out for (i). Check whether it does by using the “pedantic” versions of the PA rule and the other composition principles, repeated here for convenience:

(12') **Non-Branching Nodes (NN)**
If α is a non-branching node and β its daughter, then, for any assignment a, α is in the domain of [ ]^a if β is. In this case, [α]^a = [β]^a.

(13') **Functional Application (FA)**
If α is a branching node and {β, γ} the set of its daughters, then, for any assignment a, α is in the domain of [ ]^a if both β and γ are, and [β]^a is a function whose domain contains [γ]^a. In this case, [α]^a = [β]^a([γ]^a).

(14') **Predicate Modification (PM)**
If α is a branching node and {β, γ} the set of its daughters, then, for any assignment a, α is in the domain of [ ]^a if both β and γ are, and
\[ [\beta]^a \text{ and } [\gamma]^a \text{ are both of type } <e,t>. \text{ In this case, } [\alpha]^a = \lambda x : x \in D \text{ and } x \text{ is in the domain of } [\beta]^a \text{ and } [\gamma]^a. \text{ } [\beta]^a(x) = [\gamma]^a(x) = 1. \]

(15') **Predicate Abstraction (PA)**

If \( \alpha \) is a branching node whose daughters are a relative pronoun and \( \beta \), then \( [\alpha] = \lambda x : x \in D \text{ and } \beta \text{ is in the domain of } [\ ]^a \cdot [\beta]^a \).

Consider next a case where the definite in the relative clause contains the variable. An example is (iii).

(iii) **John is a man whose wife is famous.**

This involves so-called "pied-piping", the fronting of a larger constituent than the mere relative pronoun. Following many previous authors, we assume that pied-piping is essentially a surface phenomenon, and that the input structure for semantic interpretation is as if only the relative pronoun had moved:13

(iv) **John is a man who [t's wife is famous].**

We also assume that possessive constructions like "John's wife", "his wife", and in this case, "t's wife", are definites, headed by a non-overt equivalent of "the". Glossing over the details, we assume that the syntactic representation of "John's wife" (at the relevant level) is essentially "[the [wife (of) John]]". This can be straightforwardly interpreted by our semantic rules (assuming that "wife" has a meaning of type \(<e,<e,t>\))

The prediction – adequate, it seems – is that "John's wife" denotes John's unique wife if he has a unique wife, and nothing otherwise. So a sentence like "John's wife is famous" is truth-value-less (a presupposition failure) if John is not married to any woman, or to several.

Back now to (iv). The analysis of "t's wife" is just like that of "John's wife", except, of course, that this time the denotation depends on the assignment. Assuming the "pedantic" versions of the composition principles, compute the denotations of the essential constituents of (iv). What happens if John is not married? What if nobody (in the salient universe of discourse) is married? At what point does the computation crash in each case?

---

5.3 **Multiple variables**

5.3.1 **Adding "such that" relatives**

Quine's discussion of relative clauses also includes a remark on the "such that" construction:
The reason for permuting word order in forming relative clauses is to bring the relative pronoun out to the beginning or near it. The task can be exacting in complex cases, and is sometimes avoided by recourse to an alternative construction, the unlyrical "such that". This construction demands none of the tricks of word order demanded by "which", because it divides the two responsibilities of "which": the responsibility of standing in a singular-term position within the clause is delegated to "it", and the responsibility of signaling the beginning of the clause is discharged by "such that". Thus "which I bought" becomes "such that I bought it"; "for whom the bell tolls" becomes "such that the bell tolls for him".

The "such that" construction is thus more flexible than the "which" construction.14

So the "such that" construction has a different syntax, but essentially the same semantics as ordinary relatives. It should be a routine matter, therefore, to extend our analysis to cover it. We just need to generalize our Traces Rule to pronouns, and to rewrite PA so that it treats a "such" like a relative pronoun. (The "that" is presumably the semantically vacuous complementizer again.)

(1) **Pronoun Rule** (new addition)
   If \( \alpha \) is a pronoun, then for any assignment \( a \in D = D_e \), \([\alpha]^* = a\).

(2) **Predicate Abstraction** (revised)
   If \( \alpha \) is a branching node and \( \beta \) and \( \gamma \) its daughters, where \( \beta \) is a relative pronoun or \( \beta = "\text{such}" \), then \([\alpha] = \lambda x \in D \cdot [\gamma]^x\). 

It is now easy to calculate a semantic value for a tree like (3),

(3) 
```
   AP
     such
        CP
          that
             S
               Joe
                  VP
                    bought
                      it
```

and to prove that this "such that" phrase has exactly the same denotation as its \( wh \)-counterpart in (4).
5.3.2 A problem with nested relatives

Our reason for introducing “such that” relatives is that it allows us to look at certain examples where one relative clause is nested inside another. Let’s see what interpretation we obtain for the following NP:

(5) man such that Mary reviewed the book he wrote

This phrase has the structure in (5') and receives a meaning of type <e,t>.
We encounter no problems of interpretability here, but the interpretation we derive under our current assumptions is very wrong.

Exercise

Prove that the following predictions are made by our current semantics:

(i) If there isn’t a unique book that wrote itself, then (5) has no denotation under any assignment.

(ii) If Mary reviewed the unique book that wrote itself, then \(((5)) = [\text{man}]\).

(iii) If Mary didn’t review the unique book that wrote itself, then \(((5)) = \lambda x \in D . 0\).

The meaning predicted, then, is not the meaning that (5) actually has in English. If we can express this strange predicted meaning in English at all, we have to use quite a different phrase, namely (6).

(6) man such that Mary reviewed the book which wrote itself

In order to rectify this inadequacy, we will have to reconsider the syntax as well as the semantics of variables.

5.3.3 Amending the syntax: co-indexing

The fact that our example (5) wrongly received the meaning of (6) points us to something that seems to be wrong with our present analysis: namely, that, informally speaking, the semantics pays no attention to which trace/pronoun is related to which wh-word/“such”. Specifically, it apparently treated the “he” in (5) as if it was a trace of the “wh” right below “book”. This suggests that part of the solution might be to employ syntactic representations which explicitly encode which trace or pronoun belongs with which relative ‘pronoun or “such”.

The standard notational device for this purpose is co-indexing by means of numerical subscripts. We will henceforth assume that the syntactic structure for (5) which constitutes the input to the semantic component is a bit richer than (5’) above. It looks rather like (5’’) below.
Of course, there are infinitely many other possible indexings for the same surface phrase. We will return to this point below. For the moment, our goal is to make sure that if (5) is represented as in (5''), then it receives the intuitively correct interpretation. Merely adding co-indexing in the syntax is not, of course, sufficient by itself to attain this goal. As long as our semantic rules fail to “see” the indices, co-indexing will do nothing to determine the appropriate denotation.

5.3.4 Amending the semantics

In order to write an appropriately index-sensitive set of semantic rules, we must first redefine what we mean by an “assignment”. So far, an assignment was just an element of $D (= D_v)$. This worked fine as long as we only considered examples in which each variable was contained in at most one relative clause. In our present, more complicated example (5), we have traces and pronouns that are
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contained in both a smaller and a bigger relative clause. This situation requires
a richer notion of assignment. Our new definition is (7):

(7) A (variable) assignment is a partial function from \( \mathbb{N} \) (the set of natural
numbers) into \( D \).

For example, the following functions are assignments in this new sense:

(8) \[
\begin{array}{c}
1 \rightarrow \text{John} \\
2 \rightarrow \text{Mary}
\end{array}
\]
\[
\begin{array}{c}
1 \rightarrow \text{John} \\
2 \rightarrow \text{John}
\end{array}
\]
\[
\begin{array}{c}
2 \rightarrow \text{John} \\
5 \rightarrow \text{Mary} \\
7 \rightarrow \text{Ann}
\end{array}
\]

Given definition (7), \( \emptyset \) (the empty set) comes out as an assignment too. \( \emptyset \) is the
(only) function with \( \emptyset \) as its domain. Any function from \( \emptyset \) into \( D \), for example,
would have to be a subset of \( \emptyset \times D \), the Cartesian product of \( \emptyset \) and \( D \). Since
\( \emptyset \times D = \emptyset \), the only subset of \( \emptyset \times D \) is \( \emptyset \), hence \( \emptyset \) is the only function from
\( \emptyset \) into \( D \). Since \( \emptyset \subseteq \mathbb{N} \), \( \emptyset \) qualifies as a partial function from \( \mathbb{N} \) into \( D \), hence
as an assignment.

Since assignments are now functions, it makes sense to speak of their
"domain". We write "\( \text{dom}(a) \)" to abbreviate "the domain of the assignment \( a \)". Assignments in the new sense assign potentially different individuals to different
numerical indices. This allows us to replace our old rules for (unindexed) traces
and pronouns by a new rule that is sensitive to the index.\(^{18}\)

(9) \textbf{Traces and Pronouns Rule}

If \( \alpha \) is a pronoun or a trace, \( a \) is a variable assignment, and \( i \in \text{dom}(a) \),
then \[ \alpha_i \] \( a = a(i) \).

(9) makes the following kinds of predictions:

(10) \[
[\text{he}_2]_{[1 \rightarrow \text{Sue}, 2 \rightarrow \text{Joe}]}^{1 \rightarrow \text{Sue}} = \begin{array}{c}
1 \rightarrow \text{Sue} \\
2 \rightarrow \text{Joe}
\end{array}(2) = \text{Joe}
\]
\[
[\text{t}_1]_{[1 \rightarrow \text{Sue}, 2 \rightarrow \text{Joe}]}^{1 \rightarrow \text{Sue}} = \begin{array}{c}
1 \rightarrow \text{Sue} \\
2 \rightarrow \text{Joe}
\end{array}(1) = \text{Sue}
\]

(9) also implies that a given trace or pronoun will not have a well-defined
semantic value under just any assignment. For instance, \( \text{he}_2 \) is not in the domain
of \( [1 \rightarrow \text{Joe}], [2 \rightarrow \text{Joe}], \) or \( [\emptyset] \). As regards \( [\emptyset] \), \textit{no} pronoun or trace is in its
domain. Are any expressions at all? Yes, but only those which are in the domain
of \( [\emptyset] \) for \textit{every} \( a \). This is the case, for example, for the lexical items. Recall our
earlier convention, which we carry over into the new system.\(^{19}\) When an expression
is in the domain of $\exists x$, it is also in the domain of $\exists x'$ for all $a$. In fact, we will now think of "$\exists x$" as simply an abbreviation for "$\exists x^0$".

(11) For any tree $\alpha$, $[\alpha] := [\alpha]^0$

To have a semantic value *simpliciter* means nothing more and nothing less than to have a semantic value under the empty assignment.

As before, assignment dependency is systematically passed up the tree when we construct phrases from lexical items and one or more traces and pronouns. The composition rules NN, FA, and PM can stay just as we formulated them before (though, of course, wherever they refer to "assignments", this now means something different than it used to). With their help, we can now calculate semantic values under given assignments for many larger phrases composed of lexical items and variables. For instance, we can prove the following:

(12) $\begin{array}{c}
\text{he}_2 \\
\text{wrote} \\
t_1 \\
\end{array}$ = 1 iff Joe wrote "Barriers"

**Exercise**

Prove (12).

The Predicate Abstraction rule will have to be revised. Before we can do this, we have to define one further technical concept, that of a so-called modified (variable) assignment:

(13) Let $a$ be an assignment, $i \in \mathbb{N}$, and $x \in D$. Then $a^{x^i}$ (read: "a modified so as to assign $x$ to $i$") is the unique assignment which fulfills the following conditions:

(i) $\text{dom}(a^{x^i}) = \text{dom}(a) \cup \{i\}$,

(ii) $a^{x^i}(i) = x$, and

(iii) for every $j \in \text{dom}(a^{x^i})$ such that $j \neq i$: $a^{x^i}(j) = a(j)$.

(13) defines a possibly new assignment $a^{x^i}$ on the basis of a given assignment $a$. Clause (i) of (13) states that the domain of the new assignment contains the number $i$. If $\text{dom}(a)$ contains $i$ already, then, $\text{dom}(a) = \text{dom}(a^{x^i})$. Otherwise, the
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index i has to be added. Clause (ii) states that the new assignment $a^x_i$ is a function that maps the number i to the individual x. If $\text{dom}(a)$ contains i already, then $a^x_i$ might differ from a with respect to the individual assigned to i. Suppose $a(5) = \text{Mary}$, for example. Then $a^{\text{Ann}/5}(5) = \text{Ann}$, hence $a(5) \neq a^{\text{Ann}/5}(5)$. Clause (iii) makes sure that the new assignment $a^x_i$ is indeed like the original assignment a, except for a possible difference regarding i.

Exercise

Determine $a^{\text{Mary}/2}$ for various concrete choices of a (for example, the assignments mentioned in (8) above).

Sample answers:

$$\begin{align*}
1 \rightarrow \text{John}^{\text{Mary}/2} \quad & = \\
2 \rightarrow \text{Mary} & = \\
1 \rightarrow \text{John}^{\text{Mary}/2} & = \\
2 \rightarrow \text{John} & = \\
1 \rightarrow \text{John}^{\text{Mary}/2} & = \\
2 \rightarrow \text{Mary} & = \\
\end{align*}$$

As you can see from these examples:

The domain of $a^{\text{Mary}/2}$ always includes the number 2;
The domain of $a^{\text{Mary}/2}$ is either the same as the domain of a or larger by one element, depending upon whether a already happened to have 2 in its domain;
a^{\text{Mary}/2} is either identical to a or differs from it in one argument, depending upon whether a already happened to map 2 to Mary.

Modified assignments can be modified further, so the notation can be iterated, and we get things like this:

$$\begin{align*}
\left[\begin{array}{c}
1 \rightarrow \text{John} \\
2 \rightarrow \text{Mary}
\end{array}\right]^{\text{Mary}/2, \text{Sue}/1} & = \\
\left[\begin{array}{c}
1 \rightarrow \text{Sue} \\
2 \rightarrow \text{Mary}
\end{array}\right] & =
\end{align*}$$

The bracketing here was added for clarity. In practice, we will write

$[1 \rightarrow \text{John}]_{\text{Mary}/2, \text{Sue}/1}$
but it must be understood that this refers to the result of first modifying [1 \rightarrow John] so as to assign Mary to 2, and then modifying the result of that so as to assign Sue to 1.

Reference to modified assignments allows us to manipulate the values of certain designated variables (that is, traces or pronouns) while leaving the values of all other variables intact. The usefulness of this device will be seen when we apply our new rule of Predicate Abstraction. This is given in (16).

(16) *Predicate Abstraction Rule* (PA)

If $\alpha$ is a branching node whose daughters are $\beta_i$ and $\gamma$, where $\beta$ is a relative pronoun or “such”, and $i \in \mathbb{N}$, then for any variable assignment $\alpha$, $\int\alpha = \lambda x \in D \cdot \int\gamma$.

Every application of the Predicate Abstraction Rule targets one variable (identified by its index) in an expression $\gamma$, and defines a function of type $<e,t>$ by manipulating the value assigned to that particular variable. Let's see how this works by computing the denotation of our problematic “such that” clause in (5") (repeated from above).

(5")

```
NP
  NP
     AP
       man such_2
         CP
           that
             S
               Mary
                 VP
                   reviewed
                     DP
                       the
                         NP
                           NP
                              CP
                                 book wh_1
                                   S
                                     he_2
                                       VP
                                         wrote t_1
```
What we want to show is that the constituent “such\_2 that Mary reviewed the book wh\_1 he\_2 wrote t\_1” denotes a certain function in $D_{x<,x>}$: namely, the one that maps to 1 exactly those individuals $x$ such that Mary reviewed the book written by $x$. More precisely, we want to show that this constituent denotes this function independently of any assignment. So let’s calculate:\[21\]

\[
\begin{align*}
\llbracket \text{such\_2 that Mary reviewed the book wh\_1 he\_2 wrote t\_1} \rrbracket &= (\text{by convention } (11)) \\
\llbracket \text{such\_2 that Mary reviewed the book wh\_1 he\_2 wrote t\_1} \rrbracket^0 &= (\text{by PA}) \\
\lambda x \in D . \ & \llbracket \text{that Mary reviewed the book wh\_1 he\_2 wrote t\_1} \rrbracket^0/x_2 &= (\text{by definition } (13) \ (\text{assignment modification})) \\
\lambda x \in D . \ & \llbracket \text{that Mary reviewed the book wh\_1 he\_2 wrote t\_1} \rrbracket^{[2 \rightarrow x]} &= (\text{by vacuity of that and three applications of FA}) \\
\lambda x \in D . \ & \llbracket \text{reviewed} \rrbracket^{[2 \rightarrow x]}((\llbracket \text{the} \rrbracket^{[2 \rightarrow x]}((\llbracket \text{book wh\_1 he\_2 wrote t\_1} \rrbracket^{[2 \rightarrow x]}))((\llbracket \text{Mary} \rrbracket^{[2 \rightarrow x]})) &= (\text{by convention } (9) \ (\text{section } 5.2 \ \text{and lexical entries for review, the, Mary}) \\
\lambda x \in D . \ & \text{Mary reviewed the unique } y \text{ such that} \\
\llbracket \text{book wh\_1 he\_2 wrote t\_1} \rrbracket^{[2 \rightarrow x]}(y) &= 1 &= (\text{by PM and lexical entry for book}) \\
\lambda x \in D . \ & \text{Mary reviewed the unique } y \text{ such that} \\
y \text{ is a book and } \llbracket \text{wh\_1 he\_2 wrote t\_1} \rrbracket^{[2 \rightarrow x]}(y) &= 1
\end{align*}
\]

**Exercise**

Continue the computation started above.

### 5.4 What is variable binding?

In this section, we introduce definitions and theorems for some important notions related to variable binding. The whole section is necessarily more technical and more abstract than the previous parts of this book, and can be postponed\[22\] if you are satisfied with a merely informal understanding of the notion “variable binding” at this point.
5.4.1 Some semantic definitions

We have informally referred to traces and pronouns as "variables". What do we mean by this? The term "variable", though it clearly began its history as an (informal) semantic concept, is nowadays sometimes used in a purely syntactic sense, with explicit disclaimers that variables have any particular semantic interpretation in common. This practice may be unobjectionable in certain contexts, but not when matters of semantic interpretation and the syntax–semantics interface are at the very center of attention, as they are here in this book. We will therefore use the term "variable" (and various related terms) in a purely semantic sense. By this we don't mean that variables are not syntactic objects. They are. They are linguistic expressions. But what defines them as "variables" is not their shape or syntactic behavior, but the fact that they are interpreted in a certain way.

A variable in our sense is by definition something whose denotation varies with the assignment. More precisely:

(1) A terminal symbol $\alpha$ is a variable iff there are assignments $a$ and $a'$ such that $\langle a \rangle^a \neq \langle a \rangle^{a'}$.

Consider in the light of this definition our Traces and Pronouns Rule from above:

(2) Traces and Pronouns Rule
   If $\alpha$ is a pronoun or a trace, $a$ is a variable assignment, and $i \in \text{dom}(a)$, then $\langle \alpha \rangle^a = a(i)$.

It follows directly from that rule (and plausible assumptions about $D$) that traces and pronouns are variables in the sense of (1). Terminal symbols that are not variables are called "constants":

(3) A terminal symbol $\alpha$ is a constant iff for any two assignments $a$ and $a'$, $\langle \alpha \rangle^a = \langle \alpha \rangle^{a'}$.

Let's consider some further concepts that have to do with variables: in particular, the notions of "bound" versus "free" variables and of a "variable binder". These terms as well have been adapted to purely syntactic uses, but we are interested here in their semantic senses.

Roughly, what is meant by "variable binding" is any semantic operation which removes (or reduces) assignment dependency. By combining an expression whose denotation varies across assignments with one or more variable binders, we can create a larger expression whose denotation is assignment-invariant. This
is what happens when we build a relative clause. The embedded S-constituent (for example, "Joe likes t₁") has different semantic values under different assignments, but the complete relative clause (here "who₁ [Joe likes t₁]") has a fixed semantic value, which stays the same under all assignments. In our current semantics for English, Predicate Abstraction is in fact the only rule that accomplishes variable binding in this sense. Accordingly, the only items which qualify as variable binders are the ones which trigger the application of this rule: namely, indexed relative pronouns and "such". Here is a precise characterization of which expressions of a language L qualify as variable binders.

(4) An expression α is a variable binder (in a language L) iff there are trees β (in L) and assignments a such that

(i) β is not in the domain of [D]₀, but

(ii) some tree (of L) whose immediate constituents are α and β is in the domain of [D]₀.

To appreciate how this definition works, it is useful to be aware of the following general fact: Whenever any expression is in the domain [D]₀ for any assignment a, it is also in the domain of [D]₀' for any larger assignment a' ≥ a. So if an expression is not in the domain of [D]₀ for a given a, this can only be because the domain of a is too small, never because it is too large. Hence the kind of assignment which satisfies conditions (i) and (ii) of (4) has to be an assignment whose domain is too small to interpret β, yet big enough to interpret the next higher node.

Consider an example: How does, say, "who₁" qualify as variable binder of English under (4)? Well, pick the expression "t₁ left" and the assignment ∅, for example. They meet the conditions (4)(i) and (4)(ii) on β and a respectively, since (i) "t₁ left" is not in the domain of [D]₀, and (ii) the tree "[who₁ t₁ left]" is in the domain of [D]₀. (We can show (i) and (ii) by applying the relevant lexical entries and composition rules.) For a negative example, why is, say, "the" not a variable binder of English? Because we can prove from its lexical entry and the FA rule that whenever an NP β is not in the domain of [D]₀, then a DP consisting of the determiner "the" and the NP β isn’t either.

The related notions of "bound" and "free" occurrences of variables can be defined along similar lines. These terms (unlike the ones defined so far) apply to particular occurrences of expressions in a tree, not to the expressions (considered as types) themselves. We will see that the same variable may have bound and free occurrences in a given tree. The basic intuition behind the "free" versus "bound" distinction is this: A variable is "free" in a tree if that tree can only be interpreted under a specified assignment to this variable. For instance, the object trace t₁ is free in the trees [v₁ likes t₁], [ Mary likes t₁], [ it is not the case that he₂ likes t₁], and [ man who₂ t₁ likes t₁], since all these trees have well-defined
denotations only under assignments whose domains include 1 and thus specify a value for "t_1". By contrast, the object trace "t_1" is bound (= not free) in the trees [{_Mary likes t_1}, and [{_he_2 likes t_1}], since these trees do have denotations under assignments that exclude 1. An example with bound and free occurrences of the same variable in different places in the tree would be [{_t_1 likes the man who t_1 left}], in which the first occurrence of "t_1" is free and the second bound.

To formulate a precise definition, we need a way of referring to occurrences of expressions. Let's assume that the occurrences of each given expression \(a\) in a tree can be numbered \(a^1, a^2, \ldots\) (say, from left to right).

(5) Let \(a^n\) be an occurrence of a variable \(a\) in a tree \(\beta\).

(a) Then \(a^n\) is free in \(\beta\) if no subtree \(\gamma\) of \(\beta\) meets the following two conditions:
   (i) \(\gamma\) contains \(a^n\), and
   (ii) there are assignments \(a\) such that \(a\) is not in the domain of \([\gamma]\), but \(\gamma\) is.

(b) \(a^n\) is bound in \(\beta\) iff \(a^n\) is not free in \(\beta\).

Exercise

Apply these definitions to the examples given in the text right above.

It is a consequence of definition (5) that \(a^n\) can be bound in \(\beta\) only if \(\beta\) contains a variable binder in the sense of definition (4). In fact, something more specific follows: namely, that \(\beta\) must contain an occurrence of a variable binder in \(\beta\) which \(c\)-commands \(a^n\). Let us examine why this is so.

Suppose \(a^n\) is bound in \(\beta\). Then definition (5) implies that there is some subtree of \(\beta\) which fulfills the conditions (i) and (ii). Let \(\gamma\) be the smallest such subtree of \(\beta\). Let \(\delta\) be that daughter of \(\gamma\) which contains \(a^n\), and \(\epsilon\) that daughter of \(\gamma\) which doesn't contain \(a^n\). Let \(a\) be an assignment such that \(\gamma\), but not \(a\), is in the domain of \([\gamma]\). By assumption, \(\gamma\) is the smallest subtree of \(\beta\) which meets (i) and (ii), so \(\delta\) (which is smaller) doesn't. In particular, \(\delta\) must fail (ii), since it does meet (i) by assumption. Hence it cannot be the case that \(\delta\) is in the domain of \([\gamma]\). But this implies that \(\epsilon\) meets the definition of "variable binder" in (3). QED.

We have shown that for every bound variable there is a variable binder which is responsible for its being bound. This being so, it makes sense to define a 2-place relation "binds", which obtains between an occurrence of a variable binder and an occurrence of a variable:
(6) Let $\beta^n$ be a variable binder occurrence in a tree $\gamma$, and let $\alpha^m$ be a variable occurrence in the same tree $\gamma$ which is bound in $\gamma$.

Then $\beta^n$ binds $\alpha^m$ iff the sister of $\beta^n$ is the largest subtree of $\gamma$ in which $\alpha^m$ is free.

It follows directly from (4), (5), and (6) that (i) every bound variable is bound by a unique binder, and that (ii) no free variable is bound by anything. (By (ii), we mean more precisely: if a variable is free in $\gamma$, then nothing in $\gamma$ binds it.)

**Exercise**

What binds what in the following structure:

```
NP
  / \   
NP  AP
  /  /   
man such
   /    
CP
  /  
that
   /  
S
  /  
Mary
   /  
VP
  /  
reviewed
   /  
DP
  /  
the
   /  
NP
  /  
CP
  /  
book wh
   /  
S
  /  he
   /  wrote
    /  t
```

Show how the present definition applies to each binder–bindee pair.
5.4.2 Some theorems

The concepts we have introduced in this section provide a convenient and widely used terminology in which we can state a number of generalizations which turn out to be predicted by our current semantics for English phrase structures. For example, we can give a precise characterization of the syntactic configurations in which variables get bound. First, we can prove that there is no binding without co-indexing:

(7) If \( \beta \) binds \( \alpha \), then \( \beta \) and \( \alpha \) are co-indexed.

Exercise

Prove (7).

Using (7) and our earlier definition of "variable binding", we can state necessary and sufficient conditions on the syntactic configurations in which variables get bound:

(8) \( \beta \) binds \( \alpha \) if and only if

(i) \( \alpha \) is an occurrence of a variable,
(ii) \( \beta \) is an occurrence of a variable binder,
(iii) \( \beta \) c-commands \( \alpha \),
(iv) \( \beta \) is co-indexed with \( \alpha \), and
(v) \( \beta \) does not c-command any other variable binder occurrence which also c-commands and is co-indexed with \( \alpha \).

Exercise

Prove (8).

We can also state a precise theorem about how the extent to which the denotation of a tree depends on the choice of assignment correlates with the presence of free variables in that tree. The general version of this theorem is (9), and a special case is (10).
(9) If a tree $\gamma$ contains no free occurrences of variables indexed $i$, then for all assignments $a$ and all $x \in D$:

either $\gamma$ is in the domain of neither $[\ ]^a$ nor $[\ ]^{x^b}$, 
or it is in the domain of both, and $[\gamma]^a = [\gamma]^b$.

(10) If a tree $\gamma$ contains no free variables, then it either has no semantic value under any assignment, or else it has the same one under all of them.

Here is a proof-sketch for (9). We begin by observing that (9) holds trivially when $\gamma$ is a terminal node. Then we show that, if $\gamma$ is a non-terminal node, and if we can take for granted that all of $\gamma$'s daughters conform to (9), it follows that $\gamma$ itself will too. To do this, we consider five cases. The first case (trivial) is that none of our composition principles is applicable to $\gamma$. The second case (also trivial) is that NN applies to $\gamma$. The remaining three cases correspond to the composition principles FA, PM, and PA. We first show that if FA applies to $\gamma$, then $\gamma$ must conform to (9) whenever both its daughters do. Then we do the same for PM, and finally for PA. When we are done with all these subproofs, we have shown that, however $\gamma$ is constructed, it must conform to (9).

Why is it useful to be aware of laws like (7)–(9)? It can often save us lengthy calculations when we want to assess whether a given syntactic representation does or does not represent a given denotation. For instance, by relying on (7) and (9) we can tell very quickly that the predicate such$_1$ that Mary reviewed the book wh$_2$ he$_2$ wrote t$_2$ denotes a constant function (and thus cannot possibly single out a non-empty proper subset of the men). Using (8), we determine that the S beginning with Mary contains no free variable, and given (9), this implies that this S is either true under all assignments or false under all assignments. We need not derive predictions like this by doing all our semantic calculations from scratch with each new example. Once the appropriate general theorems have been proved, many questions about the interpretation of a given syntactic structure can be answered simply by taking a glance at it. Experienced semanticists use such shortcuts all the time, and we will increasingly do so in the later chapters of this book. Still, we must never lose sight of the fact that all predictions about which structure means what must be deducible from the lexical entries and composition principles, and from those alone. If they are not, they are not predictions at all.

### 5.4.3 Methodological remarks

We have stressed that we are interested in the semantic senses of the concepts "variable", "bound", "free", etcetera. We acknowledged in passing that other
authors use these terms in purely syntactic senses, but have made it sound as if this was a relatively recent development in the syntactic literature, and that tradition was on our side. "What tradition?," you may be wondering, especially if you have some background in formal logic. Logic books invariably give purely syntactic definitions of these terms. For instance, a typical presentation of the syntax of Predicate Logic (PL) may begin with the definition of a vocabulary, including a clause like: "the variables of PL are $x, y, z, x_1, y_1, z_1, x_2, y_2, z_2, \ldots$". So the term "variable" is defined simply by a list, or by a specification of the shapes of variables. Typical definitions of "bound" and "free" for variables in PL formulas read like this: "An occurrence of a variable $\alpha$ in a formula $\phi$ is free in $\phi$ iff this occurrence of $\alpha$ does not fall within any subformula of $\phi$ that has the form $\forall \alpha \psi$ or $\exists \alpha \psi$." If this is not a purely syntactic definition, then what is? Apart from superficial differences in terminology (for example, logicians don't say "c-command"), the standard logic-book definition of "variable binding" amounts to exactly the biconditional that we presented as a "theorem" in (8) above – except that the term "variable binder" is typically replaced by a list of the relevant symbols: namely, "$\forall \alpha$" and "$\exists \alpha$" in the case of PL.

So it looks as if we don’t really have tradition on our side when we insist that "variable" and "binding" be defined in terms of their semantic import, or when we say that (8) describes the empirical predictions of a particular semantic theory and is not true by mere definition.

However, the picture changes somewhat if we attend not just to the official definitions, but also to the informal terminological practices in the logic literature. While the explicit definitions for "variable," "bound", etcetera are syntactic definitions, there is also an informal, but very firmly established, consensus among logicians that it would be inappropriate to employ these terms for expressions that do not share essentially the same semantics. This emerges more clearly if you don’t just look at presentations of standard systems like PL, but observe the terminological choices that a logician will make when contemplating innovations or comparing different systems. No logician will call a "variable" something that is interpreted as a constant (in the sense captured by a semantic definition like our (3)). And no logician would alter the definition of "bound"/"free" (perhaps with the aim of making it simpler) without proposing exactly concomitant changes in the semantics. (The definition is clearly not simple or natural from a syntactic point of view; this in itself is an indication of the fact that the concept it aims to formalize is not a syntactic, but a semantic, one.) We are not claiming that all commonly accepted uses of the relevant terminology are compatible with the exact definitions we have formulated here. In fact, perhaps none of them are, since these particular definitions make sense only in the context of a particular formalization of the semantics of variables. Even relatively superficial technical changes would require revisions in the definitions. Still, we think that the concepts we have defined can be recognized as formal
counterparts of the informal concepts which have guided the terminological practice (if not the explicit definitions) of logicians and formal semanticists.

So, as far as the question "What tradition?" is concerned, we note that there is more to a logician's concept of variable-hood and binding than what is revealed in the letter of the standard definitions. We have tried to elucidate some of this. To what end? Why do we, as linguists, feel the need to keep syntactic and semantic terminology neatly apart where logicians have been quite content to conflate them? This has to do with a fundamental difference between formal and natural languages, or between logic and linguistic semantics. The formal languages defined by logicians are specifically designed to receive a certain semantic interpretation. Their vocabulary and syntax are deliberately set up in such a way that vocabulary items that are interpreted alike are also alike in shape and distribution. It would be pointless for the inventor of a logical language to introduce syntactically distinct vocabulary items (say, "pronouns" and "traces", or "nouns" and "verbs") only to give them identical interpretations. Therefore, the syntactic categories of these artificial languages correspond straightforwardly to classes of expressions with shared semantic properties. It is not a discovery or a matter of possible controversy that this should be so. It is therefore harmless, and even convenient, to use the very terms that informally characterize a certain semantic behavior as labels for the artificial syntactic categories. But what's harmless and convenient for logicians may be confusing and counterproductive for linguists. When it comes to theorizing about natural language, we don't know ahead of the endeavor what general properties we will find. Perhaps natural languages will turn out to have some of the design features of logical languages, such as a more or less simple correspondence (or even a one-to-one match) between syntactic and semantic categories. This could be an interesting finding, and it looks like an interesting working hypothesis to explore. But it is hardly obvious, and it is definitely controversial. Whether we want to argue for it or against it, we are better off with a terminology that allows us to state the issue than with one that equivocates.

### 5.5 Interpretability and syntactic constraints on indexing

In this section, we will reflect on the proper assignment of indices to pronouns and traces. Look at the analysis of our example from the previous section, for example.

(1) man such that Mary reviewed the book which he wrote
We succeeded in predicting that this NP can have the reading it intuitively has. That is, we found a structure for (1) which, by our revised semantic rules, demonstrably expresses the intended meaning:

(2) man such\(_1\) that Mary reviewed the book wh\(_2\) he\(_1\) wrote \(t_2\)

But we have yet to show that our grammar correctly predicts the meaning of (2) to be the only reading of (1). In other words, we have yet to make sure that we don’t generate any additional, intuitively unavailable, meanings for (1).

The prima facie problem here is that the indices, which we saw play a crucial role in determining denotation, are not overt on the surface. So in principle, the surface structure of (1) could be associated with as many different structures as there are different possible indexings. Many of the infinitely many options that arise here are semantically equivalent. Still, there are quite a few non-equivalent choices for assigning indices to the four items “such”, “wh”, “he”, and “t” in (1). What about (3), for example?

(3) man such\(_1\) that Mary reviewed the book wh\(_2\) he\(_2\) wrote \(t_1\)

(3) has a meaning that is easy to express in English: namely, the meaning of “man such that Mary reviewed the book that wrote him”. Why can (1) not be read in this way? Because, we assume, the structure in (3) is syntactically ill-formed. It is simply not made available by the syntactic component of the grammar. If it were well-formed, we claim, then (1) could be understood in the sense indicated. This is what our semantics commits us to, at any rate, and as far as we can see, it is consistent with the evidence.

From a syntactician’s point of view, the assumption that (3) is syntactically ill-formed is unobjectionable. In fact, it exhibits simultaneous violations of several commonly proposed syntactic constraints. One problem is that “such” in (3) binds a trace instead of a pronoun. That this in itself leads to ungrammaticality can be seen in simpler examples like (4).

(4) the man such\(_1\) that he/_\(t_1\) left

Another problem is that the relative pronoun in (3) binds a pronoun rather than a trace. This can also be seen to be a sufficient reason for ungrammaticality, at least when the pronoun is so close to its binder:

(5) the man who\(_1\) *he\(_1\)/_\(t_1\) left

Over longer distances, English sometimes tolerates such “resumptive pronouns” fairly well (for example, “the man who, Mary wonders where he\(_1\) went”), and other languages allow them freely, even in local configurations like (5). This sort
of cross-linguistic and language-internal variation makes it all the more plausible that our semantics is on the right track in giving pronouns and traces the same interpretation. It is not inconceivable a priori that the different distribution of pronouns and traces in English ultimately derives from some nonobvious interpretive difference between the two, or that there could be differences between the vocabularies of individual languages in this regard. But at the current state of linguistic research, there is less evidence for this than for the hypothesis that the distributional regularities have syntactic explanations compatible with an indiscriminate semantics.

Yet another thing that seems to be wrong with (3) has to do with the distribution of gender features: the masculine feature of “he” somehow clashes with the neuter feature of “book”. Again, simpler examples than (3) illustrate the phenomenon more clearly:

(6) the book such \textsubscript{1} that Mary reviewed *him\textsubscript{1}/it\textsubscript{1}

It is less clear than in the cases of (4) and (5) that what we are looking at here is a matter of syntactic well-formedness. We could treat it as that, by laying down an appropriate set of syntactic agreement principles. But lay opinion points towards a different approach. The naive answer to what is wrong with “the book such that Mary reviewed him” is that “him” can’t refer to an inanimate thing like a book, which sounds like a semantic explanation. Could this be spelled out in our theory? Taken literally, we cannot accept the naive story. Bound variables don’t refer to individuals in the first place. As we have argued at the beginning of the chapter, the “it” in the good variant of (6) does not refer to a book either. But there seems to be a way of rephrasing the basic intuition in the light of this objection: suppose “him\textsubscript{1}” cannot denote an inanimate (or non-male) individual under any assignment. More precisely, for any assignment a such that a(1) is not male, “him\textsubscript{1}” is not in the domain of \[[\text{ ]}]\textsuperscript{a}. What would this predict for (6)?

The immediate prediction is that “Mary reviewed him\textsubscript{1}” will have a truth-value only under those assignments which map 1 to a male. For assignments a which don’t meet this condition, \[[\text{Mary reviewed him\textsubscript{1}}]\textsuperscript{a} is undefined. “Mary reviewed it\textsubscript{1}”, by contrast, gets a truth-value only under those assignments which map 1 to a nonhuman. What happens, then, when we apply the predicate abstraction rule to compute the denotation of \[[\text{such \textsubscript{1} that Mary reviewed *him\textsubscript{1}/it\textsubscript{1}}]?

Here we need the pedantic version, which produces partial functions:

(7) Predicate Abstraction (pedantic version)

If \(\alpha\) is a branching node whose daughters are \(\beta\) and \(\gamma\), where \(\beta\) is a relative pronoun or “such”, and \(i\) \in \text{IN}, then for any variable assignment a, \(\llbracket \alpha \rrbracket^a = \lambda x : x \in D \text{ and } \gamma \text{ is in the domain of } \llbracket \text{ ] }^a \rrbracket^{a\gamma} \).
Given (7), we obtain:

(a) \[ \text{such}_1 \text{ that Mary reviewed him}_1 \] = 
\[ \lambda x : x \in D \text{ and } x \text{ is male. Mary reviewed } x. \]

(b) \[ \text{such}_1 \text{ that Mary reviewed it}_1 \] = 
\[ \lambda x : x \in D \text{ and } x \text{ is nonhuman. Mary reviewed } x. \]

If we combine \[ \text{book} \] with the partial function in (8a), we need the pedantic version of the Predicate Modification rule:

(9) Predicate Modification (pedantic version)
If \( \alpha \) is a branching node and \( \{ \beta, \gamma \} \) the set of its daughters, then, for any assignment \( a \), \( \alpha \) is in the domain of \( \llbracket \alpha \rrbracket^a \) if both \( \beta \) and \( \gamma \) are, and \( \llbracket \beta \rrbracket^a \) and \( \llbracket \gamma \rrbracket^a \) are both of type \( \langle e, t \rangle \). In this case, \( \llbracket \alpha \rrbracket^a = \lambda x : x \in D \text{ and } x \text{ is in the domain of } \llbracket \beta \rrbracket^a \text{ and } \llbracket \gamma \rrbracket^a. \llbracket \beta \rrbracket^a(x) = \llbracket \gamma \rrbracket^a(x) = 1. \]

Given (9), the denotation of \( \text{book such}_1 \text{ that Mary reviewed him}_1 \) is a function that does not map anything to 1, regardless of what the facts are: only male individuals are in its domain, and of those, none are books. Therefore, the definite DP \( \text{the book such}_1 \text{ that Mary reviewed him}_1 \) can never get a denotation. Combining \( \llbracket \text{book} \rrbracket \) with the function in (8b), on the other hand, yields a function which maps all books that Mary reviewed to 1, and all nonhuman things which are either not books or are not reviewed by Mary to 0 (and is undefined for humans). Under the appropriate circumstances (that is, if there happens to be a unique book reviewed by Mary), \( \text{the book such}_1 \text{ that Mary reviewed it}_1 \) thus denotes.

This is only a sketch of a semantic account of gender agreement, but it looks as if it could be made to work. For the purposes of this book, we may leave it open whether feature mismatches as in (3) and (6) are ruled out by the syntax or by the semantics.

Besides (2) and (3), there are many other conceivable indexings of (1). But (equivalent variants aside), these two were the only possibilities in which each of the two binders binds a variable. What about indexings which don’t have this property, say (10)?

(10) \( \text{man such}_1 \text{ that Mary reviewed the book \textit{wh}_1 \textit{he}_1 \textit{wrote \textit{t}_1} } \)

Here the \( \text{such}_1 \) binds no variable. (This is so despite the fact that it c-commands and is co-indexed with two variables. Those variables are already bound by the lower \( \text{wh}_1 \).) It is a so-called vacuous binder. Syntacticians have proposed that vacuous binding configurations are generally ill-formed:

(11) Each variable binder must bind at least one variable.
Again, simpler examples than (10) are more appropriate to motivate (11), since (10) violates other conditions as well (see below). Consider (12).

\[(12) \quad (a) \quad \text{*the man such that Mary is famous} \]
\[(b) \quad \text{*the man who Mary is famous} \]

Here there simply are no variables present at all, so the binders in (12a) and (12b) will be vacuous under any possible indexing. Plausibly, this is the reason why these examples are ungrammatical.

Notice that (12a) and (12b) are not uninterpretable. Our semantics predicts, in fact, that if Mary happens to be famous and there happens to be just one man in D, then (12a) and (12b) denote this unique man. Otherwise they denote nothing. So we predict that (12a) and (12b) are tantamount to the simple DP the man, except that they convey an additional presupposition: namely that Mary is famous. Similarly, we do derive a well-defined meaning for (10) – in fact, the same meaning that we derived for (1) before we introduced the refinements of section 5.3 – see above. It is not \textit{prima facie} inconceivable that natural language might have expressed such meanings, and might have expressed them in these forms. But (12a) and (12b) cannot be used with this (or any) meaning. So we need a syntactic constraint like (11) to exclude them.

(11) takes care of ruling out (10) and all other indexings of (1) in which one or both binders wind up vacuous. As we already noted, we have thus considered all possible indexings of (1), since there are only as many variables as binders in this example. Other examples, however, offer additional options, and allow us to illustrate further syntactic constraints on indexing.

For instance, neither of the two indexings we give in (13) below violates any of the syntactic constraints we have mentioned so far. Yet only one represents a meaning that is available for the corresponding surface string.

\[(13) \quad \text{the man who}_1 t_1 \text{ talked to the boy who}_2 t_2 \text{ visited him}_1/2 \]

On the starred indexing, our semantics interprets (13) as denoting the man who talked to the boy who visited \textit{himself}. This is a typical violation of a well-known binding constraint that says that a non-reflexive pronoun cannot be co-indexed with a c-commanding position in the same minimal clause. The precise formulation of the relevant constraint has been a topic of much syntactic research. Our aim here is not to contribute to this, but just to point out where binding constraints fit into our overall picture: The structures that violate such binding constraints are no less interpretable than those that conform to it. Binding constraints belong to syntax: they talk about purely formal aspects of syntactic representations, specifically about indices, c-command, and morphological properties of indexed DPs. Any semantic predictions they imply – that is,
predictions about possible and impossible meanings for a given sentence — come about indirectly, and depend crucially on the semantics that interprets indexed structures.

In sum, we argued that our semantics for variables and variable binders makes correct predictions about the meanings of a wide range of relative clause constructions, provided it is combined with suitable syntactic constraints which cut down on the number of possible indexings. The constraints which we found it necessary to appeal to all appeared to be well-established in syntactic theory.

Notes


2 Quine distinguishes between *absolute general terms* and *relative general terms*. By an “absolute (general) term” he means a 1-place predicate, i.e., an expression that is true or false of an object. In our terms, this would be an expression with a meaning of type <e,t>, e.g., an intransitive verb or common noun. By a “relative (general) term” he means a 2- (or more) place predicate. His examples are “part of”, “bigger than”, “exceeds”.

3 A *singular term* is any expression with a denotation in \( D_e \) (e.g., a proper name).

4 An “adjective” in Quine’s sense seems to be anything that can modify a noun. It needn’t have a head of the syntactic category A.

5 *Conversion* is the operation that switches the two arguments of a 2-place predicate; e.g., it maps “taller” to “shorter”, and “buy” to “be-bought-by”. *Application* is the combination of a 2-place predicate with an argument to yield a 1-place predicate; e.g., it forms “likes John” from “likes” and “John”.

6 These judgments about (1) and (2) must be understood with the usual qualification: viz., that we pretend that the domain \( D \) contains only those individuals which are considered relevant in the context in which these sentences are uttered. (2) does not really presuppose that there is exactly one house in the world. But it does presuppose that only one house is under consideration, and this presupposition makes it appreciably odd in some utterance contexts where (1) would be entirely natural.

7 We will not always write both the \( wh \)-word and “that” in the trees below.

8 Richard Montague treated pronouns, proper names, and quantifier phrases as “term phrases”. Our current NPs would correspond to Montague’s “common noun phrases”.

9 Remember that we are here concerned with restrictive relatives. The criticism that follows would not apply to nonrestrictives.

10 See section 4.5, where we discussed the analogous bracketing with a PP-modifier.

11 The fully explicit versions of (12)–(14), which take possible undefinedness into account, are as follows:
Relative Clauses, Variables, Variable Binding

If α is a non-branching node and β its daughter, then, for any assignment a, α is in the domain of \( [\alpha]^a \) if β is. In this case, \( [\alpha]^a = [\beta]^a \).

If α is a branching node and \( \{\beta, \gamma\} \) the set of its daughters, then, for any assignment a, α is in the domain of \( [\alpha]^a \) if both β and γ are, and \( [\beta]^a \) is a function whose domain contains \( [\gamma]^a \). In this case, \( [\alpha]^a = [\beta]^a(\gamma)^a \).

If α is a branching node and \( \{\beta, \gamma\} \) the set of its daughters, then, for any assignment a, α is in the domain of \( [\alpha]^a \) if both β and γ are, and \( [\beta]^a \) is a function whose domain contains \( [\gamma]^a \). In this case, \( [\alpha]^a = [\beta]^a(\gamma)^a \).

If α is a branching node whose daughters are a relative pronoun and β, then \( \alpha \) is a variable assignment, and i ∈ dom(\( a \)), then \( \alpha_i \) is in the domain of \( [\alpha]^a \), and \( (\alpha)^a = a(i) \).

Pedantic version: α is in the domain of \( [\alpha]^a \) only if α has a semantic value under all assignments, and the same one under all assignments.

Pedantic version: If α is a branching node whose daughters are β, and γ, where β is a relative pronoun or “such”, and i ∈ \( iN \), then for any variable assignment a, \( [\alpha]^a = \lambda x : x \in D \) and γ is in the domain of \( [\alpha]^a \), and \( [\gamma]^a \).

In the presentation of this calculation, we are writing mere strings between the double brackets. Strictly speaking, what should appear there, of course, are the corresponding phrase structure trees. This is merely a shorthand. We mean each string to stand for the subtree that dominates it in (5") above. So please look at that diagram (5") in order to follow the calculation.

The semantic notion of “variable binding” will be needed in chapter 10.

(4) presupposes that all trees are at most binary-branching. This is just to simplify matters. (4) is a special case of a more general definition, according to which α is a variable binder iff there are trees \( \beta_1, \ldots, \beta_n \) (for some \( n \geq 1 \)) and assignments a such that (i) none of the \( \beta \) are in the domain of \( [\beta]^a \), but (ii) a tree whose immediate constituents are α and \( \beta_1, \ldots, \beta_n \) is in the domain of \( [\beta]^a \).
We do not give a proof of this law here, but take a minute to think about why it holds true. For lexical items, it follows from convention (9) in section 5.2. For traces and pronouns, it follows directly from the Traces and Pronouns Rule. This takes care of all possible terminal nodes. Now if we take any one of our composition rules, we can show that, as long as the law is true of all the daughter nodes, the application of the rule will ensure that it is true of the mother node as well.

We use superscripts here to avoid any confusion with the subscript indices, which are parts of the variables (qua expressions) themselves.

"Subtree" is understood in such a way that β counts as a subtree of itself.

"C-command" in the sense of the definition: x c-commands y iff x is sister to a node which dominates y.
6 Quantifiers: Their Semantic Type

The only DPs treated so far (not counting predicate nominals) have been proper names, definite descriptions, pronouns, and traces. We assumed that their denotations were individuals – that is, elements of $D_e$. There are many other kinds of DPs – for example, DPs made up with a variety of determiners (“this”, “that”, “a(n)”, “every”, “no”, “many”, “few”, “most”, and so forth) and yet other types, such as “only John”. What denotations should we assign to all these DPs and to the determiners they contain? What types of denotations will be suitable in each case?

Before we address these questions directly, we will establish a couple of important negative points. First, we show that at least some DPs do not denote *individuals*. Second, we are going to see that *sets of individuals* (or their characteristic functions) are not a suitable type of denotation either. We will consider more arguments than are strictly speaking needed to make these points. It may seem like overkill, but it serves the additional purpose of drawing up a checklist against which to test positive proposals, in particular the one adopted below.

6.1 Problems with individuals as DP-denotations

There are various types of DPs for which denotations of type e seem to work well. We had some success treating proper names in this way, and also definite descriptions. Pronouns (“he”, “I”, . . . ) and demonstratives (“this book”, “that cat”) might also be accommodated if we allowed that their extensions are not fixed by the semantics once and for all, but vary from one occasion of utterance to the next. “I”, for instance, seems to denote on each occasion when it is uttered the individual who utters it. This individual then enters the calculation of the truth-value of the uttered sentence in the usual way. For instance, when Irene Heim utters “I am sleepy”, this is true just in case Irene Heim is in the extension of “sleepy”, and false otherwise. We will give more detailed
consideration to the proper treatment of such context dependency later on. Even if denotations may be fixed in part by the utterance occasion, however, individuals don’t seem to be the right kind of denotation for many DPs. Let us see why not.

6.1.1 Predictions about truth-conditions and entailment patterns

Naive intuition is not a reliable guide as to which DPs denote individuals. If asked whom or what “no man” denotes, a lay person might say that it doesn’t denote anything, but with “only John”, they might say it refers to John.

As it turns out, the latter suggestion is immediately falsified by its predictions. If \([\text{only John}] = \text{John}\), then \([\text{only John left}] = 1 \iff \text{John left}\). But if John and Sam both left, this sentence is intuitively false, not true as just predicted. Of course, this argument relies on a number of assumptions: in particular, assumptions about the constituent structure of the example and about the principles of semantic composition. It is thus not an indefeasible argument, but it is sound.

By comparison, the mere fact that it offends common sense to assign a referent to “no man”, is at best a weak reason not to do so in our theory. More decisive, in this instance as well, will be considerations pertaining to predicted semantic judgments about sentences. Such considerations can be used to refute a particular assignment of denotation for a given DP, as in the case above, where we showed that \([\text{only John}]\) cannot be John. And sometimes they can even be used to show that no denotation of a certain type will do. Let us give a few illustrations of this type of reasoning.

DPs that fail to validate subset-to-superset inferences

The following inference is intuitively valid:

(1) John came yesterday morning.
\[\therefore \text{John came yesterday.}\]

Not only is it intuitively valid, but it is predicted to be valid by any semantics that implies these three assumptions:

(i) \([\text{John}] \in D_e\)
(ii) \([\text{came yesterday morning}] \subseteq [\text{came yesterday}]^4\)
(iii) A sentence whose subject denotes an individual is true iff that individual is a member of the set denoted by the VP.
Proof: Suppose the premise is true. Then, by (iii), \([\text{John}] \in [\text{came yesterday morning}]\). Hence, by (ii) and set theory, \([\text{John}] \in [\text{came yesterday}]\). And, again by (iii), the conclusion is true. QED. Notice that no concrete assumption about which element of \(D_e\) John denotes was needed for this proof.

Now consider the parallel inference in (2).

(2) \quad \text{At most one letter came yesterday morning.}
\quad \therefore \quad \text{At most one letter came yesterday.}

This one is intuitively invalid: It is easy to imagine a state of affairs where the premise is true and the conclusion is false. For example, only one letter comes in the morning but two more in the afternoon. If we want to maintain (ii) and (iii), it therefore follows that \([\text{at most one letter}] \in D_e\), or else we could prove the validity of (2) exactly as we proved that of (1). Should we give up (ii) or (iii) to avoid this conclusion? (ii) seems well-motivated by other data: for example, the tautological status of "If John came yesterday morning, he came yesterday". And giving up (iii) would set our whole enterprise back to square one, with no plausible substitute in sight. These would be costly moves at best.

We can test other DPs in this inference schema:

(3) \quad \alpha \text{ came yesterday morning.}
\quad \therefore \quad \alpha \text{ came yesterday.}

Among the substitutions for \(\alpha\) that systematically fail to make this valid are DPs with the determiners "no" and "few", and with all determiners of the form "less than \(n\)", "at most \(n\)", "exactly \(n\)" (for some numeral "\(n\)”). For all such DPs, we thus have a strong reason to assume that their denotations are not of type \(e\). (The reverse does not hold: there are many DPs that do validate (3) but still cannot have type \(e\) denotations, for other reasons, such as those considered next.)

**DPs that fail the Law of Contradiction**

If you choose two VPs with disjoint extensions and combine first one, then the other, with a given proper name, you get two sentences that contradict each other. For example, (4) is contradictory.

(4) \quad \text{Mount Rainier is on this side of the border, and Mount Rainier is on the other side of the border.}

Again, this is a judgment predicted by any semantics committed to a few plausible assumptions. The following four suffice:
(i) \([\text{Mount Rainier}] \in D_e\)
(ii) \([\text{be on this side of the border}] \cap [\text{be on the other side of the border}] = \emptyset\)
(iii) (composition rule for subject + VP, same as above)
(iv) standard analysis of and.

We leave the proof to the reader.

The important point is again that this proof does not rely on a more specific assumption about Mount Rainier than what type of denotation it has. It will therefore generalize to analogous sentences with any other DP that denotes an individual. But many such sentences are not in fact contradictory, for example:

(5) More than two mountains are on this side of the border, and more than two mountains are on the other side of the border.

So, unless we want to mess with (ii), (iii), or (iv), we have to conclude that more than two mountains does not denote anything in \(D_e\).

Further DPs to which this argument extends are “a mountain”, “n mountains” (for any numeral “n”), “no mountain”, and lots of others.

**DPs that fail the Law of Excluded Middle**

Again we form minimal pairs of sentences whose subjects are identical and whose VPs differ. This time we choose two VPs the union of whose extensions exhausts everything there is, and we coordinate the two sentences by “or”. For example:

(6) I am over 30 years old, or I am under 40 years old.

(6) is a tautology, which we can prove if our semantics implies this much:

(i) \([I] \in D_e\)
(ii) \([\text{be over 30 years old}] \cup [\text{be under 40 years old}] = D\)
(iii) (as above)
(iv) standard analysis of or.

Our reasoning must be foreseeable at this point: Any other individual-denoting DP in place of “I” in (6) would likewise be predicted to yield a tautology. So if there are counterexamples, the DPs in them cannot denote individuals. Here is one:

(7) Every woman in this room is over 30 years old, or every woman in this room is under 40 years old.
There are other such systematic differences in the entailments, tautologies, and contradictions that we get for proper names on the one hand and for the lot of nondefinite DPs on the other. They all yield potential arguments for a distinction in semantic type. But you have the idea, and we can turn to an argument of a somewhat different kind.

6.1.2 Predictions about ambiguity and the effects of syntactic reorganization

English syntax often allows us to construct different sentences out of more or less the same lexical items. Often we will get different meanings as a result. For instance, “John saw Mary” has different truth-conditions from “Mary saw John”, and this is easily accounted for in our current theory. Other times, we get pairs of sentences whose truth-conditions coincide. This occurs with active-passive pairs, with pairs of a topicalized sentence and its plain counterpart, or with certain more stilted circumlocutions like the “such that” construction:

(8a) I answered question #7.
(8b) Question #7, I answered.

(9a) John saw Mary.
(9b) Mary is such that John saw her.
(9c) John is such that Mary saw him.

Whatever the subtler meaning differences in each of these groups, we cannot imagine states of affairs in which one member of the pair or triple is true and the other(s) false. Our semantics (given suitable syntactic assumptions) predicts these equivalences. Let us briefly sketch how before we resume our main line of argument.

First, take the unlyrical (9b) and (9c). If you treat the pronouns “her” and “him” as variables and “such that” as signaling predicate abstraction, then 

\[[\text{such that John saw her}] = \text{(the characteristic function of) the set \{x : John saw x\}}\]. This is also the value of the whole VP, and (9b) is thus predicted to be true iff Mary is in this set – which means nothing more and nothing less than that John saw her. So we have proved (9b) equivalent to (9a). The same goes for (9c).

Second, consider the topicalization construction in (8b). This may be much like the “such that” construction. In the analysis of Chomsky’s “On Wh-Movement,” the topicalized phrase in (8b) would be linked to a moved wh-phrase, and we might assume a structure of the following kind:
The CP here looks like just another predicate abstract, and indeed interpreting it this way predicts the right meaning and the equivalence with (8a).

Why did we bring up these equivalences? Because we now want you to look at analogous cases in which the names and first person pronouns of (8) and (9) have been replaced by some other kinds of DPs:

(10a) Almost everybody answered at least one question.
(10b) At least one question, almost everybody answered.

(11a) Nobody saw more than one policeman.
(11b) More than one policeman is such that nobody saw him.
(11c) Nobody is such that he or she saw more than one policeman.

Suddenly the "transformations" affect truth-conditional meaning. (10a) can be true when no two people answered any of the same questions, whereas (10b) requires there to be some one question that was answered by almost all. For instance, imagine there were ten students and ten questions; student #1 answered just question #1, student #2 just question #2, and so on, except for student #10, who didn't answer any. (10a) is clearly a true description of this state of affairs, but (10b) is false of it. (Some speakers might decide upon reflection that they actually get two distinct readings for (10a), one of which is equivalent to (10b). But even so, the fact remains that (10b) cannot express the salient reading of (10a).)

Similar comments apply to (11a), (11b), and (11c). (11c) does not require for its truth that any policeman went completely unnoticed, whereas (11b) claims that two or more did. Conversely, (11b) is true and (11c) false when everybody saw several policemen and there were some additional unseen policemen hiding
in the shadows. (11a) seems to be equivalent to (11c), and hence distinct from (11b). Even if some speakers find (11a) ambiguous, (11b) clearly differs from (11c) on any reading of either, and this is enough for us to make our point.

The point is this: The truth-conditional effects that we have just seen to be associated with certain structural rearrangements are completely unexpected if the DPs in (10) and (11) are treated like those in (8) and (9). The basic semantic properties of the “such that” and topicalization constructions directly imply that shuffling around DPs with meanings of type e in these constructions cannot possibly change truth-values. So when we do observe such changes, as between (11b) and (11c), they are one more piece of evidence that certain DPs don’t denote individuals.

A final problem with trying to assimilate the semantics of all DPs to that of proper names (closely related, as it turns out, to the preceding puzzle) is that such an approach does not anticipate certain judgments of ambiguity. For example, (13) is ambiguous in a way that (12) is not.

(12) It didn’t snow on Christmas Day.

(13) It didn’t snow on more than two of these days.

Suppose ten days are under consideration, and it snowed on exactly three of them. Is (13) true or false then? This depends on how you understand it. I could argue: “(13) is true. There was no snow on as many as seven days, and seven is more than two, so surely it didn’t snow on more than two days.” You might argue: “No, (13) is false. It snowed on three days, and three is more than two, so it did snow on more than two.” We are not going to settle our argument, because each of us is reading the sentence in a different way, and both are apparently readings allowed by English grammar. (There may be differences in the ways the two readings are most naturally pronounced, and if so, these might be traceable to some subtle difference in syntactic structure. Should this turn out to be the case, we should really group this example with the ones in the previous section; that is, it would then be another instance of how changes in structure affect truth-conditions. For our purposes here, this wouldn’t make an important difference, since one way or the other, the example highlights a fundamental semantic difference between “more than two of these days” and a proper name like “Christmas Day”.)

There is no ambiguity of this kind in (12), and there can’t really be, given that “Christmas Day” denotes a day. The three ingredients of this sentence – the name Christmas Day, the predicate it snows on, and the negation – can only be put together to produce a truth-value in one way, so there is simply no room for any possible ambiguity. If DPs like more than two days likewise denoted individuals, there ought to be no more room for ambiguity in (13). (The argument as stated
is a bit simple-minded, we admit: one can think of more than one way to semantically compose the verb snow, the temporal preposition on, a time-name, and a negation. But since all plausible alternatives come to the same thing in the case of (12) – they had better, or else we would predict an ambiguity that doesn’t in fact exist! – it’s hard to see how they could wind up with different truth-values for (13).)

6.2 Problems with having DPs denote sets of individuals

The view that all DPs denote individuals may have been a dead horse as soon as anyone gave it conscious thought. But another similarly problematic idea still comes up every now and then. P. T. Geach introduces and warns against it as follows:

Another bad habit often picked up from the same training is the way of thinking that Frege called mechanical or quantificatious thinking: *mechanische oder quantifizierende Auffassung*. I have used a rude made-up word “quantificatious” because Frege was being rude; “quantificational” and “quantifying” are innocent descriptive terms of modern logic, but they are innocent only because they are mere labels and have no longer any suggestion of quantity. But people who think quantificatiously do take seriously the idea that words like “all”, “some”, “most”, “none”, tell us how much, how large a part, of a class is being considered. “All men” would refer to the whole of the class men; “most men”, to the greater part of the class; “some men” to some part of the class men (better not ask which part!); “no man”, finally, to a null or empty class which contains no men. One can indeed legitimately get to the concept of a null class – but not this way.

I have not time to bring out in detail how destructive of logical insight this quantificatious way of thinking is.10

Exercise

Consider the following “quantificatious" fragment of English. Assume our customary composition principles. As before, the basic domains are D (= D₀) and {0, 1}, but the denotations of lexical items are now different.
[Ann] = {Ann}
[Jacob] = {Jacob}
[everything] = D
[nothing] = Ø
[vanished] = \( \lambda X \in \text{Pow}(D) \cdot X \subseteq \{y \in D : y \text{ vanished}\} \).
[reappeared] = \( \lambda X \in \text{Pow}(D) \cdot X \subseteq \{y \in D : y \text{ reappeared}\} \).

(a) Discuss the adequacy of this proposal with respect to the following types of structures:

Are the correct truth-conditions predicted in each of the three cases? Can the proposal be extended to other quantifier phrases?

(b) Which of the following inferences are predicted to be valid, given a “quantificational” semantics of the sort presented above? Which predictions are correct?

(i) Ann vanished fast.
Ergo: Ann vanished.
(ii) Everything vanished fast.
Ergo: Everything vanished.
(iii) Nothing vanished fast.
Ergo: Nothing vanished.

For this exercise, add an unanalyzed predicate vanished fast to our fragment, and assume the following relationship:

\( \{x \in D : x \text{ vanished fast}\} \subseteq \{x \in D : x \text{ vanished}\} \).

Given this assumption, determine whether you are able to prove that the premise logically implies the conclusion in each of the above examples. Note that, strictly speaking, the premises and conclusions are phrase structure trees, not just strings of words.
(c) Consider now the following passage from Geach:

In a reputable textbook of modern logic I once came across a shocking specimen of quantificational thinking. Before presenting it to you, I must supply some background. In ordinary affairs we quite often need to talk about kinds of things that do not exist or about which we do not yet know whether they exist or not; and this applies to ordinary scientific discourse as well – I once saw a lengthy chemical treatise with the title "Nonexistent Compounds." Accordingly, logicians need to lay down rules for propositions with empty subject-terms. The convention generally adopted is that when the subject-term is empty, ostensibly contrary categorical propositions are taken to be both true; for example, if there are no dragons, "All dragons are blue" and "No dragons are blue" are both true. This convention may surprise you, but there is nothing really against it; there are other equally consistent conventions for construing such propositions, but no consistent convention can avoid some surprising and even startling results.

Now my author was trying to show the soundness of this convention, and to secure that, came out with the following argument. . . . "If there are no dragons, the phrases 'all dragons' and 'no dragons' both refer to one and the same class — a null or empty class. Therefore 'All dragons are blue' and 'No dragons are blue' say the same thing about the same class; so if one is true, the other is true. But if there are no dragons to be blue, 'No dragons are blue' is true; therefore, 'All dragons are blue' is also true." I know the argument sounds like bosh; but don't you be fooled – it is bosh. ¹²

Question: What is wrong with the logician's argument?

6.3 The solution: generalized quantifiers

6.3.1 "Something", "nothing", "everything"

We have seen that quantificational DPs like "something", "everything", and "nothing" are not proper names. Within our current framework, this means that their denotations are not in \( \text{Dc} \). That is, quantificational DPs are not of semantic type \( \text{e} \). We have also seen that the denotations of quantificational DPs like "something", "everything", and "nothing" are not sets of individuals. Hence they cannot be of semantic type \( \langle \text{e}, \text{t} \rangle \). What, then, is the semantic type of quantificational DPs? The Fregean reasoning that we have been following so far
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gives a clear answer, given (a) the usual kind of phrase structure, as in (1) below, (b) our previous assumptions about the denotations for VPs and sentences, and (c) our goal to reduce semantic composition to functional application whenever possible.

(1)

Our rules of semantic composition determine that the denotations of the VP-nodes in the above structures coincide with the denotation of vanish, hence are elements of $D_{<e,t>}$. Since the denotations of the sentences must be members of $D_{e}$, we now conclude that the semantic type of quantificational DPs must be $<e,t>,t>$, provided that the mode of composition is indeed functional application, and given that type $e$ is excluded.

Here is a more intuitive way of thinking about the semantic contribution of “nothing”, “everything”, and “something” in subject position. Rather than denoting an individual or a set of individuals, “nothing” in “nothing vanished” says something about the denotation of the predicate “vanished”. It states that there is no individual of which the predicate is true; that is, there is no individual that vanished. If we replace “nothing” with “everything”, the claim is that the predicate is true of all individuals. Finally, sticking “something” in the subject position of a predicate leads to the statement that there is at least one individual of which the predicate is true. Quantificational DPs, then, denote functions whose arguments are characteristic functions of sets, and whose values are truth-values. Such functions are sometimes called “second-order properties”, first-order properties being functions of type $<e,t>$. The second-order property $[\text{nothing}]$, for example, applies to the first-order property $[\text{vanished}]$, and yields truth just in case $[\text{vanished}]$ does not apply to any individual. More recently, second-order properties are commonly referred to as “generalized quantifiers”.

The reasoning we just went through leads to the following lexical entries for “nothing”, “everything”, and “something”:

$[\text{nothing}] = \lambda f \in D_{<e,t>} . \text{there is no } x \in D_{e} \text{ such that } f(x) = 1.$

$[\text{everything}] = \lambda f \in D_{<e,t>} . \text{for all } x \in D_{e}, f(x) = 1.$

$[\text{something}] = \lambda f \in D_{<e,t>} . \text{there is some } x \in D_{e} \text{ such that } f(x) = 1.$
Quantifiers

Our semantic composition rules guarantee the right modes of combination for subject DPs with their VPs without any further machinery. If the DP is a proper name, the denotation of the VP applies to the denotation of the DP. If the DP is a quantifier phrase, the denotation of the DP applies to the denotation of the VP. The difference between the two modes of semantic composition can be illustrated by trees with semantic type annotations as in (2):

(2)

```
S, t
  └── DP, <<e,t>,t> VP, <e,t>
    │      └── N       V
    │       nothing    vanished
    └── DP, e VP, <e,t>
        └── N       V
            Mary    vanished
```

In both structures, the rule for Non-Branching Nodes (NN) determines that the denotation of the lexical item occupying the subject position is passed up to the DP-node. The difference between the denotations of the lexical items “nothing” and “John”, then, brings about a difference in the way the respective DP-nodes are semantically combined with their VPs.

Exercise

Calculate the truth-conditions for the trees in (1) above.

6.3.2 Problems avoided

Before we adopt the new higher type for quantificational DPs, we want to be sure that it will indeed avoid the problems we noted for the simpler types e and <e,t>. Let us return to the inference patterns we looked at in section 6.1.1.

Subset to superset

When $[\alpha]$ is of type e, and $\beta$ necessarily denotes a subset of $\gamma$, then $\alpha \beta$ entails $\alpha \gamma$. But At most one letter came yesterday morning does not entail At most one letter came yesterday. So we concluded that $[[\text{at most one letter}]]$ cannot be in $D_e$. 
Now we must show that $D_{<<e,\tau>,\tau}$ contains a denotation that will predict the invalidity of this inference. That is, we must show that there exists a function $f \in D_{<<e,\tau>,\tau}$ such that it is possible that $f([\text{came yesterday morning}]) = 1$, but $f([\text{came yesterday}]) = 0$. This is quite evident. We only need to point out that it is possible for the actual facts to be such that $[\text{came yesterday morning}] \neq [\text{came yesterday}]$. There are as many different functions in $D_{<<e,\tau>,\tau}$ as there are ways of mapping the elements of $D_{e,\tau}$ to $\{0, 1\}$. So for each given pair of distinct elements of $D_{e,\tau}$, there are lots of functions in $D_{<<e,\tau>,\tau}$ that map the first to 1 and the second to 0.

**Law of Contradiction**

When $[\alpha]$ is of type $e$, and $\beta$ and $\gamma$ denote necessarily disjoint sets, then "\(\alpha \land \neg \alpha \lor \gamma\)" is a contradiction. But *More than two cats are indoors and more than two cats are outdoors* is not a contradiction. So $[\text{more than two cats}]$ cannot be in $D_e$.

Now we must show that $D_{<<e,\tau>,\tau}$ contains a denotation that will predict the possible truth of this conjunction. That is, we must show that there exists a function $f \in D_{<<e,\tau>,\tau}$ such that it is possible that $f([\text{indoors}]) = 1$ and $f([\text{outdoors}]) = 1$. Obviously, there exist plenty such functions.

**Law of Excluded Middle**

When $[\alpha]$ is of type $e$, and $\beta$ and $\gamma$ denote sets whose union is necessarily all of $D$, then "\(\alpha \lor \beta \land \neg \alpha \lor \gamma\)" is a tautology. But *Everybody here is over 30 or everybody here is under 40* is not a tautology. So $[\text{everybody}]$ here cannot be in $D_e$. The proof that a suitable denotation exists in $D_{<<e,\tau>,\tau}$ is similarly trivial.

In sum, once we allow DPs to have meanings of type $<<e,\tau>,\tau>$, the problematic predictions of a theory that would only allow type $e$ (or types $e$ and $<e,\tau>$) are no longer derived. This in itself is not a tremendously impressive accomplishment, of course. There are always many boring theories that avoid making bad predictions about logical relations by avoiding making any at all.

**Truth-conditional effects of syntactic reorganization**

It is more challenging to reflect on our second argument against the uniform type $e$ analysis of DPs. We observed that certain syntactic operations which systematically preserve truth-conditions when they affect names (for example, topicalization, passive) sometimes alter them when they affect quantifiers. Does the type $<<e,\tau>,\tau>$ analysis help us predict *this* fact, and if so, why?
For reasons that we will attend to in the next chapter, the English examples we used to illustrate this issue in section 6.1.2 involve complications that we cannot yet handle. For instance, we are not ready to treat both sentences in pairs like (3a) and (3b).

(3) (a) Everybody answered many questions correctly.
     (b) Many questions, everybody answered correctly.

But we can construct a hypothetical case that shows how it is possible in principle to alter truth-conditions by the topicalization of a phrase whose meaning is of type $<<e,t>,t>>$.

Suppose we topicalized the embedded subject $\alpha$ in a structure of the form (4a), yielding (4b).

(4) (a) It is not the case that $\alpha$ is asleep.
     (b) $\alpha$, it is not the case that $t$ is asleep.

Never mind that structures like (4b) are syntactically ill-formed. We only want to show here that (4a) and (4b) can differ in truth-conditions when $\llbracket \alpha \rrbracket \in D_{<<e,t>,t>>}$, which they couldn’t when $\llbracket \alpha \rrbracket \in D_e$.

Here are two possible structures (see section 6.1.2).
If \([\alpha] \in D_e\), the two structures are provably equivalent. (*Exercise: Give the proof.*) But what happens if \([\alpha] \in D_{<e,t>,t}>\)? We will show that in this case their truth-conditions may differ. Suppose, for instance, that \(\alpha = \text{everything}\), with the lexical entry given above. Then we calculate as follows:

\[
(5) \quad [(4a)] = 1 \\
\text{iff} \\
[\text{everything is asleep}] = 0 \\
\text{iff} \\
[\text{everything}](\text{[asleep]}) = 0 \\
\text{iff} \\
[\text{asleep}](x) = 0 \text{ for some } x \in D.
\]

\[
(6) \quad [(4b)] = 1 \\
\text{iff} \\
[\text{everything}](\text{[wh, it is not the case that } t_1 \text{ is asleep}]) = 1 \\
\text{iff} \\
[\text{wh, it is not the case that } t_1 \text{ is asleep}](x) = 1 \text{ for all } x \in D \\
\text{iff} \\
[\text{it is not the case that } t_1 \text{ is asleep}]^{[1 \rightarrow x]_1} = 1 \text{ for all } x \in D \\
\text{iff} \\
[\text{t_1 is asleep}]^{[1 \rightarrow x]_1} = 0 \text{ for all } x \in D \\
\text{iff} \\
[\text{asleep}](x) = 0 \text{ for all } x \in D.
\]

The last lines of (5) and (6) express clearly distinct conditions. It is easy to imagine that the actual facts might verify the former but falsify the latter: this happens whenever some, but not all, individuals sleep.

So we have shown that topicalization of a phrase with a type \(<<e,t>,t>\) meaning can affect truth-conditions. The example was a hypothetical one, but we will soon be able to analyze real-life examples of the same kind.

### 6.4 Quantifying determiners

Having decided on the denotations for quantifying DPs like “nothing”, “everything”, or “something”, we are now ready to find denotations for the quantifying determiners “no”, “every”, “some”, and what have you. Consider the following structure:
We reason as before. Assume that

(a) The phrase structures for phrases containing quantifying determiners are as given above.
(b) The semantic type of common nouns is \(<e,t>\).
(c) The semantic type of quantificational DPs is \(<<e,t>,t>\).
(d) Determiners and NPs semantically combine via functional application.

It now follows that the semantic type of quantifying determiners is \(<<e,t>,<<e,t>,t>>\). The annotated tree (2) illustrates the composition process. (2) has two binary-branching nodes, and in each case, the mode of composition is Functional Application.

As for the lexical entries, we have the following:

\[\text{[every]} = \lambda f \in D_\langle e,t,\rangle . [\lambda g \in D_\langle e,t,\rangle . \text{for all } x \in D_e \text{ such that } f(x) = 1, g(x) = 1] \]
\[\text{[no]} = \lambda f \in D_\langle e,t,\rangle . [\lambda g \in D_\langle e,t,\rangle . \text{there is } \text{no } x \in D_e \text{ such that } f(x) = 1 \text{ and } g(x) = 1] \]
\[\text{[some]} = \lambda f \in D_\langle e,t,\rangle . [\lambda g \in D_\langle e,t,\rangle . \text{there is some } x \in D_e \text{ such that } f(x) = 1 \text{ and } g(x) = 1] \]
Exercise

Calculate the truth-conditions for (1) step by step.

Our analysis of quantifying determiners with denotations of type $<e,t>,<e,t>,t>$, is essentially the one first proposed by David Lewis in "General Semantics." A similar proposal for the treatment of quantification in natural languages was independently made by Richard Montague, and was taken up in subsequent work by Cresswell, Barwise and Cooper, and Keenan and Stavi.

6.5 Quantifier meanings and relations between sets

6.5.1 A little history

The semantics for quantifying determiners that we arrived at in the last section is a variant of the oldest known view of quantification, the so-called relational view, which can be traced back to Aristotle. About Aristotelian logic (syllogistics), the dominant paradigm up to the nineteenth century, Dag Westerståhl says the following:

The syllogistics is basically a theory of inference patterns among quantified sentences. Here a quantified sentence has the form

(1) $QXY$

where $X$, $Y$ are universal terms (roughly 1-place predicates) and $Q$ is one of the quantifiers all, some, no, not all. In practice, Aristotle treated these quantifiers as relations between terms. Aristotle chose to study a particular type of inference pattern with sentences of the form (1), the syllogisms. A syllogism has two premisses, one conclusion, and three universal terms (variables). Every sentence has two different terms, all three terms occur in the premisses, and one term, the "middle" one, occurs in both premisses but not in the conclusion. It follows that the syllogisms can be grouped into four different "figures", according to the possible configurations of variables:

<table>
<thead>
<tr>
<th>$Q_1ZY$</th>
<th>$Q_1YZ$</th>
<th>$Q_1ZY$</th>
<th>$Q_1YZ$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$Q_2XZ$</td>
<td>$Q_2XZ$</td>
<td>$Q_2ZX$</td>
<td>$Q_2ZX$</td>
</tr>
<tr>
<td>$Q_3XY$</td>
<td>$Q_3XY$</td>
<td>$Q_3XY$</td>
<td>$Q_3XY$</td>
</tr>
</tbody>
</table>
Here the $Q_i$ can be chosen among the above quantifiers, so there are $4^4 = 256$ syllogisms.

Now the question Aristotle posed – and, in essence, completely answered – can be formulated as follows:

For what choices of quantifiers are the above figures valid?

For example, if we in the first figure let $Q_1 = Q_2 = Q_3 = \text{all}$, a valid syllogism results ("Barbara" in the medieval mnemonic); likewise if $Q_1 = Q_3 = \text{no}$ and $Q_2 = \text{all}" ("\text{Celarent}"). Note that Aristotle's notion of validity is essentially the modern one: a syllogism is valid if every instantiation of $X, Y, Z$ verifying the premisses also verifies the conclusion.\(^\text{16}\)

In their history of logic, William and Martha Kneale describe a contribution by Leibniz as popularized by the eighteenth-century mathematician L. Euler:\(^\text{17}\)

Leonhard Euler's "Lettres à une Princesse d'Allemagne" (which were written in 1761 and published in St. Petersburg in 1768) must be mentioned among works of the eighteenth century that contributed something to mathematical logic. Those letters which deal with logic contain no attempt to work out a calculus, though Euler was a great mathematician; but they popularized Leibniz's device of illustrating logical relations by geometrical analogies, and this had some influence on thinkers in the next century. In particular it directed attention to the extensional or class interpretation of general statements; for Euler represented (or rather illustrated) the four Aristotelian forms of statements by three relations of closed figures according to the following scheme . . .

Every a is b

No a is b

Some a is b

And Frege (though he may be more famous for his invention of predicate logic, with its 1-place quantifiers) also endorsed the relational view of quantifiers in various places: "the words all, every, no, some combine with concept words [our NPs or VPs]. In universal and particular affirmative and negative statements we express relations between concepts and indicate the specific nature of these relations by means of those words."\(^\text{18}\)
6.5.2 Relational and Schönfinkelized denotations for determiners

On the relational theory of quantification, quantifiers denote relations between sets. For instance, “every” denotes the subset relation, and “some” denotes the relation of non-disjointness. A sentence like “Every goat is a mutt” is understood as stating that the set of goats is a subset of the set of mutts. And a sentence like “Some goat is a mutt” is understood as stating that the set of all goats is not disjoint from the set of all mutts.

If we take “relation” in its exact mathematical sense (a set of ordered pairs), then our semantics for quantifying determiners from the previous section is not literally an instance of the relational theory. But there is a very straightforward connection between our determiner denotations and relations between sets.

Here is a sample of the relations that common quantifying determiners would express on a relational theory in the strictest sense of the word:

(1) For any \( A \subseteq D \) and any \( B \subseteq D \):

\[
\begin{align*}
(a) \quad & <A, B> \in R_{\text{every}} \text{ iff } A \subseteq B \\
(b) \quad & <A, B> \in R_{\text{some}} \text{ iff } A \cap B \neq \emptyset \\
(c) \quad & <A, B> \in R_{\text{no}} \text{ iff } A \cap B = \emptyset \\
(d) \quad & <A, B> \in R_{\text{at least two}} \text{ iff } |A \cap B| \geq 2^{19} \\
(e) \quad & <A, B> \in R_{\text{most}} \text{ iff } |A \cap B| > |A - B| 
\end{align*}
\]

It is clear why we are not assuming in this book that \([\text{every}] = R_{\text{every}}, [\text{some}] = R_{\text{some}}, \text{ etcetera.} \) This would not be consistent with our goal of minimizing the number of principles of semantic composition. (If, for example, \([\text{every}] \) literally were \( R_{\text{every}} \), we could not interpret \textit{Every painting vanished} by functional application, but would need special new rules.) Nevertheless, just as it is sometimes convenient to pretend that VPs denote sets (though they really denote functions), it can be convenient to talk as if determiners denoted relations between sets. We just have to understand clearly how such talk translates back into our official theory.

To see the connection between the relations defined in (1) above and our determiner meanings of type \( \langle \mathcal{E}, \mathcal{T} \rangle, \langle \mathcal{E}, \mathcal{T}, \mathcal{T} \rangle \), it is helpful to appreciate an important analogy between quantificational determiners and transitive verbs: Both can be viewed as expressing 2-place relations, the only difference being that the latter are first-order relations (they relate individuals), while the former are second-order relations (they relate sets of individuals or characteristic functions thereof). In the section on Schönfinkelization, we saw how to construct functions in \( D_{\langle \mathcal{E}, \mathcal{T} \rangle} \) out of first-order 2-place relations. By analogous operations, we can get from second-order 2-place relations to functions in \( D_{\langle \mathcal{E}, \mathcal{I}, \mathcal{T} \rangle} \).
To consider a concrete case, how do we get from $R_{\text{every}}$ to $\llbracket \text{every} \rrbracket$? Our starting point is (1a) (repeated from above).

(1a) $R_{\text{every}} = \{<A, B> \in \text{Pow}(D) \times \text{Pow}(D) : A \subseteq B\}^{20}$

The set $R_{\text{every}}$ has a characteristic function, which we'll call "$F_{\text{every}}$":

(2) $F_{\text{every}} = \lambda <A, B> \in \text{Pow}(D) \times \text{Pow}(D). A \subseteq B.$

$F_{\text{every}}$ is a 2-place function that maps pairs of sets of individuals into truth-values. It can now be Schönfinkeled in two ways. We Schönfinkel it from left to right and call the result "$f_{\text{every}}$":

(3) $f_{\text{every}} = \lambda A \in \text{Pow}(D). [\lambda B \in \text{Pow}(D). A \subseteq B].$

We are almost there now: $f_{\text{every}}$ just about could serve as the meaning of every if it weren't for the fact that NPs and VPs denote functions, not sets. To correct this, we must replace the sets $A$ and $B$ by functions $f$ and $g$, and accordingly rewrite the subset condition to the right of "iff".

(4) $\llbracket \text{every} \rrbracket = \lambda f \in D_{<c,t>} . [\lambda g \in D_{<c,t>} . (x \in D : f(x) = 1) \subseteq (x \in D : g(x) = 1)].$

This is the lexical entry for every as we determined it in section 6.4 in a slightly different, but equivalent, formulation. In definition (4), we just used a bit more set-theoretic notation for the formulation of the value description of the $\lambda$-term.

**Exercise**

Every one of $R_{\text{every}}$, $F_{\text{every}}$, $f_{\text{every}}$, and our $\llbracket \text{every} \rrbracket$ is in principle a candidate for the denotation of the English determiner every provided that one assumes suitable composition principles to fit one's choice. And these four are not the only possibilities. Due to the systematic relations between sets and their characteristic functions, and between 2-place functions and their various Schönfinkelizations, there are lots of additional variants. Most of them are of no particular interest and don't occur in the literature; others happen to be common. In the influential paper by Barwise and Cooper (see n. 13), determiners were treated as denoting functions from $\text{Pow}(D)$ into $\text{Pow}(\text{Pow}(D))$. Your task in this exercise is to spell out this option.
(a) Give examples of Barwise and Cooper-style lexical entries for a couple of run-of-the-mill determiners.

(b) Specify the composition rules that are needed in conjunction with these lexical entries.

(c) Show that there is a one-to-one correspondence between functions from $\text{Pow}(D)$ into $\text{Pow}(\text{Pow}(D))$ and relations between subsets of $D$.

6.6 Formal properties of relational determiner meanings

If quantifiers correspond to binary relations, we may investigate whether they do or do not have standard properties of binary relations like symmetry, reflexivity, and so on. Aristotle was already interested in questions of this kind. A substantial part of the *Prior Analytics* is dedicated to investigating whether the two terms $A$ and $B$ in a statement of the form $QAB$ are “convertible”; that is, whether the quantifier involved expresses a symmetric relation. By way of illustration, let $Q$ be the determiner “some”, and $A$ and $B$ the predicates “US citizen” and “native speaker of Spanish” respectively. Since “Some US citizens are native speakers of Spanish” and “Some native speakers of Spanish are US citizens” are logically equivalent (that is, they have the same truth-conditions), we have grounds to believe that “some” is convertible. It all depends, of course, on whether the logical equivalence is unaffected by the particular choice of predicates.

Here is a list of some potentially interesting mathematical properties of relational determiner denotations. In every definition, $\delta$ stands for a determiner, and $A$, $B$, $C$ range over subsets of the domain $D$.\textsuperscript{22}

<table>
<thead>
<tr>
<th>Definiendum</th>
<th>Definiens</th>
</tr>
</thead>
<tbody>
<tr>
<td>$\delta$ is reflexive</td>
<td>for all $A : \langle A, A \rangle \in R_{\delta}$</td>
</tr>
<tr>
<td>$\delta$ is irreflexive\textsuperscript{23}</td>
<td>for all $A : \langle A, A \rangle \not\in R_{\delta}$</td>
</tr>
<tr>
<td>$\delta$ is symmetric</td>
<td>for all $A, B : \text{if } \langle A, B \rangle \in R_{\delta}, \text{then } \langle B, A \rangle \in R_{\delta}$</td>
</tr>
<tr>
<td>$\delta$ is antisymmetric</td>
<td>for all $A, B : \text{if } \langle A, B \rangle \in R_{\delta} \text{ and } \langle B, A \rangle \in R_{\delta}, \text{then } A = B$</td>
</tr>
<tr>
<td>$\delta$ is transitive</td>
<td>for all $A, B, C : \text{if } \langle A, B \rangle \in R_{\delta} \text{ and } \langle B, C \rangle \in R_{\delta}, \text{then } \langle A, C \rangle \in R_{\delta}$</td>
</tr>
<tr>
<td>$\delta$ is conservative</td>
<td>for all $A, B : \langle A, B \rangle \in R_{\delta} \iff \langle A, A \cap B \rangle \in R_{\delta}$</td>
</tr>
<tr>
<td>$\delta$ is left upward monotone\textsuperscript{24}</td>
<td>for all $A, B, C : \text{if } A \subseteq B \text{ and } \langle A, C \rangle \in R_{\delta}, \text{then } \langle B, C \rangle \in R_{\delta}$</td>
</tr>
<tr>
<td>$\delta$ is left downward monotone</td>
<td>for all $A, B, C : \text{if } A \subseteq B \text{ and } \langle B, C \rangle \in R_{\delta}, \text{then } \langle A, C \rangle \in R_{\delta}$</td>
</tr>
</tbody>
</table>
δ is right upward monotone  
for all A, B, C: if A ⊆ B and <C, A> ∈ R_δ,  
then <C, B> ∈ R_δ  

δ is right downward monotone  
for all A, B, C: if A ⊆ B and <C, B> ∈ R_δ,  
then <C, A> ∈ R_δ

---

**Exercise**

For every property defined above, try to find determiners whose denotations have it, as well as determiners whose denotations don’t.

Why have linguists been interested in classifying determiner meanings according to such mathematical properties? Is this just a formal game, or does it throw some light on the workings of natural language? Modern research within the generalized quantifier tradition has shown that some of those mathematical properties may help formulate constraints for possible determiner meanings. Keenan and Stavi, for example, have proposed that all natural language determiners are conservative. In order to see what a non-conservative determiner would look like, imagine that English “only” was a determiner rather than an adverb. The non-equivalence of “only children cry” and “only children are children that cry” would now establish that “only” is a determiner that is not conservative. Conservativity, then, is a non-trivial potential semantic universal.

Other mathematical properties have been argued to pick out linguistically significant classes of DPs. The following two exercises will give you some taste of this influential line of research. When you work on those exercises, be warned that you may not be able to come up with completely satisfactory answers. Try your best, and note any open problems. If you want to delve deeper into those areas, consult the pioneering work of Milsark, Fauconnier, and Ladusaw, and the handbook articles mentioned in note 25 at the end of the chapter for further directions.

---

**Exercise on “there”-insertion**

It has often been observed that not all kinds of NPs are allowed in “there”-insertion constructions. Here are two examples:

(i) There are some apples in my pocket.  
(ii) *There is every apple in my pocket.
Test a number of quantifiers as to their behavior in "there"-insertion constructions, and try to characterize the class of quantifiers that are permitted in this environment with the help of some formal property of determiner denotations. Consider the mathematical properties defined above.

**Exercise on negative polarity**

The adverb "ever" is an example of a so-called *negative polarity item* (NPI), so called because it seems to require a negative environment:

(i) I haven’t ever visited the Big Bend National Park.
(ii) *I have ever visited the Big Bend National Park.

However, there needn’t always be a “not” to license “ever”. For instance, the following examples are also grammatical.

(iii) Very few people ever made it across the Cisnos range.
(iv) Every friend of mine who ever visited Big Bend loved it.

Try out other sentences like (iii) and like (iv) with different determiners in place of “very few” and “every”. Which property of these determiners seems to correlate with the distribution of “ever”? Again, consider the properties defined above.

### 6.7 Presuppositional quantifier phrases

Determiners denote functions of type \(<<e,t>,<<e,t>,t>>\). Total or partial functions? So far we have tacitly assumed the former. The lexical entries we have given for *every*, *some*, *no*, *more than two*, etcetera, all define total functions. They thus guarantee that *every* \(\alpha\), *some* \(\alpha\), *no* \(\alpha\), . . . always have a semantic value, regardless of the facts (provided that \(\alpha\) itself has a semantic value\(^2\)). In other words, quantifying determiners, as we have treated them so far, never give rise to presuppositions.

But we have no good reason to assume that this is generally correct for all quantifying determiners of natural languages. Indeed, there are some persuasive examples of determiners which seem to denote partial functions.
6.7.1 "Both" and "neither"

Consider the determiners both and neither, as in both cats, neither cat. What are the intuitive truth-conditions of a sentence like (1), for instance?

(1) Neither cat has stripes.

If there are exactly two cats and neither has stripes, (1) is clearly true. If there are exactly two cats and one or both of them have stripes, (1) is clearly false. But what if there aren’t exactly two cats? For example, suppose there is just one cat and it doesn’t have stripes. Or suppose there are three cats (all equally relevant and salient) and none of them has stripes. In such circumstances, we are reluctant to judge (1) either true or false; rather, it seems inappropriate in much the same way as an utterance of the cat when there isn’t a unique (relevant and salient) cat. This suggests the following lexical entry:

(2) \( f_{\text{neither}} = \lambda A : A \in \text{Pow}(D) \& |A| = 2 \& [\lambda B \in \text{Pow}(D) . A \cap B = \emptyset] \)

Together with the definition of “presupposition” in chapter 4, (2) predicts that (1) presupposes there to be exactly two cats.

Similar judgments apply to sentences like Both cats have stripes, suggesting an analogous lexical entry for both:

(3) \( f_{\text{both}} = \lambda A : A \in \text{Pow}(D) \& |A| = 2 \& [\lambda B \in \text{Pow}(D) . A \subseteq B] \).

Exercise

Give a precise characterization of the relation between [neither] and [no] and the relation between [both] and [every].

6.7.2 Presuppositionality and the relational theory

The existence of presuppositional determiners like neither and both is actually incompatible with a strictly relational theory of quantifiers. A relation, as you recall, is a set of ordered pairs. A given ordered pair is either an element of a given relation, or else it is not. There is no third possibility. Suppose, for instance, that \( R \) is some relation between sets of individuals; that is, \( R \) is some subset of \( \text{Pow}(D) \times \text{Pow}(D) \). Then for any arbitrary pair of sets \( <A, B> \in \text{Pow}(D) \times \text{Pow}(D) \), we either have \( <A, B> \in R \) or \( <A, B> \notin R \). The characteristic function of \( R \) is a total function with domain \( \text{Pow}(D) \times \text{Pow}(D) \), and any Schönfinkelizations of it are total functions with domain \( \text{Pow}(D) \).
This means that the procedure for constructing determiner meanings from relations which we gave in section 6.5.2 always produces total functions, hence non-presuppositional meanings. In practice, this means that we cannot obtain the entries for both and neither by this construction. Put differently, we cannot fully describe the meanings of these determiners by sets of ordered pairs of sets. If we try, for example, for both, the best we can come up with is (4).

\[ R_{\text{both}} = \{<A, B> \in \text{Pow}(D) \times \text{Pow}(D) : A \subseteq B \land |A| = 2\} \]

(4) correctly characterizes the conditions under which a sentence of the form both $\alpha \beta$ is true, but it fails to distinguish the conditions under which it is false from those where it has no truth-value. If we based our lexical entry for both on (4), we would therefore predict, for example, that (5) is true (!) in a situation in which there is only one cat and I saw it.

(5) I didn't see both cats.

This is undesirable. Speakers' actual judgments about (5) in this situation fit much better with the predictions of the presuppositional analysis we adopted earlier.

While a strictly relational theory cannot describe the meanings of both and neither, an almost relational theory, on which determiner meanings are potentially partial functions from $\text{Pow}(D) \times \text{Pow}(D)$ to $\{0, 1\}$, would work fine. On such a theory, both and no, for instance, would denote the functions defined in (6).

\[ \begin{align*}
(6a) & \quad F_{\text{both}} = \lambda<A, B> : A \subseteq D \land B \subseteq D \land |A| = 2 \land A \subseteq B. \\
(6b) & \quad F_{\text{no}} = \lambda<A, B> : A \subseteq D \land B \subseteq D \land A \cap B = \emptyset.
\end{align*} \]

In the case of a nonpresuppositional determiner $\delta$, $F_{\delta}$ happens to be the characteristic function of some relation on $\text{Pow}(D)$, but not when $\delta$ is presuppositional. In practice, the label "relational theory" is also applied to such an almost relational theory.

The fact that presuppositional determiners do not correspond to relations in the strictest sense gives rise to some indeterminacy as to how standard properties of relations apply to them. Consider, for example, the notion of irreflexivity. As a property of relations between subsets of $D$, it may be defined as in (7a) or (7b).

\[ R \text{ is irreflexive} \]

\[ \begin{align*}
(7a) & \quad \ldots \text{iff for all } A \subseteq D, <A, A> \notin R. \\
(7b) & \quad \ldots \text{iff for no } A \subseteq D, <A, A> \in R.
\end{align*} \]

(7a) is the definition we employed above; (7b) would have been a fully equivalent choice, and no less natural. Suppose we now ask ourselves whether neither is an irreflexive determiner. The initial answer that we come up with is that this
is not a well-defined question. The only definition that we have for "irreflexivity" as a property of determiners is the one in section 6.6. But this cannot be applied to \textbf{neither}, since there is no such thing as $R_{\text{neither}}$. We just learned that it is not possible to define such a relation.

We might stop right here and agree henceforth that the concept of irreflexivity is applicable only to nonpresuppositional quantifiers. But this may not be desirable if this concept plays some role in our semantic theory. For example, Barwise and Cooper (see note 13) have claimed that it plays a central role in the analysis of the English "there"-construction. If we are interested in this sort of claim, we have an incentive to look for a natural extension of the definition of irreflexivity that will allow it to apply to both total and partial determiner meanings.

So let's replace the "R" in (7) by an "F", which stands for a possibly partial function from $\text{Pow}(D) \times \text{Pow}(D)$ into $\{0, 1\}$. How should we rewrite the rest? There are, in principle, two possibilities:

\begin{enumerate}
\item[(8)] $F$ is irreflexive
\begin{enumerate}
\item[(a)] ... iff for all $A \subseteq D$, $F(A, A) = 0$.
\item[(b)] ... iff for no $A \subseteq D$, $F(A, A) = 1$.
\end{enumerate}
\end{enumerate}

When $F$ happens to be total – that is, when $F$ is the characteristic function of some relation – then (8a) and (8b) are equivalent, and they amount to exactly the same thing as (7). More precisely, if $R$ is irreflexive in the sense of (7), then $\text{char}_{R^{31}}$ is irreflexive in the sense of both (8a) and (8b): and if $\text{char}_{R}$ is irreflexive in the sense of either (8a) or (8b), then $R$ is irreflexive in the sense of (7). (Exercise: Prove this.) This being so, both (8a) and (8b) qualify as natural extensions of the basic concept of irreflexivity. However, the two are \textit{not equivalent!} They do coincide for total functions $F$, but they diverge for partial ones. To see this, consider the function $F_{\text{neither}}$:

\begin{enumerate}
\item[(9)] $F_{\text{neither}} = \lambda <A, B> : A \subseteq D \& B \subseteq D \& |A| = 2 \& A \cap B = \emptyset$.
\end{enumerate}

If we adopt (8b), $F_{\text{neither}}$ qualifies as irreflexive, but if we adopt (8a), it doesn't. (Exercise: Prove this.)

So we have to make a choice between (8a) and (8b). From a purely formal point of view, the choice is entirely arbitrary. The standard choice in the linguistic literature is (8b), for reasons that have to do with the intended empirical applications (see exercise below). (8b) is the more liberal definition of the two, in the sense that it makes it easier to qualify as irreflexive: Any $F$ that is irreflexive according to (8a) is also irreflexive under (8b), but not vice versa. This fact becomes more transparent if we reformulate (8b) as follows:

\begin{enumerate}
\item[(10)] $F$ is irreflexive
\begin{enumerate}
\item iff for all $A \subseteq D$ such that $<A, A> \in \text{dom}(F) : F(A, A) = 0$.
\end{enumerate}
\end{enumerate}
Exercise

Prove that (10) is equivalent to (8b).

Similar choices arise when other mathematical properties are extended from relations between sets to potentially partial determiner meanings. In many such cases, one of the *prima facie* reasonable definitions has been chosen as the standard definition in the linguistic literature. For instance, here are the official extended versions for selected concepts from our list in section 6.6. As before, "A", "B", and "C" range over subsets of D, and \( \delta \) stands for a determiner.

<table>
<thead>
<tr>
<th>Definiendum</th>
<th>Definiens</th>
</tr>
</thead>
<tbody>
<tr>
<td>( \delta ) is reflexive</td>
<td>for all A such that (&lt;A, A&gt; \in \text{dom}(F_\delta)) : (F_\delta(A, A) = 1)</td>
</tr>
<tr>
<td>( \delta ) is irreflexive</td>
<td>for all A such that (&lt;A, A&gt; \in \text{dom}(F_\delta)) : (F_\delta(A, A) = 0)</td>
</tr>
<tr>
<td>( \delta ) is symmetric</td>
<td>for all A, B : (&lt;A, B&gt; \in \text{dom}(F_\delta) ) iff (&lt;B, A&gt; \in \text{dom}(F_\delta)), and (F_\delta(A, B) = 1) iff (F_\delta(B, A) = 1)</td>
</tr>
<tr>
<td>( \delta ) is conservative</td>
<td>for all A, B : (&lt;A, B&gt; \in \text{dom}(F_\delta) ) iff (&lt;A, A \cap B&gt; \in \text{dom}(F_\delta)), and (F_\delta(A, B) = 1) iff (F_\delta(A, A \cap B) = 1)</td>
</tr>
<tr>
<td>( \delta ) is left upward monotone(^{32} )</td>
<td>for all A, B, C : if A ( \subseteq ) B, (F_\delta(A, C) = 1), and (&lt;B, C&gt; \in \text{dom}(F_\delta)), then (F_\delta(B, C) = 1)</td>
</tr>
<tr>
<td>( \delta ) is left downward monotone</td>
<td>for all A, B, C : if A ( \subseteq ) B, (F_\delta(B, C) = 1), and (&lt;A, C&gt; \in \text{dom}(F_\delta)), then (F_\delta(A, C) = 1)</td>
</tr>
<tr>
<td>( \delta ) is right upward monotone</td>
<td>for all A, B, C : if A ( \subseteq ) B, (F_\delta(C, A) = 1), and (&lt;C, B&gt; \in \text{dom}(F_\delta)), then (F_\delta(C, B) = 1)</td>
</tr>
<tr>
<td>( \delta ) is right downward monotone</td>
<td>for all A, B, C : if A ( \subseteq ) B, (F_\delta(C, B) = 1), and (&lt;C, A&gt; \in \text{dom}(F_\delta)), then (F_\delta(C, A) = 1).</td>
</tr>
</tbody>
</table>

Exercise

Go back to the exercises on "there"-sentences and negative polarity in section 6.6, and reconsider your answers in the light of the present discussion of presuppositional determiners.

6.7.3 Other examples of presupposing DPs

Our observations about the presupposition of both cats generalize to DPs of the form the two cats, the three cats, the four cats, etcetera. A sentence like
The twenty-five cats are in the kitchen is felt to presuppose that there are twenty-five (relevant) cats, and to assert that all (relevant) cats are in the kitchen. Our semantics should thus predict that, for any numeral \( n \) and any NP \( \alpha \):

\[
(11) \quad \text{the } n \ \alpha \ \text{has a semantic value only if } \|\alpha\| = n. \\
\quad \text{Where defined, } \|\text{the } n \ \alpha\| = \lambda A. \ [\alpha] \subseteq A.
\]

How do we ensure this prediction by appropriate lexical entries in conjunction with the usual composition principles?

The answer to this question depends on what we take to be the syntactic structure of DPs of this form. If we could assume that their constituent structure is [DP [D the n] NP], and could treat the n as a lexical unit, it would be straightforward. We would only have to write lexical entries like those in (12), then.

\[
(12) \quad \begin{array}{l}
F_{\text{the three}} = \lambda<A, B> : |A| = 3. \ A \subseteq B \\
F_{\text{the four}} = \lambda<A, B> : |A| = 4. \ A \subseteq B \\
e\text{tc.}
\end{array}
\]

This analysis is quite clearly not right. The definite article and the numeral are evidently lexical items in their own right, and we would like to explain how their separate meanings contribute to the meaning of the DP as a whole. Moreover, the constituent structure of such DPs seems to be [the [n NP]] rather than [[the n] NP]. Unfortunately, we cannot attend to these facts, since we will not be able to go into the semantics for plural NPs in this book. Well aware of its limitations, then, we will assume the \textit{ad hoc} analysis in (12).

**Exercise**

Suppose we treated \textbf{the cat} and \textbf{the cats} as elliptical surface variants of \textbf{the one cat} and \textbf{the at least two cats} respectively. How does this treatment compare to the semantics for singular definites that we gave in chapter 4?

Another type of DP that gives rise to presuppositions is the so-called partitive construction, exemplified by \textit{two of the five cats}, \textit{none of the (at least two) dogs}, etcetera. Assuming again a very \textit{ad hoc} syntactic analysis, we could capture their presuppositions by lexical entries for the complex determiners along the lines of (13).

\[
(13) \quad F_{\text{none of the (at least two)}} = \lambda<A, B> : |A| \geq 2. \ A \cap B = \emptyset
\]
Exercise

Return to the exercises on there-sentences and NPIs in section 6.6. What predictions do the solutions you proposed imply for DPs of the forms considered in this section, e.g., the seven cats, none of the cats?

6.8 Presuppositional quantifier phrases: controversial cases

In the previous section, we considered some types of DPs which carry presuppositions about the cardinality of their restrictors. For most of them, we did not provide a serious compositional analysis to pinpoint the exact lexical sources of these presuppositions. But whatever their sources may turn out to be, the fact that the DPs as a whole carry the presuppositions in question is rather apparent and uncontroversial.

With these observations in the background, we now turn to a family of much more controversial claims about the presuppositions of quantified structures in natural languages.

6.8.1 Strawson’s reconstruction of Aristotelian logic

The meanings we have been assuming throughout this chapter for sentences with every, some (a), and no are the ones that were promoted by the founders of modern logic. We repeat once more the definitions of the relevant second-order relations:

\[
\begin{align*}
(1) \quad (a) \quad R_{\text{every}} &= \{<A, B> : A \subseteq B\} \\
(b) \quad R_{\text{some}} &= \{<A, B> : A \cap B \neq \emptyset\} \\
(c) \quad R_{\text{no}} &= \{<A, B> : A \cap B = \emptyset\}
\end{align*}
\]

As has been noted by many commentators on the history of logic, the definitions in (1) are not consistent with certain assumptions about the semantics for every, some, no that were part of (at least some versions of) Aristotelian logic. In many works in this tradition, generalizations such as the ones in (2) were considered valid.
(2) For any predicates \( \alpha, \beta \):

(i) every \( \alpha \beta \) and no \( \alpha \beta \) is a contradiction.

(ii) some \( \alpha \beta \) or some \( \alpha \text{ not } \beta \) is a tautology.

(iii) every \( \alpha \beta \) entails some \( \alpha \beta \).

(iv) no \( \alpha \beta \) entails some \( \alpha \text{ not } \beta \).

We give a concrete English instance for every of the sentence schemata is (2(i))–(2(iv)):

(3) (a) Every first-year student in this class did well and no first-year student in this class did well.

(b) Some cousin of mine smokes, or some cousin of mine doesn’t smoke.

(c) Every professor at the meeting was against the proposal.

\[ \therefore \text{ Some professor at the meeting was against the proposal.} \]

(d) No student presentation today was longer than an hour.

\[ \therefore \text{ Some student presentation today wasn’t longer than an hour.} \]

According to the Aristotelian laws in (2), (3a) is contradictory, (3b) is a tautology, and (3c) and (3d) are valid inferences. The predictions of the modern (“classical”) analysis in (1), are otherwise: (a) is contingent; specifically, it is true when there are no first-year students in this class. (b) is likewise contingent; it is false when the speaker has no cousins. The premise in (c) doesn’t entail the conclusion: the former is true but the latter false when there were no professors at the meeting. Likewise for (d): its premise is true, but its conclusion false if there were no student presentations today.

Which of the two sets of predictions is correct for English? As we will see, the empirical evidence bearing on this question is surprisingly difficult to assess. Before we take a closer look at it, let’s present a concrete semantic analysis of the determiners every and some that predicts the validity of the Aristotelian laws in (2). Such an analysis is suggested in the following passage from Strawson:

Suppose someone says “All John’s children are asleep”. Obviously he will not normally, or properly, say this, unless he believes that John has children (who are asleep). But suppose he is mistaken. Suppose John has no children. Then is it true or false that all John’s children are asleep? Either answer would seem to be misleading. But we are not compelled to give either answer. We can, and normally should, say that, since John has no children, the question does not arise. . . .

. . . The more realistic view seems to be that the existence of children of John’s is a necessary precondition not merely of the truth of what is said, but of its being either true or false. . . .
... What I am proposing, then, is this. There are many ordinary sentences beginning with such phrases as "All . . .", "All the . . .", "No . . .", "None of the . . .", "Some . . .", "Some of the . . .", "At least one . . .", "At least one of the . . ." which exhibit, in their standard employment, parallel characteristics to those I have just described in the case of a representative "All . . ." sentence. That is to say, the existence of members of the subject-class is to be regarded as presupposed (in the special sense described) by statements made by the use of these sentences; to be regarded as a necessary condition, not of the truth simply, but of the truth or falsity, of such statements. I am proposing that the four Aristotelian forms [i.e., "all $\alpha \beta$", "no $\alpha \beta$", "some $\alpha \beta$", "some $\alpha$ not $\beta$", as they are interpreted in Aristotelian logic] should be interpreted as forms of statements of this kind. Will the adoption of this proposal protect the system from the charge of being inconsistent when interpreted? Obviously it will. For every case of invalidity, of breakdown in the laws [of Aristotelian logic], arose from the non-existence of members of some subject-class being incompatible with either the truth or the falsity of some statement of one of the four forms. So our proposal, which makes the non-existence of members of the subject-class incompatible with either the truth or the falsity of any statement of these forms, will cure all these troubles at one stroke. We are to imagine that every logical rule in the system, when expressed in terms of truth and falsity, is preceded by the phrase "Assuming that the statements concerned are either true or false, then . . ." Thus . . . the rule that [all $\alpha P$] entails [some $\alpha \beta$] states that, if corresponding statements of these forms have truth-values, then if the statement of the [form all $\alpha \beta$] is true, the statement of the [form some $\alpha \beta$] must be true; and so on.

At the beginning of this quote, it is not immediately apparent which set of English determiners Strawson means to make claims about. His primary example all John's children is probably among those less controversial candidates for a presuppositional analysis which we already surveyed in the previous section. Possessive DPs like John's children are standardly analyzed as covert definite descriptions (with structures essentially of the form the children (of) John), and all + definite is taken to be a partitive. (So all John's children is a surface variant of all of John's children, the optionality of of being an idiosyncratic property of all.) As far as this particular example goes, then, Strawson's claims may not go beyond what is commonplace in contemporary formal semantics, and the same goes for none of the, some of the, and at least one of the, which he lists a few paragraphs down. But in this same list he also includes plain all, no, some, and at least one. So it is clear that he means his proposal to extend beyond the partitives, and we concentrate here on its application to the simple determiners he mentions, plus every (which we assume Strawson would not distinguish from all in any respect relevant here).
According to Strawson then, at least some occurrences of English every, some, etcetera behave as if their lexical entries were not as in (1) above, but rather as below:

\[(4)\]
\[(a)\] \(F_{\text{every}} = \lambda A, B : A \neq \emptyset . A \subseteq B.\)
\[(b)\] \(F_{\text{no}} = \lambda A, B : A \neq \emptyset . A \cap B = \emptyset.\)
\[(c)\] \(F_{\text{some}} = \lambda A, B : A \neq \emptyset . A \cap B \neq \emptyset.\)

These entries validate the Aristotelian generalizations in (2), provided, as Strawson notes, that there are suitable definitions of the basic semantic properties. Here is a proposal:

**Basic semantic properties**

\(\phi\) is a *tautology* iff the semantic rules alone establish that, if \(\phi\) is in the domain of \([\_]\), then \([\phi]\) = 1.

\(\phi\) is a *contradiction* iff the semantic rules alone establish that, if \(\phi\) is in the domain of \([\_]\), then \([\phi]\) = 0.

\(\phi\) *entails* \(\psi\) iff the semantic rules alone establish that, if \(\phi\) and \(\psi\) are both in the domain of \([\_]\) and \([\phi]\) = 1, then \([\psi]\) = 1.

So it is reasonable to hypothesize that, to the extent that native speakers' judgments about examples like (3) conform to the predictions in (2), this is due to their using the entries in (4) rather than those in (1).

### 6.8.2 Are all determiners presuppositional?

Strawson was not directly engaged in natural language semantics, and it is impossible to attribute to him any very specific claim in that domain. He did say that at least some uses of simple English determiners carried existence presuppositions, but this certainly doesn't imply that (4a), (4b), and (4c) are the lexical entries for English every, no, some. At best it implies that some English utterances are correctly translated into a symbolic language whose determiners have the semantics of (4). As regards English itself, this leaves many possibilities open. Perhaps (4) represents certain readings, among others, of lexically ambiguous items of English. Or perhaps it's not the lexical denotations of the determiners at all that are responsible for the relevant presuppositions, but other ingredients of the structures in which they occur. Linguists inspired by Strawson's discussion have explored various concrete options in this regard. We will begin here with a hypothesis that is simpler and more radical than most of the other proposals that are out there. By considering some standard objections to it, we will develop an appreciation for a variety of hypotheses that are currently under debate.
Following up on Strawson's remarks, let us look at the following Presuppositionality Hypothesis, versions of which have been argued for by James McCawley, and more recently by Molly Diesing.38

(5) **Presuppositionality Hypothesis**

In natural languages, all lexical items with denotations of type $<e,t>,<e,t>,t>$ are presuppositional, in the sense of the following mathematical definition (where $\delta$ is a lexical item of the appropriate semantic type, such as a determiner):

$\delta$ is presuppositional iff for all $A \subseteq D$, $B \subseteq D$ : if $A = \emptyset$, then $<A, B> \notin \text{dom}(F_{\delta})$.

According to this hypothesis, determiner denotations like those defined in (1) are not possible in natural languages at all, and *a fortiori* cannot be the denotations of English *every*, *all*, *some*, *at least one*, *a*, *no* on any of their readings. The closest allowable denotations are the minimally different presuppositional ones defined in (4). Similarly, the denotations of most other determiners that we have considered in this chapter must be slightly different from what we have assumed. The following would be some revised lexical entries that conform to the Presuppositionality Hypothesis:

(6) (a) $F_{\text{few}} = \lambda A, B : A \neq \emptyset \land |A \cap B|$ is small.
   (b) $F_{\text{most}} = \lambda A, B : A \neq \emptyset \land |A \cap B| > \frac{1}{2}|A|$.
   (c) $F_{\text{at least three}} = \lambda A, B : A \neq \emptyset \land |A \cap B| \geq 3$.
   (d) $F_{\text{at most three}} = \lambda A, B : A \neq \emptyset \land |A \cap B| \leq 3$.

How does this hypothesis fare with respect to the linguistic facts? There are some observations that appear to support it, as even its harshest critics concede. Lappin and Reinhart,39 for instance, report that their informants judge 7(a) below to be a presupposition failure rather than a true statement. (The informants were all aware, of course, that America has never had kings, and made their judgment on the basis of this piece of factual knowledge.) McCawley40 already reported similar intuitions about 7(b):

(7) (a) All/every American king(s) lived in New York.
   (b) All unicorns have accounts at the Chase Manhattan Bank.

These judgments are predicted by the lexical entry for *every* in (4) (and a parallel entry for *all*), whereas the one in (1) would predict judgments of "true". So we have observations here which, *ceteris paribus*, favor (4) over (1), and thus support the Presuppositionality Hypothesis. But for other examples, the predictions seem not to be borne out. Here are some more data from Reinhart:41
(8) (a) No American king lived in New York.
    (b) Two American kings lived in New York.

Regarding (8a) and (8b), only about half of Reinhart's informants judged them presupposition failures on a par with (7). The other half judged (8a) true and (8b) false. So this looks as if some people employed the presuppositional entries in (4) and (6), whereas others employed the standard ones from (1) and previous sections. Reinhart also contrasts (7) with (9).

(9) (a) Every unicorn has exactly one horn.
    (b) Every unicorn is a unicorn.

Her informants judged (9a) and (9b) true without hesitation, as if in this case, unlike with (7), they employed the standard nonpresuppositional entry of every.

So the evidence seems to be mixed. Can we detect some systematic pattern? Lappin and Reinhart (following earlier authors, especially Barwise and Cooper and de Jong and Verkuyl) endorse two descriptive generalizations. The relevant difference between (7) and (9), they maintain, is that (9) need not be taken as a description of the actual world, whereas (7) cannot naturally be taken any other way. They formulate this first generalization roughly as follows:

(10) In non-contingent contexts, speakers' judgments about presupposition failure and truth-value conform to the standard (nonpresuppositional) analyses of determiners.

The notion of a "non-contingent context", of course, cries out for further clarification, and we will return to this shortly.

If (10) succeeds in distinguishing (7) from (9), the difference between (7) and (8) must lie elsewhere, both presumably being understood as "contingent" statements in the relevant sense. Apparently, what is decisive here is the choice of determiner. (7) and (8) are completely alike up to their determiners, and by testing a larger sample of additional determiners in the same kind of context, it emerges that the dividing line is the same as that which determines the grammaticality of there be sentences. Determiners disallowed in the there construction (every, almost, every, not every, most, and, unsurprisingly, uncontroversially presuppositional ones like both and neither) give rise to presupposition failure judgments in examples like (7). The ones that are alright in there sentences (no, numerals, few, many), when placed in the same sentence frame, evoke the mixed reactions found with (8). If (following Milsark) we define a "strong" determiner as one that is barred from there sentences, and a "weak" determiner as one that's allowed there, we can state Lappin and Reinhart's second descriptive generalization as follows:
In contingent contexts, strong determiners evoke judgments that conform to the presuppositional analysis, whereas weak determiners give rise to mixed judgments that conform sometimes to the presuppositional and sometimes to the standard analysis.

We will assume that both of these empirical generalizations are at least roughly on the right track. Either one of them seems to threaten the Presuppositional Hypothesis by outlining a set of *prima facie* counterexamples. Suppose we nevertheless wanted to defend this hypothesis. Let us take a closer look, first at (10) and then at (11), to see whether such a defense is possible and what it would commit us to.

### 6.8.3 Nonextensional interpretation

The fact that (9a) and (9b) are spontaneously judged true is unexpected if every carries a Strawsonian existence presupposition. Similar facts were acknowledged by previous proponents of a presuppositional semantics for every: in particular, Diesing, de Jong and Verkuyl, and Strawson himself. Let's look at some of their examples and how they responded to the challenge. We quote from de Jong and Verkuyl:

> ... we claim that the standard interpretation of universal quantification is not based upon the most regular use of *all* in natural language, but rather upon the marked use of this expression: its conditional use. Due to its conditional structure,

\[(12) \forall x (Px \rightarrow Qx)\]

expresses a specific relation between P and Q. Such an interpretation of

\[(13) \text{All ravens are black}\]

is favored by the fact that the set of all ravens is a subset of the set of black entities, *which is not based on observation, but on induction or hypothesis. Blackness is taken as a property inherent to ravens*, as long as no counterexample shows up. Examples such as (13) must be treated as marked cases in comparison with *contingent sentences* such as (14) and (15).

\[(14) \text{All seats are taken.}\]
\[(15) \text{All men are ill.}\]

We use the term “marked” here in the linguistic sense. Sentence (13) is a clear example of a statement having *the status of a law* – or a hypothesis,
an opinion or a belief – that is firmly settled in science, in biological theory and also in our everyday naive physics. . . . In general, sentences like (14) and (15) are fully nonsensical if there are no men or seats in the context of use. This seems to be due to the fact that there is no inherent relation between seats and the property “to be taken”, or between men and the property “being ill”. As a consequence (12) cannot serve as a basis for this interpretation of (14) and (15). However, suppose that (ultrafeminist) science discovers that (15) is a law of nature. In that case, the interpretation of (15) is on a par with (13). So it depends on whether a certain sentence functions as a lawlike statement in a theory (or a more or less consistent set of everyday assumptions), when the conditional use gets the upper hand. . . . We regard the lawlike use of sentences as marked because we do not think it is a property of natural language that there are theories, whether scientific or embedded in our everyday opinions. Summarizing, all can be used in lawlike sentences as well as in contingent statements. Both contexts impose different interpretations on all. Only in hypothetical contexts can all be interpreted without presuppositions on the size of $[N]$. In contingent statements the use of all requires a non-empty noun denotation.

De Jong and Verkuyl and Lappin and Reinhart, though they come down on opposite sides about the two competing lexical entries for all (every), agree essentially on the terms of the debate. Not only do they give similar descriptions of the relevant intuitions, they also assume that, first, the judgments about (9a), (9b), and (13) are incompatible with a (unambiguously) presuppositional analysis of every/all, and that, second, they support the standard nonpresuppositional analysis. But are these two assumptions so evidently correct? Strawson, it turns out, argued long ago that the second was quite mistaken:44

. . . There are, in fact, many differences among general sentences. Some of these differences have been exploited in support of the claim that there are at least some general sentences to which the negatively existential analysis (“$(x)(fx \supset gx)$”) is applicable.43 For example, it may be said that every one of the following sentences viz.,

(16) All twenty-side rectilinear plane figures have the sum of their angles equal to $2 \times 18$ right angles

(17) All trespassers on this land will be prosecuted

(18) All moving bodies not acted upon by external forces continue in a state of uniform motion in a straight line
might be truthfully uttered; but in no case is it a necessary condition of their truthful utterance that their subject-class should have members. Nor can it be said that the question of whether or not they are truthfully uttered is one that arises only if their subject-class has members. . . . These facts, however, are very inadequate to support the proposed analysis. If the proposed analysis were correct, it would be a sufficient condition of the truthful utterance of these sentences that their subject-classes had no members; for \( \neg(\exists x)(fx) \) entails \( (x)(fx \supset gx) \). But this is very far from being the case for these, or for any other, general sentences.

Let us consider this important point more carefully. If Strawson is right, it was no more than a coincidence that the “true”-judgments which Reinhart found with (9a) and (9b) conformed to the standard analysis of every. If we vary the predicate, we will find just as many “false”-judgments:

\[
(9) \begin{align*}
(c) & \text{ Every unicorn has exactly two horns.} \\
(d) & \text{ Every unicorn fails to be a unicorn.}^{46}
\end{align*}
\]

Similarly, if we make suitable alterations in the predicates of Strawson’s examples, their intuitive truth-values go from true to false:

\[
(18') \text{ All moving bodies not acted upon by external forces decelerate at a rate of } 2.5 \text{ m/sec}^2.
\]

The same knowledge of physics that makes us assent to (18), regardless of whether we believe that there actually exist any bodies not acted upon by external forces, will make us dissent from (18').

The same point can be made by minimally varying the quantifier, say from all/every to no:

\[
(9) \begin{align*}
(e) & \text{ No unicorn has exactly one horn.} \\
(18'') & \text{ No moving body not acted upon by external forces continues in a state of uniform motion in a straight line.}
\end{align*}
\]

If the “true”-judgments on (9a) and (18) are claimed as evidence for the standard analysis of every/all, shouldn’t the “false”-judgments on (9e) and (18'') be counted as evidence against the standard analysis of no?

Summarizing Strawson’s point, the judgments that speakers have about truth and falsity of lawlike quantificational statements do not support the standard analysis of quantifiers any more than the presuppositional one. Both analyses
Quantifiers seem to make systematically wrong predictions in this domain, albeit different ones. The presuppositional analysis errs by predicting that emptiness of the restrictor’s extension suffices to render all such statements truth-value-less. The standard analysis errs by predicting that emptiness of the restrictor’s extension suffices to verify all of them. It looks like both analyses miss the point of what really determines the intuitive truth-values of these statements.

We might leave it at this and simply set the semantics of lawlike quantificational sentences aside. What data we have looked at concerning their truth-conditions turned out to be simply irrelevant to the question we set out to answer: namely, whether natural language exhibits a universal constraint against non-presuppositional determiners. But this is not entirely satisfactory. If none of the semantic treatments of quantifiers that we are currently entertaining is applicable to lawlike statements, then shouldn’t we abandon them all, instead of wasting our time comparing them to each other?47

Fortunately, we can do better. Diesing maintains her version of the Presuppositionality Hypothesis in spite of the apparently conflicting data in connection with lawlike statements that she is well aware of. Here is a possible story that is in the spirit of Diesing’s reaction to the challenge.48

So far, we have only seen negative characterizations of what people do when they decide that, say, (9a) is true and (9c) is false. They don’t, we saw, treat the fact that there are no actual unicorns as evidence one way or the other. What do they treat as evidence, then? It’s not really so hard to give at least the rough outlines of a positive answer. There is a certain body of mythology that we have acquired together with the word unicorn. This mythology specifies a set of possible worlds in which there exist unicorns and in which these unicorns have certain properties and not others. All unicorns in the worlds that instantiate the myth have one horn, for instance, and none of them in any of these worlds have two horns. This, it seems, is the intuitive reason why (9a) is true and (9c) is false. If we consider not the set of actual unicorns (which is empty), but rather the set of mythologically possible unicorns (which consists of all the unicorns in the possible worlds which are truly described by the relevant myth), then it turns out that this set is (i) nonempty, (ii) a subset of the set of possible individuals with exactly one horn, and (iii) disjoint from the set of possible individuals with two horns.

This account suggests that the interpretation of (9a) and (9c) does, after all, involve a run-of-the-mill interpretation of every: namely, the one in (1) or (4). Either makes the correct prediction that (9a) is true and (9c) false, once we assume that the quantifier does not quantify over actual unicorns, but rather over mythologically possible unicorns. We do not have to commit ourselves to a particular technical realization of this idea here. Diesing assumes that lawlike statements are implicitly modalized. Quantificational DPs in lawlike statements, then, would be under the scope of a modal operator, and this is why quantification
is over possible individuals. Suppose some such analysis is on the right track. It appears, then, that what's special about the class of statements that Lappin and Reinhart call "non-contingent" and de Jong and Verkuyl "lawlike" is not a special interpretation of the quantifier at all. Rather, it is a special property of the environment the quantifier finds itself in.

Analogous stories can be told about de Jong and Verkuyl's biological law (13) or Strawson's sentences (16)-(18). The point about (18), we might say, is that it quantifies over all physically possible bodies not acted on by external forces, not just the actual ones. And (17) quantifies over all possible trespassers in the worlds that conform to the intentions of the land-owner at least as well as any world with trespassers in it can conform to them. In every one of these examples, we make our truth-value judgments by considering such sets of (wholly or partially) non-actual individuals.

At this point, it looks as if the truth-conditions of lawlike statements such as (9a)-(9e), (13), (16)-(18), and (18'), (18'') are, after all, consistent with both the presuppositional or the standard analysis of every/all and no. The existence presuppositions of the former are automatically fulfilled once the restrictor is interpreted in a suitably large domain of possible individuals, and then both analyses make the same, correct, predictions. It seems, then, that the data about lawlike quantificational statements that we have considered so far have no bearing one way or the other on the decision between presuppositional and standard lexical entries for the quantifiers. But this time around, the conclusion does not come packaged with reasons to suspect that both are wrong. Rather, both remain in the running, and we are ready to continue our search for decisive evidence.

Diesing reports an argument due to Kratzer that intends to show that the presuppositional analysis is not just compatible with, but even necessary for, an adequate treatment of quantifying determiners in lawlike and other modalized statements. The bottom line of the argument is that unless we adopt a presuppositional analysis, quantifying determiners within the scope of modal operators might give rise to the Samaritan Paradox, a paradox that is well known to scholars of modal logic and conditionals. Although we cannot go into the technical details of the argument here, we think an informal sketch may be useful, even in the absence of the necessary background from modal logic.

The Samaritan Paradox comes up with sentences like (19a)-(19c):

(19)  (a) The town regulations require that there be no trespassing.
      (b) The town regulations require that all trespassers be fined.
      (c) The town regulations require that no trespassers be fined.

Suppose we live in a world in which (19a) and (19b) are true, but (19c) is false. Intuitively, there is nothing wrong with such an assumption. Yet any theory
that combines the standard modal analysis for verbs like “require” with the nonpresuppositional analysis of “all trespassers” and “no trespassers” would predict otherwise. (19a) says that there is no trespassing in any possible worlds that are compatible with the actual town regulations (these are the possible worlds in which no actual town regulation is violated). If “all trespassers” and “no trespassers” are nonpresuppositional, (19b) and (19c) would both be true. (19b) is true just in case all trespassers are fined in all possible worlds that conform to the actual town regulations. Since there is no trespassing in any of those worlds, there are no trespassers, and (19b) comes out true. And so does (19c). Consequently, we can’t draw the desired distinction between (19b) and (19c).

The way we seem to understand (19b) and (19c) is that we temporarily suspend the regulation that there be no trespassing. But how come we suspend this regulation? Because we are temporarily assuming that there are trespassers. Where does this assumption come from? If “all trespassers” and “no trespassers” are presuppositional, we have an answer. The presupposition that there are trespassers could play a systematic role in picking out the set of possible worlds we are considering. For this to be an acceptable answer, however, we have to give a good account of the difference between the nonpresuppositional “no trespassing” in (19a) and the presuppositional “no trespassers” in (19c). The next section will look into this issue.

6.8.4 Nonpresuppositional behavior in weak determiners

We now turn to the challenge that the Presuppositionality Hypothesis faces from the behavior of weak determiners. Recall Reinhart and Lappin’s generalization (11):

(11) In contingent contexts, strong determiners evoke judgments that conform to the presuppositional analysis, whereas weak determiners give rise to mixed judgments that conform sometimes to the presuppositional and sometimes to the standard analysis.

The salience of the standard (nonpresuppositional) interpretation is known to be affected by pragmatic and grammatical factors, and in some examples it is entirely natural. Consider these:

(20) (a) No phonologists with psychology degrees applied for our job.
(b) Two UFOs landed in my backyard yesterday.
(c) At most, twenty local calls from this number were recorded.
These sentences might very well be used in conversations where speaker and hearer are fully conscious of the possibility that there may be nothing that satisfies the restrictor. For instance, if I believe that there just aren't any phonologists with psychology degrees and am trying to convince you of this, I might use (20a) to cite circumstantial evidence for my conviction. If it later turns out that indeed there are no phonologists with psychology degrees, I will not feel any pressure to rephrase my statement. Quite to the contrary, I may then reiterate: “That’s what I thought: there aren’t any phonologists with psychology degrees. No wonder that none applied.”

The situation with (20b) is a bit different. Unlike the previous sentence, this one cannot be used by a sincere speaker who believes that the restrictor is empty. After all, if there are no UFOs, then (20b) cannot be true. But it seems quite clear intuitively that it can be false in this case. Imagine that you and I have an ongoing disagreement about the existence of UFOs: I believe they exist, you do not. (20b) could be uttered in the course of one of our arguments about this matter: I might use it as evidence for my position. If you want to defend yours then, you will argue that (20b) is false. So in this case as well, the non-emptiness of the restrictor seems not to be a presupposition of (20b).

The case of (20c) is more similar to that of (20a), in that the non-existence of any local calls from this number would make it true (rather than false). Imagine I just got the phone bill and there is no extra charge for local calls. According to my contract with the phone company, the first twenty calls every month are covered by the basic rate. If I utter (20c) in this situation, I will not be taken as prejudging the question of whether there were any local calls from this number at all.

In every one of these cases, the standard, nonpresuppositional entries for the determiners seems to fit our intuitions much better than their presuppositional alternatives, and the Presuppositionality Hypothesis might accordingly look undesirable. This conclusion is further reinforced when we broaden the scope of our examples to include quantifiers in positions other than subject position.53 An especially natural environment for nonpresuppositional interpretations is the post-copular position of there sentences. Reinhart,54 for instance, reports that even those informants who perceived (8a) and (8b) as presupposition failures had no hesitation in judging (21a) and (21b) as respectively true and false.

(8) (a) No American king lived in New York.
(b) Two American kings lived in New York.

(21) (a) There were no American kings in New York.
(b) There were two American kings in New York.
In this particular construction, strong determiners are not grammatical in the first place. But in non-subject positions where both weak and strong determiners can occur, we see quite clearly that there is indeed a minimal contrast between weak and strong determiners. Strong determiners receive presuppositional readings regardless of position, whereas weak ones need not. Zucchi cites the following paradigm from Lumsden:

\[(22) \quad \text{If you find every mistake(s), I'll give you a fine reward.}\]

\begin{itemize}
  \item most
  \item many
  \item a
  \item no
  \item three
\end{itemize}

The examples with the strong determiners (every, most) convey that the speaker assumes there to be mistakes, whereas the ones with weak determiners (many, a, no, less than 2) sound neutral in this regard.

What do we conclude from this (cursory) survey of data? An obvious possibility is to maintain the Presuppositionality Hypothesis in its radical form and explore the possibility that weak determiners might be affected by a type ambiguity, as considered by Partee.\(^56\) In chapter 4, we looked at predicative uses of indefinites like \textit{a cat} in \textit{Julius is a cat}, and concluded that when so used, indefinites are of semantic type \(<e, t>\). Maybe all nonpresuppositional uses of weak DPs can be assimilated to the predicative use. This is the line taken by Diesing,\(^57\) who builds on insights from Discourse Representation Theory.\(^58\) Weak DPs may or may not be interpreted as generalized quantifiers. Only if they are, are they presuppositional. If weak DPs are ambiguous, we now need to say something about the observed distribution of possible readings. Diesing invokes a special principle, her Mapping Hypothesis, to this effect.

Much current research explores ways of doing away with principles like the Mapping Principle. De Hoop\(^59\) follows Diesing in assuming a type ambiguity for weak DPs, but links the different semantic types to different cases (in the syntactic sense). She then proposes to derive the distribution of the two types of weak DPs from syntactic principles governing the two types of cases, plus independent principles governing the interpretation of topic and focus. Other recent attempts in the literature reject the idea that weak DPs are ambiguous as to their semantic type, and argue that topicality and focus all by themselves are able to account for the presuppositionality facts.\(^60\) As things stand, there is no consensus yet.
Notes

1 Remember that we are assuming (following Abney) that phrases like "the mouse" are DPs (determiner phrases) headed by a determiner whose sister node is an NP. Since proper names, pronouns, and traces behave syntactically like phrases headed by an overt determiner, they are classified as DPs as well. Our DPs correspond to Montague's "term phrases". Our NPs correspond to Montague's "common noun phrases".

2 In this chapter, we will indulge in a lot of set talk that you should understand as a sloppy substitute for the function talk that we would need to use to be fully accurate.

3 Russell wouldn't disagree here, we presume, even as he criticizes Meinong for his "failure of that feeling for reality which ought to be preserved even in the most abstract studies," and continues: "Logic, I should maintain, must no more admit a unicorn than zoology can" ("Descriptions," in A. P. Martinich (ed.), The Philosophy of Language, 2nd edn (New York and Oxford, Oxford University Press, 1990), pp. 212–18, at p. 213). Such pronouncements can sound dogmatic out of context, but they don't, of course, carry the burden of Russell's argument.

4 Once again, we pretend that VPs denote sets when they really denote the corresponding characteristic functions.

5 The Law of Contradiction states the validity of "not (p and not p)".

6 The Law of Excluded Middle states the validity of "p or not p".


8 This statement would have to be qualified if we considered DPs that contain variables that are free in them. In that case, movement may affect variable binding relations, and thus affect truth-conditions, even if the moving phrase is of type e. We may safely disregard this qualification here, since it seems obvious that the DPs under consideration ("at least one question", "more than one policeman", etc.) have no free variables in them.

9 Days are, of course, included among the elements of D. They are more abstract entities than chairs and policemen, but objects nonetheless.


11 Recall that Pow(D), the power set of D, is defined as \{X : X ⊆ D\}.

12 Geach, Logic Matters, pp. 58ff.


14 D. Lewis, "General Semantics," in D. Davidson and G. Harman (eds), Semantics of Natural Languages (Dordrecht, Reidel, 1972), pp. 169–218. Lewis has intensional denotations instead of our extensional ones, but this has little impact on the treatment of quantifier phrases. Lewis's paper was first presented at the third La Jolla Conference on Linguistic Theory in 1969.

15 R. Montague, Formal Philosophy, ed. R. M. Thomason (New Haven, Yale University Press, 1974); M. J. Cresswell, Logics and Languages (London, Methuen, 1973); Barwise and Cooper, "Generalized Quantifiers"; E. L. Keenan and J. Stavi,


18 This quote is from G. Frege, "Über Begriff und Gegenstand" ["On concept and object"], *Vierteljahresschrift für wissenschaftliche Philosophie*, 16 (1892); our emphasis. See also *Grundgesetze der Arithmetik*.

19 For any set A, |A| is the cardinality of A, i.e., the number of members of A.

20 For any set A, Pow(A) (the power set of A) is the set of all subsets of A.

21 For any sets A and B, A × B (the Cartesian Product of A and B) is defined as \{<x, y> : x ∈ A and y ∈ B\}.

22 These definitions don't directly mention the quantifier denotation [0], but only the corresponding relation R_φ. Given the previous section, however, it is a routine exercise to translate every definiens into equivalent conditions on F_φ, f_φ, and [0].

23 Note that "irreflexive" doesn't mean "nonreflexive".

24 For this and the following three monotonicity properties, a number of other terms are also in common use. "Upward monotone" is also called "upward-entailing" or "monotone increasing", and "downward monotone" is called "downward-entailing" or "monotone decreasing". "Left upward monotone" is called "persistent", and "left-downward monotone", "anti-persistent". Authors who use these terms "persistent" and "anti-persistent" tend to use "monotone increasing" or "upward monotone" in the narrower sense of "right upward monotone" (and similarly for "monotone decreasing" and "downward monotone").


26 Keenan and Stavi, "Semantic Characterization." Conservativity is Barwise and Cooper's live-on property. Barwise and Cooper, "Generalized Quantifiers" conjecture that every natural language has conservative determiners.


28 If the restrictor happens to lack a semantic value, as in every woman on the escalator in South College, then of course the whole DP doesn't get one either. This is evidently not due to any lexical property of the *determiner* (here *every*).

29 We are not talking here about "both" and "neither" when they occur as part of the discontinuous coordinators "both ... and ...", "neither ... nor ..." (as in "Both John and Bill left", "Neither John nor Bill left"). These are not determiners and do not interest us here.
Strictly speaking, the lexical entry has to define \( \text{[neither]} \), of course. But you know how to get \( \text{[neither]} \) from \( f_{\text{neither}} \).

Recall that \( \text{char}_R \) is the characteristic function of \( R \).

Regarding the four monotonicity concepts, there is less of a consensus about how to extend them to partial determiners. For instance, an alternative to the present definition of "left upward monotone" would be the following:

\[
\text{left upward monotone} \quad \text{for all } A, B, C : \text{ if } A \subseteq B \text{ and } F_{\delta}(A, C) = 1, \text{ then }<B, C> \in \text{dom}(F_{\delta}) \text{ and } F_{\delta}(B, C) = 1
\]

Analogous alternative definitions may be entertained for the other three monotonicity properties. As you may discover in the exercise below, it is not completely evident which choice is preferable.


Our primary source here is P. F. Strawson, *Introduction to Logical Theory* (London, Methuen, 1952), who in turn refers to Miller, *The Structure of Aristotelian Logic*. According to Horn, however, Aristotle's own writings are not explicit enough to endorse all the "laws of the traditional system" as listed by Strawson (*Introduction*, pp. 156–63), and Aristotle's commentators have fleshed out the system in partially disagreeing directions. See L. Horn, *A Natural History of Negation* (Chicago, University of Chicago Press, 1989), esp. ch. 1, sect. 1.1.3: "Existential Import and the Square of Opposition," pp. 23–30. We do not want to get into questions of exegesis and history here. Strawson's proposal is interesting to us in its own right, even if (as Horn argues) it contradicts explicit statements by Aristotle and the majority of his medieval followers.

Don't be confused by the fact that "the classical analysis" refers to the modern one, not to any of those that date from antiquity. This is the customary terminology in the contemporary literature. The same perspective is reflected in the labels "standard" or "orthodox", which also mean the modern analysis.


By "subject-class", Strawson means the set denoted by the restrictor (i.e., \( \alpha \) in the schemata in (2)).

J. D. McCawley, "A Program for Logic," in Davidson and Harman (eds), *Semantics*, pp. 498–544; M. Diesing, "The Syntactic Roots of Semantic Partition" (Ph.D. dissertation, University of Massachusetts, Amherst, 1990); *idem, Indefinites* (Cambridge, Mass., MIT Press, 1992). McCawley associates all quantifying determiners other than "any" with existential presuppositions. Diesing ("Syntactic Roots" and *Indefinites*) connects the presuppositionality of quantifying DPs to the presence of a syntactically represented restrictive clause and the ability to undergo the syntactic operation of Quantifier Raising (see our chapters 7 and 8 for a discussion of quantifier raising). In M. Diesing and E. Jelinek, "Distributing Arguments," *Natural Language Semantics, 3/2* (1996), pp. 123–76, the connection with type theory and generalized quantifiers is made explicit. Combining the two works by Diesing with
that by Diesing and Jelinek, we conclude, then, that we are justified in ascribing to
Diesing the view that all DPs that express generalized quantifiers are presuppositional
in the sense under discussion here. Be this as it may, the arguments for or against
the presuppositional hypothesis that we are about to discuss apply directly to
Diesing’s proposals concerning the presuppositionality of quantifying DPs, even under

39 S. Lappin and T. Reinhart, “Presuppositional Effects of Strong Determiners: A Process­
ing Account,” Linguistics, 26 (1988), pp. 1022–37; T. Reinhart, Interface Strategies,
OTS Working Paper TL-95-002 (Utrecht University, 1995), ch. 4: “Topics and the
Conceptual Interface.”
40 McCawley, “Program for Logic.”
41 Reinhart, Interface Strategies.
42 Barwise and Cooper, “Generalized Quantifiers”; F. de Jong and H. Verkuyl, “Gen­
eralized Quantifiers: The Properness of their Strength,” in J. van Benthem and A. ter
Meulen, Generalized Quantifiers in Natural Language (Dordrecht, Foris, 1985),
pp. 21–43.
43 De Jong and Verkuyl, “Generalized Quantifiers,” pp. 29f., emphasis added and
eamples renumbered.
Reinhart, Interface Strategies, cites the examples from this passage.
45 Strawson’s label “negatively existential” for the standard analysis of every derives
from the fact that “∀x (fx → gx)” is equivalent to “¬∃x (fx & ¬gx)”. Notice also
that Strawson uses the so-called Principia Mathematica notation for predicate logic,
where “(x)” is the symbol for “∀x”, and “→” for “→”.
46 Lappin and Reinhart, “Presuppositional Effects,” and Reinhart, Interface Strategies,
actually cite (i):

(i) Every unicorn is not a unicorn.

They report that their informants judge (i) false and imply that this judgment con­
forms to the predictions of the standard analysis. But this sentence is predicted false
by the standard analysis only if the negation takes scope over the subject quantifier.
On the other scope-order, it is predicted true. In (i), the latter scope-order may be
excluded or dispreferred for independent reasons (see A. S. Kroch, “The Semantics
of Scope in English” (Ph.D. dissertation, MIT, 1974). Our (9d) is constructed so as
to avoid the potential scope ambiguity.
47 We might avoid this conclusion if we agreed with de Jong and Verkuyl that the
lawlike use of quantified sentences is somehow “marked”, i.e., beyond the proper
domain of linguistic theory. But this does not seem right to us (and here we agree
with Lappin and Reinhart). It is true that it is not “a property of natural language
that there are theories,” but – more to the point – it does appear to be a property
of natural language that theories can be expressed in it.
48 Diesing, Indefinites, pp. 95ff.
49 The argument was developed by Kratzer in class lectures. See Diesing, Indefinites,
p. 96.
50 For an overview and further references consult L. Aqvist, “Deontic Logic,” in Gabbay
and Guenthner, Handbook, pp. 605–714. See also A. Kratzer, “Modality,” in von
Stechow and Wunderlich, Semantik, pp. 639–50.
51 See also our chapter 12, where an intensional semantics is introduced, and the
semantics of attitude verbs is discussed.
The observations we are sketching here are found in a large number of works, many of which build on the work of Milwarck. See Milwarck, "Existential Sentences in English"; and idem, "Toward an Explanation of Certain Peculiarities of the Existential Construction in English," Linguistic Analysis, 3/1 (1977), pp. 1–29. (Linguistic Analysis misspells his name as "Milwarck.")

The following discussion includes examples with quantifying DPs in object position that we cannot yet interpret in a compositional way. Quantifying DPs in object position are the topic of the next chapter. We are confident, however, that the point of the examples can still be appreciated.

Reinhart, Interface Strategies.


Diesing, "Syntactic Roots"; idem, Indefinites.


7 Quantification and Grammar

7.1 The problem of quantifiers in object position

Almost all instances of quantifying DPs that we have looked at so far were in subject position. And there was a good reason for it. Compare (1a) to (1b):

(1) (a) [dp every linguist] [vp offended John].
(b) John [vp offended [dp every linguist]].

(1a) is true just in case the set of linguists is included in the set of those who offended John. We have seen how to arrive at the correct truth-conditions in a compositional way. The determiner denotation relates two sets. The first set (the restrictor set) is provided by the common noun “linguist”, the second by the VP “offended John”. But what if “every linguist” occurs in object position as in (1b)? Shouldn’t we assume that “every” still denotes a relation between sets? But then, which two sets? The restrictor set is the set of linguists, and it is provided by the common noun as before. The second set should be the set of all those who were offended by John. But this set is not denoted by any constituent in (1b). This is, in a nutshell, the problem of quantifiers in object position.

The dilemma becomes more dramatic if we consider sentences with multiple quantifier phrases:

(2) Some publisher offended every linguist.

(2) has two readings. On one reading, the claim is that there is at least one publisher who offended every linguist. The other reading is compatible with a situation where every linguist was offended by a possibly different publisher. Set theory lets us express the two readings:

(2') (a) \{x : x is a publisher\} \cap \{x \in \{y : y is a linguist\} \subseteq \{z : z offended x\}\} \neq \emptyset
(b) \{x : x is a linguist\} \subseteq \{x \in \{y : y is a publisher\} \cap \{z : z offended x\} \neq \emptyset\}. 
But how can we compute such statements in a compositional way from plausible syntactic structures?

From the perspective of our type-theoretic framework, the problem of quantifiers in object position presents itself as a type mismatch. What happens if we try to interpret (1b), for example? Recall that \([	ext{every}]\) is of type \(\langle e, t \rangle, \langle e, t \rangle, t \rangle\), and \(\text{linguist}\) of type \(\langle e, t \rangle\). These combine by Functional Application (FA) to yield a denotation of type \(\langle e, t \rangle, t \rangle\) for the quantifier phrase. \(\text{[offend]}\) is of type \(\langle e, \langle e, t \rangle \rangle\). Unfortunately, denotations of this type do not combine with those of \(\langle e, t \rangle, t \rangle\), either as function and argument or as argument and function. So FA yields no value for the VP, nor does any other principle apply. We are stuck.

We mentioned in the last chapter that the relational theory of quantification that we have adopted here is the oldest known theory of quantification, dating back at least to Aristotle. The problem of quantifiers in object position is almost as old. Medieval scholars tried to solve it, but failed, and so did many logicians and mathematicians in more modern times. A solution was eventually found by Frege. Frege discovered the notation of quantifiers and variables, and thereby “resolved, for the first time in the whole history of logic, the problem which had foiled the most penetrating minds that had given their attention to the subject.”

Modern linguistic theories fall into different camps, depending on their approach to the problem of quantifiers in object position. There are those who assume in the spirit of Frege that sentences are constructed in stages, and that at some stage, the argument positions of predicates might be occupied by traces or pronouns that are related to quantifier phrases via a syntactic relationship. The relationship is movement (Quantifier Lowering or Quantifier Raising) in Generative Semantics and in Chomsky’s Extended Standard Theory and its offspring, and the operation of “Quantifying in” in Montague Grammar. Other semanticists avoid displacement of quantifier phrases, and try to interpret all arguments of predicates \textit{in situ}. The displacement of quantifier phrases may be simulated in the semantics by storing their denotation in a so-called Cooper Store, or the flexibility of type theory may be used to overcome type mismatches. Variable-free versions of predicate logic (Combinatory Logic) led to variable-free versions of natural language semantics, as in the work of Szabolcsi, Steedman, Cresswell, and Jacobson. In what follows, we have picked an example of an \textit{in situ} approach and a particular instantiation of a movement approach for further discussion and comparison.

7.2 Repairing the type mismatch \textit{in situ}

We will now consider a way of overcoming the problem of quantifier phrases in object position by leaving the quantifier phrase in place.
7.2.1 An example of a "flexible types" approach

On the "flexible types" approach,\(^6\) we try to solve the problem of quantifier phrases in object position by optionally changing the semantic type of the quantifier phrase or of the verb. We will illustrate the first possibility here. You are invited to try out the second on your own.\(^7\)

Let quantifier phrases be multiply ambiguous. For 1-place and 2-place predicates as arguments we have the following two entries\(^8\) (extending this approach to any n-place predicate is straightforward – see exercise below):

\[
\begin{align*}
\text{[everybody}_1\text{]} &= \lambda f \in D_{<c,t>} . \text{for all persons } x \in D, f(x) = 1. \\
\text{[everybody}_2\text{]} &= \lambda f \in D_{<c,c,t>} . [\lambda x \in D . \text{for all persons } y \in D, f(y)(x) = 1]. \\
\text{[somebody}_1\text{]} &= \lambda f \in D_{<c,t>} . \text{there is some person } x \in D \text{ such that } f(x) = 1. \\
\text{[somebody}_2\text{]} &= \lambda f \in D_{<e,c,t>} . [\lambda x \in D . \text{there is some person } x \in D \text{ such that } f(y)(x) = 1].
\end{align*}
\]

For an example, consider the phrase structure tree (1):

\[
\begin{align*}
\text{[[ everybody}_1 \ [v_p \text{ offended somebody}_2\text{]]]} &= 1 \\
\text{iff} \\
\text{[[ everybody}_1\text{]]} (\text{[[ v_p \text{ offended somebody}_2\text{]]]}(x) &= 1 \\
\text{iff} \\
\text{for all persons } x, \text{[[ v_p \text{ offended somebody}_2\text{]]]}(x) &= 1
\end{align*}
\]

The truth-conditions of (1) can be calculated as follows:

\[
\begin{align*}
\text{[[ everybody}_1 \ [v_p \text{ offended somebody}_2\text{]]]} &= 1 \\
\text{iff} \\
\text{[[ everybody}_1\text{]]} (\text{[[ v_p \text{ offended somebody}_2\text{]]]}(x) &= 1 \\
\text{iff} \\
\text{for all persons } x, \text{[[ v_p \text{ offended somebody}_2\text{]]]}(x) &= 1
\end{align*}
\]
for all persons \( x \), \( \text{[somebody}_2\text{][[offended]]}(x) = 1 \)

iff

for all persons \( x \), there is some person \( y \), such that \( \text{[offended]}(y)(x) = 1 \)

iff

for all persons \( x \), there is some person \( y \), such that \( x \) offended \( y \).

This proposal implies that all English quantifier phrases are multiply ambiguous. They all have multiple syntactic representations. Since the ambiguity is systematic, it's a benign case of ambiguity. On this proposal, the syntax of English doesn't have to specify that, say, \( \text{everybody}_1 \) can only occur in subject position, and \( \text{everybody}_2 \) can only occur in object position. As far as the syntax goes, any quantifier phrase with any subscript can occur in any position. If a quantifier phrase appears in the wrong position, its mother node is not interpretable, and the structure is automatically ruled out.

Exercise 1

Specify the semantic rule for \( \text{nobody}_2 \), and calculate the truth-conditions for the phrase structure tree corresponding to \( \text{somebody}_1 \) greets \( \text{nobody}_2 \).

Exercise 2

Design a new entry for “everybody” (\( \text{everybody}_3 \)) that makes it possible to interpret phrase structure trees of the following kind:

```
  S
 / \  
Ann VP
 /   
V   PP
|     |
introduced  everybody  to Maria
```

Up to now in this section, we have limited ourselves to quantifying DPs that consist of one word, like “everybody”, “something”. But of course we must deal
with the full range of quantifying DPs of the form [Det NP]. These are not listed in the lexicon, so if they are ambiguous, their ambiguity must be traceable to a lexical ambiguity in one of their constituents. The natural place to locate this ambiguity is in the determiner.

For instance, the DP "a linguist" ought to have an alternate meaning of type \( \langle e, et \rangle, et \rangle \) so that we can use it in object position in place of "somebody" in (2). If we keep the familiar meaning for linguist of type \( \langle e, t \rangle \), this means that the determiner must be of type \( \langle et, \langle e, et \rangle, et \rangle \) in this case. Specifically, we need the following entry:

\[
[2] (a_2) = \lambda f \in D_{\langle e, t \rangle} \cdot [\lambda g \in D_{\langle e, et, et \rangle} \cdot [\lambda x \in D \cdot \text{for some } y \in D, f(y) = 1 \text{ and } g(y)(x) = 1]]
\]

Similarly for other determiners.

If every English determiner is multiply ambiguous in this way, then we would obviously be missing a generalization if we simply listed each reading for each determiner in the lexicon. Speakers of English presumably need not learn each reading separately: once they know what a determiner means in a subject DP, they can infer its meaning in object position without further evidence. So they must know some general rule by which the alternate readings of an arbitrary determiner are predictable from its basic type \( \langle et, \langle et, et \rangle \rangle \) meaning. In other words, there has to be a \textit{lexical rule} like the following:

\[
(3) \text{For every lexical item } \delta_1 \text{ with a meaning of type } \langle et, \langle et, et \rangle \rangle, \text{ there is a (homophonous and syntactically identical) item } \delta_2 \text{ with the following meaning of type } \langle et, \langle e, et \rangle, \langle e, t \rangle \rangle:
[\delta_2] = \lambda f \in D_{\langle e, et \rangle} \cdot [\lambda g \in D_{\langle e, et, et \rangle} \cdot [\lambda x \in D \cdot [\delta_1](f)(\lambda z \in D \cdot g(z)(x))]].
\]

(3) automatically provides an alternate meaning for any arbitrary determiner. The only determiner entries we need to list individually are those of the simplest type \( \langle et, \langle et, t \rangle \rangle \).

7.2.2 \textbf{Excursion: flexible types for connectives}

Flexible types were first proposed not for quantifiers but for connectives like "and" and "or". These seem to be able to coordinate phrases of a variety of different syntactic categories, with meanings of a corresponding variety of semantic types. For example:

\[
(4) [s_s \text{ John stays at home}] \text{ and } [s_s \text{ Mary works}].
\]

\[
(5) \text{ Ann will be } [r_{pp} \text{ in the garden}] \text{ or } [r_{pp} \text{ on the porch}].
\]
Bill \([v, \text{writes}] \text{ and } [v, \text{reads}]. \text{ Portuguese.}\)

A few books] or [dp, à lot of articles]] will be read.

Suppose the structure of the coordinate phrase in each example is as follows:

\[
\begin{array}{c}
X \\
\Delta \\
\{ \text{and} \} \\
X \\
\{ \text{or} \} \\
\Delta
\end{array}
\]

where \(X = S, \text{ PP, V, or DP}\)

Suppose further that \text{and} and \text{or} have their familiar meanings from propositional logic or, more accurately, appropriate Schönfinkelizations thereof:

\[
\begin{align*}
\text{[and]} &= \begin{bmatrix}
1 & 1 \\
0 & 0 \\
1 & 0 \\
0 & 0
\end{bmatrix} \\
\text{[or]} &= \begin{bmatrix}
1 & 1 \\
0 & 1 \\
1 & 0 \\
0 & 0
\end{bmatrix}
\end{align*}
\]

Sentence (4), where \(X = S\), is then straightforwardly interpretable. But (5) and (6) are not, because the denotations of PP and V in these examples are of type \(<e, t>\) and \(<e, <e, t>\rangle\) respectively, therefore unsuitable to combine with \text{[and]} or \text{[or]}.

A natural response is to posit a systematic ambiguity: There is not just one word “and”, but a family of related words \text{and}_1, \text{and}_2, \text{and}_3, \text{etcetera} (and likewise for “or”). The most basic members of each family, \text{and}_1 and \text{or}_1, are of type \(<t, t, t>\rangle\) and have the meanings defined in (9). The next ones are of type \(<e, t, t, t>\rangle\; \text{for instance, } \text{or}_2 \text{ is interpreted as follows.}

\[
\begin{align*}
\text{[or}_2\text{]} &= \lambda f \in D_{<e, t>} \cdot [\lambda g \in D_{<e, t>} \cdot [\lambda x \in D \cdot ([\text{or}_1](f(x))(g(x)))]].
\end{align*}
\]

This is the homonym of “or” that we must employ in (5) to make the sentence interpretable. Try it out to verify that it predicts the appropriate truth-conditions! Homonyms of yet more complicated types are needed for (6) and (7).

---

**Exercise 1**

Define appropriate meanings for \text{and}_3 and \text{or}_4 to be used in (6) and (7).
Exercise 2

What we have said here about the 2-place connectives "and" and "or" carries over to some extent to the 1-place connective negation. But the picture is less clear here because the word not has a much more limited syntactic distribution. Among the following examples parallel to (4)—(7) above, only (iv) is grammatical.

(i) *[s Not [s John stays at home]].
   *Not does John stay at home.

(ii) ?Ann will be [pp not [pp on the porch]].

(iii) *Bill [v not [v reads]] Portuguese.

(iv) [op Not [op everything]] will be read.

The most common position for not in English is none of the above, but rather as in (v).

(v) John doesn't stay at home.

What denotations for not do we need to interpret (iv) and (v)?

The point of this excursion into the semantics of "and" and "or" was to lend plausibility to the view that flexible types are a systematic phenomenon in natural language, and not just an ad hoc device we need for the treatment of quantifiers.

7.3 Repairing the type mismatch by movement

We have briefly looked at one way of assigning intuitively correct interpretations to VPs which contain quantificational objects. On that approach, the objects were left in place. We will now pursue in more depth an approach which maintains our original assumption that the determiner is unambiguous and can only be combined with two 1-place predicates. Since the overt syntactic structure of "John offended every linguist"
John VP
offended DP
every linguist

does not contain two such predicates, it follows that this structure cannot be the input to the semantic component. Rather, this sentence must have another structural description under which “every” combines with two constituents each having a denotation of type \(<e,t>\). Such a structure can be created by *moving* the DP “every linguist”.

To have a concrete proposal to work with, suppose here (and in the chapters to come) that we have a model of grammar like the inverted Y model of the Revised Extended Standard Theory and Government Binding Theory:

According to this model, semantic interpretation applies to representations on a level of Logical Form (LF), which is transformationally derived from S(urface) Structure (SS). The DP “every linguist” in (1) will move out of its VP and adjoin to S in the derivation from SS to LF. This movement operation, then, might feed semantic interpretation, but not necessarily phonetic realization, and might therefore be invisible. Like all movement operations, it leaves a trace. So the structure created by this movement has at least the following ingredients:

(2)
What we see in (2) is not quite enough yet to make an interpretable structure. Traces, on our assumptions from chapter 5, must bear an index to be interpretable. And since in a complete sentence every variable must be bound, the index on the trace has to be matched by an index on a variable binder somewhere. We propose that (3) below is a more complete and accurate representation of the structure created by moving “every linguist”.

(3)

```
(3)  S
    /\  
   DP  1  S
  /     |
 every linguist  John  VP
   |
  offended t1
```

The indexed trace here is straightforwardly interpreted by our Traces and Pronouns Rule as a variable. The adjoined index right below the moved phrase is supposed to be the variable binder. This requires a slight generalization of our Predicate Abstraction Rule, which currently only covers variable binders of the form “such_i”, “wh_i”, “who_i” etcetera. The fact that the rule had to mention particular lexical items was undesirable to begin with, since it went against the spirit of type-driven interpretation that does not permit composition principles that mention lexical items. When you look at the Predicate Abstraction Rule a bit more closely, however, you will notice that all it needs to see is the index on the relative pronouns, not their particular shapes. Let’s therefore adopt the following revision of the Predicate Abstraction Rule, with the understanding that “such” and the relative pronouns are to count as semantically vacuous items.

(4)  **Predicate Abstraction Rule (PA)**

Let \( \alpha \) be a branching node with daughters \( \beta \) and \( \gamma \), where \( \beta \) dominates only a numerical index i. Then, for any variable assignment \( a \), \( \langle \alpha \rangle^a = \lambda x \in D . \langle \gamma \rangle^a_{x^i} \).

The interpretation of structure (3) is now straightforward. Inside the VP, the transitive V of type \(<e, <e, t>\) composes by FA with the trace of type e, yielding a VP meaning of type \(<e, t>\). This composes (again by FA) with the subject's meaning, here of type e, to yield a type t meaning for the lower S. Concretely, the interpretation obtained for this lower S-node is this:

For any \( a \): \[ \langle \text{John offended } t_1 \rangle^a = 1 \text{ iff John offended } a(1). \]
At the next higher node, PA applies, and yields the following meaning of type \(<e,t>\).

\[
\lambda x \in D. \text{John offended } x.
\]

This is a suitable argument for the quantifying DP’s meaning of type \(<<e,t>,t>\), so FA can apply to the top S-node, and we obtain:

\[[(3)] = 1 \text{ iff for every } x \text{ such that } x \text{ is a linguist, John offended } x.\]

We have obtained the correct result, and we have been able to derive it here without resorting to a type-shifted homonym of “every”. We used exactly the denotation for “every linguist” that this DP has in subject position. There was no type mismatch at any point in the tree, thanks to the movement operation that applied to the object. This operation effected two crucial changes: it provided the transitive verb with an argument of type \(e\), and the moved quantifier phrase with an argument of type \(<e,t>\).

So far, so good. But the structure in (3) is not exactly what syntacticians imagine to be the output of movement. When we introduced the idea of moving the object in (2), you probably expected a representation like (5) rather than (3).

The difference between (3) and (5) is that in (5) the higher index forms a constituent with the moved phrase, whereas in (3) it forms a constituent with
the moved phrase's scope (= sister). Is this a substantive difference or just an insignificant variation of notation?

From the point of view of semantics, the difference is clearly substantive (or else we wouldn't have bothered to depart from familiar custom): the constituency in (5) would not have been interpretable by means of our existing inventory of composition principles (or minor revisions thereof). How about from the point of view of syntax? There it is less obvious that it makes a real difference whether we assume (3) or (5). Certain principles that refer to co-indexing will have to be trivially reformulated, but otherwise, how could we possibly tell the difference? One apparent difference in predictions concerns the question of what happens when a moved phrase moves further (as in successive cyclic movement). (5) leads us to expect that the whole unit including the index moves, as in (6), whereas (3) would seem to imply that a new binder–trace relationship is created (possibly marked by a new index j ≠ i), as in (7).

These surely *look* different, but it is harder than one might think to come up with empirical evidence that would bear on the choice. To our knowledge, the issue has not been investigated, and there are no obvious reasons why (3) (and (7)) wouldn't be just as suitable for the purposes of syntactic theory as (5) (and (6)). So we will assume henceforth that whenever we find representations like (5) in the syntactic literature, we can simply treat them as abbreviations for representations like (3).
7.4 Excursion: quantifiers in natural language and predicate logic

The movement analysis that we just looked at makes it possible for us to maintain that a quantifier word such as every or some always combines with two 1-place predicates to make a sentence. For instance, the sentence every cat is hungry is interpreted by combining the denotations of every, cat, and is hungry. And the denotation of the sentence Ann fed every cat is computed by combining the denotations of every, cat, and a predicate abstract of the form \[i \{Ann fed tj\}\]. In this respect, English quantifiers are quite unlike the quantifiers 'if and 3 of standard predicate logic (henceforth PL). The latter seem to have the syntax of sentence operators. For instance, a sentence like \(\exists x [\text{hungry}(x)]\) is built by prefixing \(\exists x\) to a sentence \(\text{hungry}(x)\).

Thus we see two salient differences between English determiners and PL quantifiers. One difference concerns the types of arguments they combine with: English determiners combine with (1-place) predicates, PL quantifiers with sentences. The other difference has to do with the number of arguments each needs in order to make a complete sentence: English determiners need two arguments, PL quantifiers only one.

As we will see in this section, the first of these differences is relatively superficial, and can easily be removed by a minor departure from the standard syntactic and semantic treatment of PL. But the second difference is more substantive, and points to an important property of natural languages that is not shared by PL.

7.4.1 Separating quantifiers from variable binding

As we have mentioned, the syntax of PL is customarily set up in such a way that a quantified expression like \(\exists x [\text{hungry}(x)]\) is built up from a sentence \(\text{hungry}(x)\) by adding a prefix \(\exists x\). But this is not the only possible way of doing it. One could equally well think of \(\exists x [\text{hungry}(x)]\) as built up from the bare quantifier \(\exists\) and an expression \(x [\text{hungry}(x)]\).

Customary parse:  

```
     sentence
    /     \  
  \exists     sentence
    /     \  
hungry     x
```
Semantically, one would then want to treat the constituent \( x \text{ [hungry}(x)\text{]} \) as a predicate abstract: It would denote the set \( \{x : x \text{ is hungry}\} \). The appropriate semantic rule would be analogous to our Predicate Abstraction Rule.

This way, a bare quantifier such as \( \exists \) or \( \forall \) can be seen to combine with one predicate to form a sentence. It is thus a 1-place second-order predicate. In other words, a bare quantifier expresses a property of sets of individuals. For example, \( \exists \) expresses the property of being a non-empty set, and \( \forall \) the property of including all elements of \( D_c \). Semantic rules to this effect would be as follows.

Let \( \alpha \) be any predicate. Then

(i) \( \exists \alpha \) is true iff \( \alpha \) denotes a non-empty set;

(ii) \( \forall \alpha \) is true iff \( \alpha \) denotes \( D \).

This reinterpretation of the syntax and semantics of PL does not alter anything substantial. We still generate all the same PL sentences (except with slightly different structures), and they get exactly the same truth-conditions as before. It does, however, eliminate the first of our two differences between English determiners and PL quantifiers: both now take predicates rather than sentences as arguments.

### 7.4.2 1-place and 2-place quantifiers

There remains the difference in the number of arguments: an English quantificational determiner requires two predicates to form a complete sentence, a PL quantifier needs only one. This difference is actually quite an obstacle to a mechanical translation of English into PL, and it has a lot to do with the most common mistakes that people make in a beginning logic class. Let's reflect a little on what we do when we are asked to symbolize English sentences in PL.

Consider two simple English sentences and their PL translations:

(1) (a) Some cat is gray.
(b) \( \exists x \ [\text{cat}(x) \& \text{gray}(x)] \)
(2) (a) Every cat is gray.
(b) \( \forall x [\text{cat}(x) \rightarrow \text{gray}(x)] \).

The transition from (a) to (b) involves an evident distortion of syntactic structure, and even the insertion of new lexical items (the connectives & and \( \rightarrow \)). Why are such changes necessary?

One way of putting the point is as follows. The PL quantifiers \( \exists \) and \( \forall \) don't really symbolize the English determiners some and every; rather, they correspond to certain complete DPs of English: namely, the DPs some individual and every individual. So when we translate English sentences into PL, we must first paraphrase them in such a way that all instances of some and every combine only with the noun individual. Combinations of the form D + noun for any other noun must be paraphrased away. For instance, we can get rid of some cat in (1a) by paraphrasing (1a) as (1c), and we can eliminate every cat from (2a) by paraphrasing it as (2c) or (equivalently, (2c')).

(1) (c) Some individual is [both a cat and gray].
(2) (c) Every individual is [gray if a cat].
(c') Every individual is [either not a cat or gray].

If it weren't for the existence of paraphrases of this kind, we wouldn't be able to express the meanings of (1a) and (2a) in PL at all.

In a way it is a lucky coincidence that such paraphrases exist. It depends on the particular lexical meanings of some and every. The meaning of English some, for instance, ensures that, for two arbitrary predicates \( \alpha \) and \( \beta \), some \( \alpha \beta \) is equivalent to some individual \([\alpha \text{ and } \beta]\) (and not, for example, to some individual \([\text{not } \alpha] \text{ or } \beta]\), except for certain specific choices of \( \alpha \) and \( \beta \). Likewise, it is a fact about the particular meaning of English every that every \( \alpha \beta \) is always equivalent to every individual \([\text{not } \alpha] \text{ or } \beta]\) (but normally has very different truth-conditions from every individual \([\alpha \text{ and } \beta]\)).

A question that suggests itself at this point is the following: Do all English determiners support systematic equivalences of this kind? More precisely, if \( \delta \) is an arbitrary determiner of English, is there going to be some way of forming out of two predicates \( \alpha \) and \( \beta \) some complex predicate \( F(\alpha, \beta) \) by means of and and not, in such a way that (for arbitrary \( \alpha \) and \( \beta \)) \([\delta \alpha \beta] \) is equivalent to \([\delta \text{ individual(s)}] F(\alpha, \beta)\)?

It turns out that the answer to this question is “no”. The standard counterexample is the determiner most. It has been proved that there is no way to construct a predicate \( F(\alpha, \beta) \) out of \( \alpha \) and \( \beta \) by conjunction and negation so that most \( \alpha \beta \) is always equivalent to most individuals \( F(\alpha, \beta) \). We will not attempt to prove this general claim here — you will have to take our word for
But the following exercise should at least give you some intuitive appreciation of it.

Exercise

Suppose we add a quantifier \( M \) to PL and give it the following interpretation:

\[ M \alpha \text{ is true iff } \alpha \text{ is true of more individuals than it is false of.} \]

(We presuppose here the alternative parse introduced in 7.4.1.) An appropriate English gloss for \( Mx \ [F(x)] \) would be “most individuals are F”.

Now consider the following two proposed PL symbolizations for the English sentence *Most cats are gray*.

(i) \[ Mx \ [cat(x) \& \ gray(x)] \]

(ii) \[ Mx \ [cat(x) \rightarrow \ gray(x)] \]

Neither of them is right. Explain why not. For each symbolization, describe a possible state of affairs where its truth-value deviates from the intuitive truth-value of the English sentence.

What exactly is the significance of the fact that natural languages have determiners like *most*? One thing that it implies is that English *most* sentences cannot be translated into PL. This in itself, however, would not be so interesting if it could be traced to a mere limitation in the *lexicon* of PL. As it turns out, there are other English quantifiers that cannot be translated into PL: for instance, *finitely many*. As logicians have proved, an English statement of the form

\[ (3) \ [\text{finitely many } \alpha] \beta \]

cannot be symbolized in PL. But in this case, the problem is not that the English determiner *finitely many* takes two arguments rather than just one. \( (3) \) is equivalent to \( (4) \).

\[ (4) \ \text{finitely many individuals } [\alpha \text{ and } \beta] \]

(For instance, *Finitely many natural numbers are smaller than 100* has exactly the same truth-conditions as *Finitely many individuals are natural numbers and smaller than 100.*) In this case, the problem is simply that PL has so few
quantifiers: namely, just $\exists$ and $\forall$. This is, so to speak, a mere limitation in the lexicon of PL, and it could be relieved, without any change to the syntax and compositional semantics of PL, by adding further quantifier symbols. For example, we might add to $\exists$ and $\forall$ such symbols as $\neg \exists$, $\exists !$, $\exists ^{-}$, and $\neg \exists ^{-}$, and define their semantics as follows:

(5) Let $\alpha$ be a 1-place predicate. Then:

(i) $\exists \neg \alpha$ is true iff $\exists \alpha$ denotes the empty set;
(ii) $\exists ! \alpha$ is true iff $\exists \alpha$ denotes a singleton set;
(iii) $\exists ^{-} \alpha$ is true iff $\exists \alpha$ denotes an infinite set;
(iv) $\neg \exists ^{-} \alpha$ is true iff $\exists \alpha$ denotes a finite set.

By enriching the lexicon of PL in this way, we could make expressible such English statements as (3). But we still couldn't express most. The problem with most runs deeper. It is an irreducibly 2-place quantifier, and adding more 1-place quantifiers to PL will therefore not help.

7.5 Choosing between quantifier movement and in situ interpretation: three standard arguments

We began this chapter by observing a problem with quantifiers in object position. Given the semantic type for quantificational DPs that we had come up with in chapter 6, we predicted that they were interpretable only when they were sisters to a phrase of type $<e,t>$. This prediction was falsified by the distribution of quantificational DPs in English — or at least, it was falsified by their surface distribution. We presented two different responses to this problem. One response took for granted that the surface distribution of quantifiers essentially reflected their distribution at the level at which semantic interpretation takes place, and it consisted of positing a systematic type ambiguity in the lexical meanings of determiners. The other response was to stick by the prediction and draw appropriate conclusions about the syntactic behavior of quantifiers. For the time being, the choice between them seems to be open. There is a trade-off between positing lexical type-shifting rules and positing movement. Pending independent motivation for either one or the other, the choice seems to be just a matter of taste.

But maybe if we broaden our scope of linguistic phenomena, we will find that there is some independent evidence for one choice over the other? Indeed, this has been claimed to be the case. Specifically, it has been argued that the
movement approach has a decided advantage in dealing with three phenomena: scope ambiguity, antecedent-contained deletion, and bound-variable anaphora.

We will present some versions of these standard arguments in this section. But we should note at the outset that their ultimate force is very difficult to assess. For one thing, proponents of the \textit{in situ} approach have developed ever more sophisticated answers to these arguments, and the most recent theories that they have developed differ from the movement-based theories in so many respects at once that comparison becomes a very global and difficult matter. And an even more important difficulty is that the issue is not really an all-or-nothing choice. It is entirely conceivable that there exist both quantifier movement and mechanisms for \textit{in situ} interpretation (of nonsubject quantifiers), and that the grammars of natural languages employ one option in some constructions and the other in others. In fact, the majority of scholars who have done influential work in this domain have (explicitly or implicitly) taken such mixed positions of one sort or another.\footnote{So we could not possibly purport here to give decisive evidence in favor of a pure movement approach. The best we can hope for is to give you some preliminary idea of what sorts of considerations are pertinent, and what price you have to pay for a particular decision.}

7.5.1 \textit{Scope ambiguity and \textquotedblleft inverse\textquotedblright scope}

If all quantifiers are interpreted in their surface positions, then a given surface structure with two or more quantifiers in it can receive only one reading, unless we admit ever more complicated types, going far beyond the flexible types options we have considered so far.\footnote{Take sentence (1):}

\begin{enumerate}
\item (1) Somebody offended everybody.
\end{enumerate}

(1) has two readings, not just one. It can mean that there is somebody who offended everybody. Or else it can mean that for everybody there is somebody that s/he offended. On either of the proposals for \textit{in situ} interpretation that we discussed above, we predict only the reading where there is somebody who offended everybody.

Once we move quantifiers, however, it is trivial to derive several distinct and truth-conditionally non-equivalent LFs from a given SS. For instance, we can derive at least the LFs (3) and (4) from the SS of sentence (2).

\begin{enumerate}
\item (2) The company sent one representative to every meeting.
\end{enumerate}
(3) $\text{S}$

$\text{DP} \rightarrow \text{D NP}$

$\text{every meeting}$

$\text{S} \rightarrow \text{DP .}$

$\text{DP} \rightarrow \text{D NP}$

$\text{one representative}$

$\text{S} \rightarrow \text{DP VP}$

$\text{VP} \rightarrow \text{V NP PP}$

$\text{the company}$

$\text{VP} \rightarrow \text{V NP PP}$

$\text{sent} \ t_1 \ \text{to} \ t_2$

(4) $\text{S}$

$\text{DP} \rightarrow \text{D NP'}$

$\text{one representative}$

$\text{S} \rightarrow \text{DP .}$

$\text{DP} \rightarrow \text{D NP}$

$\text{every meeting}$

$\text{S} \rightarrow \text{DP VP}$

$\text{VP} \rightarrow \text{V DP PP}$

$\text{the company}$

$\text{VP} \rightarrow \text{V DP PP}$

$\text{sent} \ t_1 \ \text{to} \ t_2$
If we apply our semantic rules to these two structures, they assign them distinct truth-conditions. We can prove this, for example, by considering the following state of affairs:

(5) There is exactly one company, c.
There are exactly two representatives, r₁ and r₂.
There are exactly three meetings, m₁, m₂, and m₃.
c sent r₁ to m₁, r₂ to both m₂ and m₃, and nobody else to anything else.

We can show that one of the LFs above is true in this scenario and the other one is false. More specifically:

(6) Claim 1: Given the facts in (5), \( \llbracket (3) \rrbracket = 1 \).
Claim 2: Given the facts in (5), \( \llbracket (4) \rrbracket = 0 \).

Here is a proof of claim 2:

(i) Work out the extension of the larger predicate abstract in (4):²⁵
Let's abbreviate that predicate abstract by "α":
\( \alpha := 1[every \ meeting 2[the \ company \ sent \ t₁ \ to \ t₂]] \)
Now let \( x \in D \) be an arbitrary individual. Then:
\[ \llbracket (\alpha) \rrbracket (x) = 1 \]
iff
\[ \llbracket (every \ meeting 2[the \ company \ sent \ t₁ \ to \ t₂]) \rrbracket^{\otimes x/n} = 1 \]
iff
\[ \llbracket (every \ meeting) (\llbracket (2[the \ company \ sent \ t₁ \ to \ t₂]) \rrbracket^{(1 \rightarrow x)}(y) = 1 \]
iff
for every \( y \in \{m₁, m₂, m₃\} \):
\[ \llbracket (2[the \ company \ sent \ t₁ \ to \ t₂]) \rrbracket^{(1 \rightarrow x)}(y) = 1 \]
iff
for every \( y \in \{m₁, m₂, m₃\} \):
\[ \llbracket (the \ company \ sent \ t₁ \ to \ t₂) \rrbracket^{(1 \rightarrow x)/2} = 1 \]
iff
for every \( y \in \{m₁, m₂, m₃\} \):
\[ \llbracket (the \ company \ sent \ t₁ \ to \ t₂) \rrbracket^{(1 \rightarrow x)}(y)(c) = 1 \]
iff
for every \( y \in \{m₁, m₂, m₃\} \):
\[ \llbracket (sent) (\llbracket (t₁) \rrbracket^{(1 \rightarrow x)}(y)) (\llbracket (t₂) \rrbracket^{(1 \rightarrow x)}(y)) (c) = 1 \]
iff
for every \( y \in \{m₁, m₂, m₃\} \):
\[ c \ sent \ (\llbracket (t₁) \rrbracket^{(1 \rightarrow x)}(y)) \ to \ (\llbracket (t₂) \rrbracket^{(1 \rightarrow x)}(y)) \]
iff
for every \( y \in \{m₁, m₂, m₃\} \):
\[ c \ sent \ x \ to \ y \]
iff
c sent x to m₁, m₂, and m₃.
According to (5), no individual x satisfies this condition, so we have determined:
For no \( x \in D : \llbracket \alpha \rrbracket (x) = 1. \)

(ii) It follows that there is no \( x \in D \) such that \( \llbracket \text{representative} \rrbracket (x) = 1 \) and \( \llbracket \alpha \rrbracket (x) = 1. \)
According to the lexical entry of \( \llbracket \text{one} \rrbracket ^{26} \), this implies that \( \llbracket \text{one} \rrbracket (\llbracket \text{representative} \rrbracket (\llbracket \alpha \rrbracket )) = 0. \)
But \( \llbracket \text{one} \rrbracket (\llbracket \text{representative} \rrbracket (\llbracket \alpha \rrbracket )) = \llbracket (4) \rrbracket \), so we have proved our claim 2.

---

**Exercise**

Prove claim 1 of (6). Annotate each step in your proof with references to all the rules and definitions that you are applying.

May^{27} argued that the case for quantifier movement is even stronger when we consider not just examples like (2), but also examples like (7).

(7) One apple in every basket is rotten.

May’s point about example (7) is that its *most natural* reading (perhaps even its only reading) cannot be generated by an *in situ* approach, but is straightforwardly predicted by quantifier movement.

Consider what we get on the *in situ* analysis with flexible types: “in” has the same type of meaning as a transitive verb, \(<e, et>\). So “every” must have its type \(<et, <e, et>, et>\) meaning here. Thus we get:

\[
\llbracket \text{in every basket} \rrbracket = \lambda x . \text{for every basket y, } x \text{ is in y}
\]

We proceed to the next node up by Predicate Modification and get:

\[
\llbracket \text{apple in every basket} \rrbracket = \lambda x . \text{x is an apple and for every basket y, } x \text{ is in y}
\]

Finally we apply the meanings of “one” and “rotten” and derive the following truth-condition for (7):

\[
\llbracket (7) \rrbracket = 1 \text{ iff there is at least one } x \in D \text{ such that } x \text{ is an apple and } x \text{ is in every basket and } x \text{ is rotten.}
\]
This is definitely not the salient meaning of (7). What (7) rather seems to mean is that every basket contains one rotten apple. By moving the quantifier phrase “every basket”, we can easily generate this meaning:

(7')

\[\text{S} \rightarrow \text{DP} \rightarrow \text{I} \rightarrow \text{S} \\rightarrow \text{DP} \rightarrow \text{is rotten} \]

\[\text{DP} \rightarrow \text{D} \rightarrow \text{NP} \rightarrow \text{PP} \rightarrow \text{P} \rightarrow t_1 \rightarrow \text{in} \]

Exercise

Calculate the truth-conditions of (7').

7.5.2 Antecedent-contained deletion

Another phenomenon that is a potential problem for an in situ interpretation of quantifier phrases is so-called Antecedent-Contained VP Deletion, illustrated by the following example.

(8) I read every novel that you did.

Antecedent-Contained VP Deletion is an instance of VP deletion. Here are some more examples of VP deletion.

(9) I read War and Peace before you did.

(10) I went to Tanglewood even though I wasn't supposed to.
(11) You may very well put this experience behind you, but you shouldn’t think that you really have to.

Suppose that in the construction of those sentences, a VP is deleted in the derivation from SS to PF. This VP, then, is not pronounced, but is available for semantic interpretation. The deletion is licensed by a preceding VP that has the same shape. With sentences (9)–(11), it is easy to see what the licensing VP would be. In (9), it’s “read War and Peace”. In (10), it’s “go to Tanglewood”, and in (11), it’s “put this experience behind you”. But now go back to (8). What has to be deleted in (8) is a VP consisting of the verb read and a trace that is bound by the (index introduced by) the relative pronoun. The surface structure of (8) does not seem to provide an antecedent VP that looks the same. If the object quantifier phrase in (8) is allowed to move out, however, we obtain the right antecedent VP. The LF for (8) will now look at follows:

\[
\begin{array}{c}
\text{(8')} \\
\text{IP} \\
\text{DP} \\
\text{every} \\
\text{NP} \\
\text{NP} \\
\text{novel} \\
\text{CP} \\
\text{that} \\
\text{C} \\
\text{IP} \\
\text{you} \\
\text{PAST} \\
\text{VP} \\
\text{read} \\
\text{t_1} \\
\end{array}
\]

So in examples involving Antecedent-Contained VP Deletion, we seem to need movement of object DPs anyway. So there are reasons to assume such movement is in principle available, and we might as well use it to resolve the type mismatch in simple sentences with quantified objects as well. An additional mechanism of type-shifting is, at best, redundant.
7.5.3 **Quantifiers that bind pronouns**

Consider the following sentences:

(12) Mary blamed herself.

(13) No woman blamed herself.

(14) Every woman blamed herself.

Sentences (12)–(14) contain instances of reflexive pronouns. Reflexive pronouns are necessarily anaphoric. If a pronoun is used anaphorically, its value is determined by its antecedent. If the antecedent is a proper name, then a pronoun that is anaphorically related to it might simply inherit the proper name’s referent as its semantic value. But what if the antecedent is a quantifier phrase? (13) is not synonymous with (13’), and (14) is not synonymous with (14’). Hence we don’t seem to be able to claim that reflexives always inherit the denotation of their antecedent.

(13’) No woman blamed no woman.

(14’) Every woman blamed every woman.

Reflexives are not the only pronouns that can be anaphorically related to quantifier phrases. Pronouns like he, she, it have such uses as well. This is shown by the following examples:

(15) No man noticed the snake next to him.

(16) We showed every woman a newspaper article with a picture of her.

Again, it would not do to say that these pronouns simply inherit the denotations of their antecedents. (15) does not mean the same as “No man noticed the snake next to no man”.

So how should we interpret these reflexives and pronouns? It seems that they behave as bound variables. We have successfully treated some other cases of bound-variable pronouns in the chapter on relative clauses. Can’t we just combine our treatment of pronouns from that chapter with the present treatment of quantifier phrases? Let’s try.

On the quantifier movement approach, the matter is straightforward. Although subject quantifiers are not forced to move in order to avoid a type mismatch, there is no reason why they shouldn’t be allowed to move. Suppose we exercise
this option, and also choose to co-index the pronoun with the trace left by the moved quantifier. This leads to the following representations.

\[ (17) \]

\[ S \]
\[ \rightarrow \]
\[ DP \]
\[ \rightarrow \]
\[ every \, woman \, 1 \, S \]
\[ \rightarrow \]
\[ DP \quad \rightarrow \quad VP \]
\[ \rightarrow \quad t_1 \quad V \quad DP \]
\[ \rightarrow \quad \text{blamed} \quad \text{herself}_1 \]

\[ (18) \]

\[ S \]
\[ \rightarrow \]
\[ DP \]
\[ \rightarrow \]
\[ no \, man \, 1 \, S \]
\[ \rightarrow \]
\[ DP \quad \rightarrow \quad VP \]
\[ \rightarrow \quad t_1 \quad V \quad DP \]
\[ \rightarrow \quad \text{noticed} \quad \text{D} \quad \text{NP} \]
\[ \rightarrow \quad \text{the} \quad \text{NP} \quad \text{PP} \]
\[ \rightarrow \quad \text{\textasciitilde} \text{snake} \quad \text{P} \quad \text{DP} \]
\[ \rightarrow \quad \text{next to} \quad \text{him}_1 \]

The rules for interpreting these structures are already all in place. Look up, in particular, our Traces and Pronouns Rule. Although we did not have reflexive
pronouns in mind when we first wrote that rule, our formulation just mentions a "pronoun"; so it will apply equally to "herself₁" in (17) and "him₁" in (18) – provided only that the lexicon classifies both as pronouns.

---

**Exercise**

Calculate the predicted truth-conditions for one of the structures (17), (18).

On a pure *in situ* approach to quantifiers, it is less obvious how to derive the appropriate meanings. Suppose the pronoun is co-indexed with its antecedent, but that antecedent is an unmoved DP:

\[(19)\]

\[ S \]
\[ DP₁ \]
\[ \text{no woman blamed herself₁} \]

Does (19) express the correct truth-conditions?

The interpretation of the VP-node in (19) is straightforward. Given that we are interpreting "herself₁" as a variable, by the Traces and Pronouns Rule, it receives the following assignment-dependent meaning:

For any a: \[ \text{[blamed herself₁]}^a = \lambda x . x \text{ blamed a(1)}. \]

But how do we continue up to the S-node? If we simply compose the VP meaning with the meaning of the DP "no woman", ignoring the index on the latter, we wind up with this:

For any a: \[ (19]^a = 1 \text{ iff no woman blamed a(1)}.\]

This cannot be what we want, since it does not give us an assignment-independent truth-value for (19).

Somehow, we must interpret (19) in such a way that the variable 1 gets bound. Where is the variable binder that does this? The only candidate is the index on the DP "no woman". In light of our earlier discussion, let's try to think of (19) as an abbreviation for (19').
(19')

Is this an interpretable structure? We get without problems to the (unlabeled) node above VP. According to our current formulation of PA, this denotes the following function:

\[
\lambda x . [\text{blamed herself}_1]^{[1 \to x]}
\]

This turns out to be a function of type \(<e_1, e_2, t_1>\). (Exercise: Calculate this function, and find an English word or phrase that actually denotes it.) So it is not a suitable argument for the basic type \(<e_1, e_2, t_1>\) meaning of "no woman", but forces us to employ the alternate type \(<e_2, e_3, e_4>\). But then the whole sentence in (19') receives a meaning of type \(<e_1, t_1>\) instead of a truth-value! (Exercise: Calculate the predicted meaning of (19'), and find an English word or phrase that actually denotes it.)

So reading (19) as short for (19') did not work in this case. What other options do we have? Well, one thing we can certainly do is take (19) at face value and introduce a new composition rule specifically for structures of this form. Here is a proposal that yields the correct predictions for our example:

(20) If \(\alpha\) has the daughters \(\beta_1\) and \(\gamma\), where \(\llbracket \beta \rrbracket^a \in D_{e_1, t_1}\) for all \(a\), then, for any assignment \(a\):

\[
[\alpha]^a = [\beta]^a(\lambda x . [\gamma]^{a_{\text{sw}}}(x))
\]

Applying this to (19'), we obtain:

\[
[19'] = 1\quad \text{iff (by (20))}
\]

\[
[\text{no woman}] (\lambda x . [\text{blamed herself}_1]^{a_{\text{sw}}}(x)) = 1\quad \text{iff (by meaning of "no woman")}
\]

there is no woman \(y\) such that \([\lambda x . [\text{blamed herself}_1]^{a_{\text{sw}}}(x)](y) = 1\)
iff (by definition of λ-notation)
there is no woman \( y \) such that \([\text{blamed herself}]^{\eta'}(y) = 1\)
iff (by meaning of VP, as calculated earlier)
there is no woman \( y \) such that \( y \) blamed \( a^{\eta'}(1) \)
iff (by definition of assignment modification)
there is no woman \( y \) such that \( y \) blamed \( y \).

This works, but the addition of a new composition rule is costly, and, *ceteris paribus*, we would prefer an alternative that avoids it. Moving the quantified subject seems to be just such an alternative. By doing that, we were able to interpret the examples with bound-variable pronouns with no more, semantic rules than we had already motivated before.

---

**Exercise**

It might be suspected that the problem we have just outlined for the *in situ* approach is really an artifact of an inadequate view of pronouns. Wouldn't things be easier if we didn't think of "herself" as a variable in the first place? Then we wouldn't need to worry about getting it bound. In this exercise, you get to explore a non-variable semantics of reflexives which at first sight is quite attractive.\(^{33}\)

Maybe a reflexive is a special sort of second-order predicate, as in the following lexical entry:

\[ [\text{herself}] = \lambda f \in D_{<e, <e, l>}. \lambda x \in D. f(x)(x) \]

This meaning is of type \(<<e, et>, et>\) so as to be able to take the \(<e, et>>\) type meaning of a transitive verb like *blame* as its argument and yield an \(<et>\) type VP meaning as the value.

(a) Show how sentence (13) above is interpreted on this proposal.
(b) What does the proposal predict for the following example?

**No woman bought a book about herself.**

Draw a tree for this example; then work through it under the flexible types theory that was presented above, using suitable readings for the type-ambiguous determiners "no" and "a".
Notes


3 R. Cooper, “Montague's Semantic Theory and Transformational Syntax” (Ph.D. dissertation, University of Massachusetts, Amherst, 1975), idem, Quantification and Syntactic Theory (Dordrecht, Reidel, 1983).


7 That is the one which Richard Montague used in his paper “On the Proper Treatment of Quantification in Ordinary English” (“PTQ”), in R. H. Thomason (ed.), Formal Philosophy. Selected Papers by Richard Montague (New Haven, Yale University Press, 1974), pp. 247–70. Abstracting away from matters irrelevant to the present issue, Montague's lexical entries for transitive verbs specify denotations of type \(<\text{<e,t>},\text{t}>\),\text{<e,t>}>\). Exercise: Define a meaning of this type for the verb “offend”, and show that your proposal correctly predicts the truth-conditions of sentence (1).

8 The subscripts on the quantifier phrases help distinguish between homonyms of different semantic types. They shouldn't be confused with the indices we use for the interpretation of pronouns and traces.
9 Here and below, we sometimes leave out the innermost angled brackets and commas in complicated type expressions. For instance, 
"<<e,et>,et>" abbreviates "<<e,<e,et>,<e,et>>". This notational practice improves readability without introducing any ambiguity.

10 Actually, this is not quite right: (3) as formulated generates only variants of type <et,<<e,et>,<e,et>>, not those of even more complex types which we would need, e.g., in the innermost object position of a 3-place verb like "introduce". It is possible, however, to formulate a more general version of (3) which recursively defines meanings for a whole infinite supply of alternate determiners of all the types that could possibly be required. If you are interested, you can consult, for instance, Mats Rooth's 1985 UMass Ph.D. thesis, "Association with Focus."

11 For an extensive discussion, see E. L. Keenan and L. Faltz, *Boolean Semantics* (Dordrecht, Reidel, 1983).

12 We assume binary branching because our system of composition rules is not equipped to deal with ternary branching nodes. If instead we extended the system so that it could deal with ternary branching, this should not substantially affect the point about type flexibility that we are interested in at the moment.


14 At least not in most of the syntactic literature. We are aware of one exception: Edwin Williams and Henk van Riemsdijk, for reasons of their own, have actually proposed syntactic structures more like (3) than like (5), not only for QR but also for wh-movement. See, e.g., Edwin Williams, "A Reassignment of the Functions of LF," *Linguistic Inquiry*, 17/2 (1986), pp. 265–99, for an argument why "Who left?" should have structure (i) instead of (ii).

15 The best we could do if forced to interpret (5) is to add a new composition principle which wraps up predicate abstraction and functional application of quantifier to argument all in one operation. *Exercise:* Formulate that principle.


17 We are using "sentence" here in the sense of what logicians call a "formula". "Sentences" in the narrower sense of logic are only those formulas that don't contain any variables free. Our non-standard terminology may not be optimal for the purposes of logicians, but it is more natural from a linguist's point of view, and it facilitates the comparison between natural language and logic that interests us here.
Or, more accurately, the characteristic function of this set. As usual, we will indulge in set talk.

Where the predicate individual is to be understood in its most general sense, in which it is true of every element of D.

The equivalence of (2c) and (2c') is a consequence of the semantics of if. As you recall from propositional logic, if \( \phi \), \( \psi \) is equivalent to \( \text{not} \ \phi \) or \( \psi \).

We could also say here: "by means of and, not, or, if, and iff". But this wouldn't make any difference, because it is known that or, if, and iff are definable in terms of and not. (See Gamut, Logic, vol. 1, ch. 2, or any other introduction to propositional logic.)


See Hendriks, “Type Change." Cooper stores are another option; see Cooper, Quantification.

In some of the steps of this derivation, we have collapsed more than one rule application. If necessary, fill in the intermediate steps.

In the first step of our calculation, we presuppose that \( [\alpha] = [\alpha]^\emptyset \). To make sense of this, remember that the empty set \( \emptyset \) is a special case of a partial function from \( \text{IN} \) to \( D \), hence qualifies as an assignment in the technical sense of our definition from chapter 5. \( \emptyset \) is the (extremely partial) assignment which has no variable at all in its domain. It is generally the case that a tree will have a well-defined semantic value under the assignment \( \emptyset \) just in case it has the same semantic value under all assignments. So "\( \{\phi\} \)" can always be replaced by "\( \{\phi\}^\emptyset \)".

We assume here that one has the following entry: \( \text{[one]} = \lambda x \in D_{\text{ex}, \text{t}} \cdot [\lambda g \in D_{\text{ex}, \text{t}} \cdot \text{there is at least one } x \in D \text{ such that } f(x) = g(x) = 1] \).


We can leave open for the moment whether it is a grammatical (though implausible) reading, or not even that. See ch. 8.

May dubbed this the “inversely linked” reading, because the scope relations here are in some sense the “inverse” of the surface hierarchy: the superficially more embedded quantifier has the widest scope.

VP deletion will be discussed in chapter 9.

The account of Antecedent-Contained Deletion we just sketched presupposes that information about the presence of a suitable antecedent VP at LF is available during the derivation from SS to PF. Apparent "communication" between LF and PF is a currently much debated issue that comes up independently in the area of focus interpretation and pronunciation. As for Antecedent-Contained Deletion, some authors assume that DPs are not truly in their VP-internal positions at PF any more. They have already moved up the tree, and occupy a position within a functional projection. See Hornstein, Logical Form, for example.

This argument is originally due to Ivan Sag, Deletion and Logical Form, MIT (Ph.D. dissertation, MIT, 1976; published by Garland, New York, 1980), and

33 More sophisticated versions of this analysis have been extended to all pronouns (not just reflexives). Some of them are serious competitors for the quantifier movement solution we adopt in this book, and thus potentially invalidate the objection to *in situ* approaches we have raised here. See especially the works of Pauline Jacobson, cited in n. 5.
8 Syntactic and Semantic Constraints on Quantifier Movement

That quantifying DPs are (sometimes or always) interpreted in positions that are different from their argument positions and related to them by movement is an idea that has a long tradition in generative syntax. Carden¹ (working in a framework where surface structure was transformationally derived from the input to semantic interpretation) proposed a rule of Quantifier Lowering, and an analogous raising rule was assumed by Chomsky² to affect quantifiers in the derivation from Surface Structure to Logical Form. At each step in the evolution of syntactic theory, syntacticians have sought to describe quantifier movement within a general theory of movement and to deduce as many of its properties as possible from basic principles of syntax. An up-to-date overview and assessment of this enterprise would be way beyond the scope of this book. Our purpose here is merely to clarify some of the respects in which certain questions that have been asked about the syntax of quantifiers depend on assumptions about their semantics.

From the perspective which we are taking in this book, we expect quantifier movement to be constrained from two independent directions: by interpretability and by the laws of syntax. Ideally, we expect to find that every observation about where quantifiers move can be fully explained by the assumptions that (i) every derivation must terminate in an interpretable structure and (ii) quantifier movement is subject to exactly the same syntactic laws as every other kind of movement. Which structures are interpretable depends, of course, on our semantic analysis. Specifically, it depends on our inventory of composition rules and lexical type-shifting rules. We will adopt as a working hypothesis the most restrictive theory in this regard. Our composition rules are only Functional Application, Predicate Modification, and Predicate Abstraction, and there are no type-shifting rules at all. This means that fewer potential structures will be interpretable than on alternative assumptions.

Let us consider in this light some of the questions that appear in the syntactic literature on quantifier movement. Many of these were discussed by May,³ in the first extensive study of quantifiers within a framework similar to contemporary transformational syntax. While May's particular answers in that work have mostly been revised or replaced by now, the questions themselves are still pertinent.
8.1 Which DPs may move, and which ones must?

May posited a special rule of Quantifier Raising (QR), formulated as follows:

(1) Adjoin Q (to S).

"Q" stands for "quantifier", so this rule explicitly targeted quantificational DPs, and the question of it applying to, say, a proper name or a pronoun did not even arise. Soon after May's writing, however, syntacticians questioned the existence of specific rules like (1) and began to explore a theory of grammar in which there is only a completely general rule "Move α". From this perspective, all DPs should be equally movable, and if adjunction to S is an option for any of them, it should in principle be as available for those which do not contain a quantificational determiner as for those that do. So if the syntax generates a structure like (2), why wouldn't it also generate (3)?

(2)

```
S
  DP 1 S
   every linguist John VP
       offended t₁
```

(3)

```
S
  DP 1 S
   Mary John VP
       offended t₁
```

There is no problem with interpretability. As you can easily verify, (3) has well-defined truth-conditions, equivalent to those of "John offended Mary" with "Mary" left in situ. Given this equivalence, it is difficult to see how we could have direct empirical evidence for or against the hypothesis that DPs of type e can undergo all the same movements as those of type <<e,t>,t>. The simplest assumption, then, is that they can, and this is what we will assume henceforth.

As regards quantificational DPs, May assumed that his rule (1) was, like all transformational rules, in principle optional, but that constraints on the
output of LF derivations made its application effectively obligatory. One of the constraints he posited was essentially the version of the $\Theta$-Criterion that we discussed in section 3.4. As we noted there, this $\Theta$-Criterion requires every function-denoting node to be sister to something that can be interpreted as its argument. Specifically, it requires every phrase of a type $\langle e,t \rangle$ to have a sister of type $e$. In this regard, it is stronger than our Interpretability Principle, which allows phrases of type $\langle e,t \rangle$ to take sisters of type $e$, type $\langle e,t \rangle$, or type $\langle et,t \rangle$. In cases like (2), the $\Theta$-Criterion and Interpretability both force movement, because two nodes with types $\langle e,et \rangle$ and $\langle et,t \rangle$ respectively cannot be sisters under either constraint. But the same does not hold for all quantifiers in argument positions. We have seen that quantifying DPs in subject position are interpretable in situ, and this holds more generally for whatever position realizes the highest (outermost) argument of a predicate. For instance, if unaccusative verbs like "arrive" take their only argument as an object, then there are structures in which quantifying DPs are straightforwardly interpretable in situ in object position. The stronger $\Theta$-Criterion adopted by May, by contrast, forces all quantifiers to vacate their base positions by LF, even those that are sisters to nodes of type $\langle e,t \rangle$ to begin with.

It would be interesting to obtain some empirical evidence that distinguishes between these two predictions, but this may be difficult, perhaps even impossible. Leaving a quantifier in situ never gives rise to truth-conditions which could not also be obtained by subjecting it to (short) movement. So the two theories are indistinguishable in their immediate semantic predictions. If they can be teased apart at all, it will require some imaginative use of indirect evidence. We are not aware that any arguments have been presented one way or the other, and we will leave the matter open.

Exercise

Let $\alpha$ and $\beta$ be constituents with denotations of type $\langle et,t \rangle$ and type $\langle e,t \rangle$ respectively, and let $i \in |N$ be such that $\beta$ contains no free occurrences of variables indexed $i$. Then the following two trees have the same denotation under every assignment.

```
   α  β
   /   \
  α   i
     \
    t   β
```

Prove this.
8.2 How much moves along? And how far can you move?

The “Q” in May’s QR rule was meant to stand for a feature that characterizes lexical items like “every”, “some”, “no”, etcetera. As formulated in (1), this rule therefore could in principle apply by moving just a quantificational determiner by itself, leaving the rest of the DP stranded, as in the derivation below:

(1) John fed every bird.

\[
\begin{array}{c}
\text{LF:} \\
\text{S} \\
\text{every} \\
1 \\
\text{S} \\
\text{John} \\
\text{VP} \\
\text{fed} \\
\text{DP} \\
\text{t}_1 \text{ bird}
\end{array}
\]

May offered a syntactic explanation for the unavailability of this LF. But from our point of view, such an explanation is redundant, because we are not dealing with an interpretable structure here in the first place. The trace’s type e meaning combines with the noun’s type \(<e,t>\) meaning to yield a truth-value (1) as the meaning of the DP “t1 bird”. This cannot be composed with the type \(<e,et>\) meaning of the verb, and thus the VP and all higher nodes are uninterpretable.

Exercise

The argument we just gave took for granted that traces are always interpreted as variables of type e. This is the only interpretation for traces that our semantics so far has made available, and thus we have indeed shown that structure (1) is not interpretable under our current theory. But our current theory, of course, is quite preliminary, and it might well turn out that we must eventually revise it in some way that has the side effect of making structures like (1) interpretable after all. Indeed, it has been proposed that traces can in principle receive interpretations of any semantic type, not just e. Motivation for this assumption comes from the analysis of structures in which, for example,
APs, predicative PPs, VPs, or Vs have been moved. Consider the following topicalization structures:

(i) \([_{\text{AP}} \text{ Brahms}], I \text{ adore } [_{\text{AP}} \ t_1].\)

(ii) \([_{\text{PP}} \text{ on the porch}], \text{ she isn't } [_{\text{PP}} \ t_1].\)

(iii) \([_{\text{AP}} \text{ hard-working}], \text{ he is } [_{\text{AP}} \ t_1].\)

(iv) \(\ldots \text{ and } [_{\text{VP}} \text{ buy the couch}], \text{ I did } [_{\text{VP}} \ t_1].\)

How can we interpret these structures? One answer might be: not at all – they are uninterpretable, and therefore the moved phrases must be lowered back into their base positions before the semantic component applies. But another response would be to generalize our rules for traces and abstraction so that they can handle these structures as they stand. Here is how we might go about this. Let's assume that an index is not just a number, but a pair of a number and a semantic type. We correspondingly need a revised definition of "assignment".

(v) A **variable assignment** is a partial function \(a\) from the set of indices to the set of all denotations, such that, for every \(<i,\tau> \in \text{dom}(a)\), \(a(i,\tau) \in D_\tau\).

The semantic rules stay essentially the same, except that the formulations of Pronouns and Traces and of Predicate Abstraction need trivial adjustments to fit the more complicated notion of an index:

(vi) If \(\alpha\) is a trace or pronoun, and \(i\) and \(\tau\) are a number and a type respectively, then, for any assignment \(a\), \([\alpha_{<i,\tau>}]^a = a(i,\tau)\).

(vii) If \(\alpha\) is a branching node with daughters \(\beta\) and \(\gamma\), where \(\beta\) (apart from vacuous material) dominates only an index \(<i,\tau>\), then, for any assignment \(a\): \([\alpha]^a = \lambda x \in D_\tau . [\gamma]^a_{\epsilon_{<i,\tau>}}\).

(a) Show how this proposal applies to examples (i)–(iv). (Choose one of (ii)–(iv) to illustrate.) Also show briefly that it subsumes our previous treatment of relative clauses and quantifier movement structures.

(b) What does this proposal predict regarding the potential structure (1) that we considered above?

If there is a movement operation like QR, we expect it to be subject to the very same locality constraints that are obeyed by other, better-known movement operations. Many authors have noted a close parallel between constraints on \(wh\)-movement and constraints on quantifier scope, for example. The constraints
needed are unlikely to be derivable from Interpretability. Structures with quantifiers that have been moved too far are perfectly interpretable. An unconstrained rule of QR would allow us to derive two distinct interpretable LFIs for sentence (2), for example, as you will find out when you do the exercise below.

(2) Most accidents that nobody reported were minor.

---

**Exercise**

(a) Display both LFIs for (2). Give informal, but unambiguous paraphrases of the readings represented by each, and describe a situation in which their truth-values differ.

(b) Do the readings predicted coincide with those that are actually available for (2)? If not, what sort of remedy is appropriate to improve the empirical adequacy of our theory? Think of suitable syntactic constraints you may be familiar with.

---

### 8.3 What are potential landing sites for moving quantifiers?

Built into May's formulation of QR was a stipulation that quantifiers always adjoin to nodes of the category S. Again, we consider the question whether the effects of this stipulation can be reduced (in part or in whole) to the requirement of Interpretability.

Look at the following schematic representation of a structure that is derived by adjoining a quantifying DP $\alpha$ to a node of some category $X$:

\[
\begin{array}{c}
\alpha \quad \text{(type } \langle e, t, t \rangle) \\
\quad \text{(type } \langle e, 't' \rangle) \\
\end{array}
\begin{array}{c}
i \\
X \quad \text{(type } \tau) \\
\begin{array}{c}
\quad \text{(type } \tau) \\
\quad \text{(type } \tau) \\
\end{array}
\end{array}
\]

\[
\begin{array}{c}
\text{......} \\
\text{t; ......} \\
\end{array}
\]
“τ” stands for the type of X here, whatever it may be. The node dominating X and the adjoined index is subject to the Predicate Abstraction Rule, which assigns it a meaning of type \( <e,\tau> \). Up to this point, the structure is interpretable regardless of what type τ is. (Notice that our formulation of PA places no condition on the type of γ.) However, at the next node up (the higher X-node), we obtain a value only if \( \tau = t \). Any other choice for τ would make it impossible to apply \( \alpha \) (type \( <\langle e,\tau>,t> \)) to its sister (type \( <e,\tau> \)) (and any other ways to compose these two nodes by our rules are unavailable to begin with). It follows thus that quantifier movement can only target adjunction sites whose semantic type is t.13

How does this prediction relate to May’s assumption that QR can only adjoin to the category S? The two are clearly not equivalent. They might coincide in practice, however, if independent factors conspire to make S (= IP) the only syntactic category with denotations of type t. In the majority of examples we have analyzed so far, this happened to be the case.14 But we have no good reason to be sure it will hold up when we consider more refined syntactic analyses and/or a wider range of constructions. There may well be non-S categories with meanings of type t (for instance, small clauses, or VPs with internal subjects), and if there are, we expect quantifier movement to be able to adjoin to them.

The assumption that QR always adjoins to S was never uncontroversial. Montague had already proposed (the equivalent of) quantifier adjunction to VP and NP,15 and the syntactic literature soon converged on a consensus that at least these two additional options were indeed needed, and probably adjunction to DP as well.16 Let us review some of the evidence that was presented in this literature and assess it from our current perspective.

8.4 Quantifying into VP17

8.4.1 Quantifiers taking narrow scope with respect to auxiliary negation

One fairly straightforward argument for quantifier adjunction to VP can be constructed if we combine our current general assumptions with certain particular ones about the syntax and semantics of negation. Concretely, let’s assume that auxiliary negation in English occupies an I(nflectional)-node, at LF as well as on the surface. And let’s furthermore assume that it has the familiar denotation of type \( <t,t> \):

\[
\text{\texttt{not}} = \begin{bmatrix} 1 \rightarrow 0 \\ 0 \rightarrow 1 \end{bmatrix}
\]
Such a syntax and semantics of negation imply that the VP in the LF of a sentence like “John didn’t leave” must have a denotation of type t. A natural way to ensure this is to generate the subject “John” inside the VP and let it leave a trace there when it moves to its surface subject position:

(2)                     
                   IP
                      / 
                   / 1   
                   /    /
               John   I
                  /     /
(1) not         VP
            /
        t₁    leave

The interpretation of (2) is straightforward (“do”, we take it, is vacuous). Consider now (3) and (4).

(3) Joe didn’t invite a professor.

(4) Al didn’t attend more than two meetings.

These examples are judged to be ambiguous. (3) can either mean that Joe didn’t invite any professor. Or else it can mean that there was some professor that he failed to invite. Similarly, (4) can mean that Al attended no more than two meetings; or else that there were more than two meetings from which he was absent. The second reading in each case is the one that we generate if we adjoin the object DP to IP:

(4a)                     
                   IP
                      / 
                   / 1   
                   /    /
               Al    I
                  /     /
(1) not         VP
            /
        t₁    attend t₂

     DP  2
     /  
     /    /
 more than two meetings


To obtain the first-mentioned reading, we must adjoin instead to VP:

The LF in (4b) could not have been generated with a QR rule that allows only adjunction to IP (= S). But the reading it represents is available (in fact, even preferred) in sentence (4). If our background assumptions are correct, we therefore have a straightforward empirical argument for quantifier adjunction to VP. The most reasonable general hypothesis at this point is that quantifiers can adjoin freely to any syntactic category, as long as the output of the derivation is interpretable.

8.4.2 Quantifying into VP, VP-internal subjects, and flexible types

In our effort to show that quantifiers do sometimes adjoin to VP, we have so far been presupposing our working hypothesis that there are no flexible types (or other devices for in situ interpretation). We have also taken for granted that VPs (at least in the cases we have been looking at) contain a subject position, and therefore are semantically of type t. Neither of these two assumptions is taken for granted, however, in most of the literature that deals with the issue of possible adjunction sites for QR. This must be borne in mind when one reads that literature and tries to assess the internal logic and the current relevance of its arguments. There is actually a rather intricate relation among one's answers on these three questions: (i) whether quantifiers can adjoin to VP, (ii) whether VPs are of type t or of type <e,t>, and (iii) whether flexible types are available. The three questions are logically independent of each other, but not all possible sets of answers to them are equally defensible.
In much of the early 1980s literature on the syntax of quantifier movement, the idea of VP-internal subjects was not even entertained, and it was implicitly taken for granted that VP meanings were of type \(<e,t>\). On that assumption, how compelling is the putative evidence for VP adjunction that we have reviewed in this section? Let's look again at the case of quantifier scope below auxiliary negation, as in (4).

(4) Al didn't attend more than two meetings.

Assuming (as before) that the scope of negation is fixed in some I(nflectional) position, what does the ambiguity of (4) tell us about possible adjunction sites for QR?

To get any interpretable LF out of (4) at all, we must now assume that "not" has a type-shifted meaning of type \(<et,et>\). Then, if "more than two meetings" adjoins to IP, we derive the reading on which there are more than two meetings from which Al was absent. How about the other reading? Can we still argue that we need adjunction to VP to generate that reading? Let us see what adjunction to VP would give us here:

(4c)

Is this even interpretable? Yes, but only if we employ flexible types for quantifiers, in this case, a meaning of type \(<<e,et>,et>\) for "more than two meetings"! If we do that, we can calculate as follows:

\[
[\text{more-than-two}_2 \text{ meetings } 1[\text{attend } t_1]]
\]
\[
= [\text{more-than-two}_2][[\text{meetings}]](1[\text{attend } t_1])
\]
\[
=[^{(1)}] \lambda x \in D . [\text{more-than-two}_2][[\text{meetings}]](\lambda y \in D . [1[\text{attend } t_1]](y)(x)) = \lambda x \in D . \text{there are more than two } y \in D \text{ such that } [[\text{meeting}]](y) = 1 \text{ and } [1[\text{attend } t_1]](y)(x) = 1
\]
\[
= \lambda x \in D . \text{there are more than two meetings } y \text{ such that } [1[\text{attend } t_1]](y)(x) = 1
\]
\[=^{[2]} \lambda x \in D . \text{there are more than two meetings } y \text{ such that } \llbracket \text{attend } t_1 \rrbracket^{[1 \rightarrow y]}(x) = 1 \]
\[= \lambda x \in D . \text{there are more than two meetings } y \text{ such that } \llbracket \text{attend} \rrbracket^{y}(\llbracket t_1 \rrbracket^{[1 \rightarrow y]})(x) = 1 \]
\[= \lambda x \in D . \text{there are more than two meetings } y \text{ such that } x \text{ attended } y \]
\[= \lambda x \in D . \text{ } x \text{ attended more than two meetings} \]

In checking this calculation, pay attention especially to the steps marked by superscripts on the "=". Step [1] eliminates more-than-two\textsubscript{2} (of type \llbracket<et>,<et>,t>\rrbracket) in favor of its basic homonym more-than-two\textsubscript{1} (of type \llbracket<et>,<et>,t>\rrbracket, whose lexical entry is then used in the following step). Step [2] is licensed by the Predicate Abstraction Rule, which tells us that \llbracket 1[\text{attend } t_1] \rrbracket^{\emptyset}(y) = \llbracket \text{attend } t_1 \rrbracket^{y}. (\text{It also exploits the fact that } \llbracket a.D \rrbracket = \llbracket a \rrbracket^{a} \text{ and that } \emptyset^{y} = [1 \rightarrow y].)

Given the denotation we have just calculated for the higher VP in (4e), it can be further shown that (4c) as a whole does indeed represent the desired reading of (4) (that is, that it is equivalent to (4b).) But it would be premature to conclude from this that quantifiers must be able to adjoin to VP. Since we needed flexible types for the interpretation of (4c) anyway, we might as well have left the quantifier \textit{in situ}:

\[
(4d)
\]

\[
\quad \begin{array}{c}
\text{IP} \\
\text{Al} \\
\text{(did) not} \\
\text{attend} \\
\text{more than two meetings}
\end{array}
\]

\[
(4d) \text{ is just as interpretable as (4c), and receives exactly the same truth conditions! (Exercise: Show this.)}
\]

So it looks at this point as if we cannot make a case for quantifying into VP unless we take a stand against flexible types, which in turn commits us to VP-internal subjects. If we want to show that VP adjunction is needed \textit{even if we allow flexible types}, we must come up with more complicated examples. One way to do this is to bring in bound-variable pronouns. Consider the following variant of (4):

\[
(5) \text{ Al didn't return every clock to its owner.}
\]
(5) has a reading where the direct object every clock has narrower scope than the negation, but at the same time is the antecedent of the bound-variable pronoun it. What is the LF for this reading? As we saw earlier, if the pronoun is to be bound, it must be inside a PA configuration and co-indexed with a trace. So every clock cannot have stayed *in situ* in this LF; rather, it must have adjoined to a node which is high enough to dominate both its own trace and the pronoun. The only such nodes in (5) are VP, I, and IP. But since the scope of every clock also is narrower than that of the negation in the intended reading, it cannot have raised as high as I or IP. This leaves adjunction to VP as the only possibility. In other words, (5) has a reading that is expressed by the LF in (5), *but not by any LF derivable without adjoining to VP.*

\[
(5')
\]

---

**Exercise**

Another way to show that flexible types are not sufficient to do away with quantifying into VP is to construct examples with *two* quantifiers in the VP. Spell out this argument.

In sum, then, the need for quantifier adjunction to VP can be established even if we remain neutral on the availability of flexible types (or other methods of *in situ* interpretation) – provided, of course, that we accept the background assumptions that we have made here about negation, as well as the background assumptions that we needed earlier to argue for quantifier movement in the first place.
8.5 Quantifying into PP, AP, and NP

In our earlier discussion of so-called inversely linked readings we considered the sentence "One apple in every basket is rotten" and showed how its salient reading could be derived by adjoining the embedded quantifier "every basket" to IP. We noted in passing that this sentence might also have another reading, under which it asserted the existence of a rotten apple that is simultaneously inside every basket. The pragmatic oddity of this meaning makes it difficult to decide whether we actually want our grammar to generate it. But if we look at other examples with an analogous structure, it is quite clear that analogous readings are perfectly grammatical, and therefore we have to make sure our grammar predicts them.

8.5.1 A problem of undergeneration

Consider:

(1) No student from a foreign country was admitted.

(1) is naturally understood as asserting that no student from any foreign country was admitted.

Let us first convince ourselves that this is not the inversely linked reading, and that in fact it cannot be represented by any LF in which quantifiers only adjoin to S. Here is what we get if we adjoin the more embedded DP (here, "a foreign country") to IP:\18

(2)
(2) is an interpretable LF, but it does not represent the meaning of (1) that we are interested in. For (2) to be true, there needs to be only one foreign country from which no students were admitted. So (2) might be true even when there are many students from foreign countries who were admitted. The reading of (1) that we described above, by contrast, is false as soon as there is even one student who is from a foreign country and is admitted.\textsuperscript{19}

Is there a better candidate than (2) for the LF of this intended reading? If only type t nodes are available adjunction sites for quantifier movement, there is little else we can do. Raising the “no” DP above the “a” DP, as in (3), won’t help.

![Diagram of (3)]

(3) violates both the prohibition against vacuous binding (the variable binder “1” binds nothing!) and the prohibition against unbound traces (“\textit{t}_1” is free!).\textsuperscript{20} The types all fit together fine, but the truth-conditions of (3) turn out to be dependent on the assignment. Calculate them, and you will see that they don’t even bear a remote resemblance to the ones we are after.

**Exercise 1**

Prove the following two claims about (3):

(i) For any assignment \(a\), (3) is not in the domain of \(\llbracket\ \rrbracket^a\) unless the domain of \(a\) includes 1.

(ii) For any assignment \(a\) such that (3) is in the domain of \(\llbracket\ \rrbracket^a\), \(\llbracket(3)\rrbracket^a = 1\) iff there are either no foreign countries, or no student from \(a(1)\) was admitted (or both).
Exercise 2

It may seem that the problems with the unbound trace and vacuous binder in (3) could have been avoided if we had allowed a different constituent structure for the subject DP of (3): namely, \([_{\text{DP}} \text{no student}] \ [_{\text{PP}} \text{from a foreign country}]\) instead of \([_{\text{DP}} \text{no} \ [_{\text{NP}} \text{student}] \ [_{\text{PP}} \text{from a foreign country}]\)]. Then we could have QR’d the \textit{no} DP above the \textit{a} DP without creating unbound traces or vacuous binders:

![Diagram]

Show that this is not a solution to our problem.

We have seen that, once we have QR’d “a foreign country” out of the containing DP “no student from t”, we cannot QR the latter any higher than the former. We can then at best QR it to adjoin to its immediately dominating IP-node, but that, of course, gives us no meaning distinct from leaving it \textit{in situ}.

So if we want to find an interpretable LF for (1) that’s truth-conditionally distinct from (2) yet interpretable, it seems that we will have to avoid QR’ing “a foreign country” out of its containing DP. But where could we adjoin it \textit{within} that DP? The PP- and NP-nodes have meanings of type \(<e,t>\), and the containing DP itself has a meaning of type \(<et,t>\). So these are all excluded as adjunction sites by the Interpretability principle. We have thus exhausted all the possible LF derivations, and have not found any one suited for the natural reading of (1) that we described at the beginning of this section.

This conclusion, of course, depends on our tacit assumption that we don’t have the option of flexible types. If we do allow those, the picture changes completely. Then, all we need to do to obtain the desired interpretation is to leave everything \textit{in situ}: 
Exercise 1

Identify the type-shifted meaning for “a” that we need to interpret (4), and show that the predicted meaning is the one we want.

Exercise 2

Huang, following Fiengo and Higginbotham (see n. 16), proposed that the intended reading of (1) could be generated by adjoining the smaller DP to the NP containing it:
(a) Show that (i) is equivalent to (4), provided that we use a type-shifted meaning for "a".

(b) Show that adjunction of "a foreign country" to PP would yield yet another LF with the same meaning as (4) and (i).

(c) Can you think of examples that show that QR must be allowed to adjoin to PP or NP, even if flexible types (and hence in situ interpretation) are also available? (Hint: Recall what we have said about the analogous question regarding adjunction to VP.)

8.5.2 PP-internal subjects

May was aware that examples like (1) were a problem for his hypothesis that QR always adjoined to S, and he proposed a solution. He hypothesized that in the derivation from S structure to the input of the semantic component, PP modifiers were sometimes allowed to rewrite as clausal structures. If this operation were applied to the PP "from a foreign country" in (1), it would turn it into something essentially isomorphic to the relative clause "who is from a foreign country". Now this relative clause, of course, contains an S-node, and hence an adjunction site for QR.

Exercise

Show that sentence (i) has a straightforwardly interpretable LF which represents exactly the reading that we are trying to generate for (1).

(i) No student who is from a foreign country was admitted.

(By "straightforwardly interpretable", we mean, of course, interpretable without the use of flexible types.)

May's idea, then, was to account for the reading that (1) shares with (i) by adding enough internal structure within the DP "no student from a foreign country" so that it, too, would contain an internal IP-node for QR to adjoin to. The attractiveness of this idea does not depend on any commitment to a syntactic characterization of the possible landing sites for QR. We can equally appreciate it from our present perspective, where the decisive factor is semantic interpretability, and quantifiers must therefore adjoin to nodes with meanings of type t. From this perspective, we don't need an S-node between the embedded and the containing DP; but we do need something with type t, hence something more, it seems, than the superficially present PP and NP, whose types are <e,t>.
We are reminded at this point of the assumption of VP-internal subjects that made it possible to have an additional site for quantifier scope without any syntactic structure building. It is natural to ask, then, whether it would serve our present purposes to posit subject positions in other categories, such as PP or NP, as well. Indeed, there are arguments for such covert subject positions in the syntactic literature. So let’s explore this possibility.

Suppose, for instance, there is a subject position within the PP headed by “from” in example (1):

(5)

```
DP
  no
  NP
    student
  PP
    ?
  P
    from
  DP
    a foreign country
```

What occupies this subject position? On the surface, it is evidently an empty category. Could it be a trace? Then something would have had to move from there, but it’s hard to see what that would be. Let’s assume, rather, that it is a base-generated empty category: that is, an empty pronoun usually referred to as “PRO”:

(5')

```
DP
  no
  NP
    student
  PP
    PRO
    P
    from
  DP
    a foreign country
```
How are we going to interpret the empty pronoun? An empty pronoun presumably gets the same semantic interpretation as an overt one; in particular, it can be a variable. This, at any rate, would seem to be the most promising hypothesis for us here; it would imply that the PP has a subject of semantic type e, hence that the PP's type is t, which makes it a good adjunction site for the quantifier "a foreign country".

So far, so good; but there remains a big hole in this analysis: If the PRO is a variable, it needs a binder. The meaning of sentence (1) as a whole is clearly independent of a variable assignment, so its LF cannot contain any free variables. What could possibly be binding this PRO, though? It is not contained within a predicate abstraction configuration, and there is no obvious way to create one around it by performing some suitable movement.

Here is a possible solution. Suppose that PRO is not, after all, a variable (like overt pronouns), but rather is semantically vacuous. Now this assumption looks at first like a step in the wrong direction. If PRO is vacuous, it doesn't provide an argument for [from], and hence the PP won't get a meaning of type t. We might as well have left it out altogether then, it seems. – But that's not quite right. There is a difference between generating no subject position at all and generating one that is semantically vacuous. Such an item is still visible to the syntax, hence can be subjected to movement. Look what happens then:

\[(6)\]

\[
\begin{align*}
\text{DP} & \quad \text{no} \\
\text{NP} & \quad \text{student} \\
\text{PP} & \quad \text{PRO} \quad 1 \\
\text{PP} & \quad t_1 \\
\text{DP} & \quad \text{from} \\
\text{a foreign country}
\end{align*}
\]

(6) is derived from \((5')\) by short movement of the PRO: PRO has adjoined to its immediately dominating PP, leaving an indexed trace in its previous site and giving rise to the insertion of a co-indexed variable binder at the adjunction site. (This is just how movement always operates, on our assumptions.) As a result of this movement, the lower PP now has a nonvacuous subject, the indexed
trace. It also has a binder for this subject. There are no unbound variables in (6): the trace has a binder, and the PRO, being semantically vacuous, is not a variable in need of a binder in the first place.

Now all that we have left to do to make this structure interpretable is QR “a foreign country”. The lower PP provides a suitable adjunction site, and the LF we thus obtain for the whole sentence (1) is (7):

(7)

We leave it to the reader to calculate the interpretation of (7), and thereby verify that it captures the intended reading of (1) – without committing us to flexible types.

In sum, we saw how positing an underlying subject position in PP may help us account for the NP-internal scope readings of DPs embedded within other DPs, while still maintaining the smallest possible inventory of composition principles and lexical meanings for quantifiers. This analysis did not have to stipulate anything apart from lexical properties of PRO: PRO was taken to be a DP that is visible to the syntax, but not to the semantics. Interpretability did the rest. Being a DP, PRO could undergo QR, just like any other DP. In order for the whole structure to be interpretable, PRO had to undergo QR.

8.5.3 Subjects in all lexically headed XPs?

Once we entertain a subject position not only in VP but also in PP, it is natural to do the same for the remaining categories that were traditionally analyzed as
1-place predicates: namely, APs and NPs. Regarding APs, it is easy to construct examples analogous to our (1), except with an AP modifier instead of the PP modifier:

(8) No student [\_\_\_ interested in more than one topic] showed up.

(8) has a natural reading paraphrasable as "no student who was interested in more than one topic showed up". We can generate this reading by giving the AP a PRO subject and then adjoining both this PRO and the quantifier "more than one topic" to the AP. This analysis is exactly parallel to the one we just applied to (1), so we need not elaborate further.

Are there examples which require a subject in NP?\(^\text{22}\) (9) looks like a case in point.

(9) No owner of an espresso machine drinks tea.

This example differs from (1) in that the PP here is presumably an argument of the noun, not a modifier. In other words, the "of" in (9) (unlike the "from" in (1)) is semantically vacuous, and the meaning of "owner" is of type <e,et>. Positing a subject in PP here would not be helpful: due to the vacuity of "of," there is no chance of getting a type t meaning for the PP anyway. How, then, do we generate the meaning of (9) (paraphrasable as "no one who owns an espresso machine drinks tea")? A PRO subject in the NP seems the obvious solution. The LF then looks as in (9'), and you are invited to show that this means what it should.

(9')
Exercise 1

Examples (1) and (8) do not really show that we need subjects in PP or AP as well.

(a) Show that the relevant readings of (1) and (8) can also be accounted for if NPs have subjects, but PPs and APs don't.
(b) Can you think of examples that establish the need for a PP- or AP-internal subject, even if NP-internal subjects are also available?

Exercise 2

Show how the assumption of empty subjects in PPs and APs helps with the interpretation of pronouns in DPs of the following kind:

(i) No woman in her (own) office . . .
(ii) No man attentive to his (own) needs . . .

8.6 Quantifying into DP

May criticized his earlier analysis of inversely linked readings and proposed an alternative, on which the embedded quantifier was not adjoined to S but to the larger DP. For instance, instead of the LF structure (2) that he had assumed for the (most natural) reading of (1), he posits the one in (3).

(1) One apple in every basket is rotten.

(2) 

\[
\text{IP} \quad \text{1} \quad \text{IP} \\
\text{DP} \quad \text{every basket} \quad \text{DP is rotten} \\
\text{one NP} \quad \text{apple PP} \\
\text{in t}_1
\]
Before we turn to the reasons why May (and others) thought that (3) might be preferable to (2), let us first see whether (3) is interpretable at all and expresses (as May assumes) the same truth-conditions as (2).

It is evident that (3) is not interpretable on the most restrictive assumptions that we have been trying to maintain up to now: The DP one apple in \( t_i \) has a meaning of type \( \langle \text{et,t} \rangle \). The node above it that includes the variable binder 1, then, gets a meaning of type \( \langle e, \langle \text{et,t} \rangle \rangle \). (This follows as usual from the semantics of Predicate Abstraction.) This cannot combine with the standard type \( \langle \text{et,t} \rangle \) meaning of \( \text{every basket} \). So we need flexible types. Specifically, we need here a meaning for \( \text{every basket} \) that is of type \( \langle e, \langle e, \langle \text{et,t} \rangle \rangle, \langle \text{et,t} \rangle \rangle \). Let's define such a meaning which not only has the desired type but also gives rise to the intended truth-conditions for (3):^25

\[
(4) \quad [\text{every basket}] = \lambda f \in D_{\langle e, \langle \text{et,t} \rangle \rangle} . [\lambda g \in D_{\langle \text{et,t} \rangle} . \text{for every basket } x, f(x)(g) = 1]
\]

Exercise

Apply (4) in the interpretation of (3) and thereby convince yourself that (3) is equivalent to (2).

Having shown that (3) is in principle viable as an alternative to (2), we turn to the question of whether there is any evidence for such LFs.
8.6.1 Readings that can only be represented by DP adjunction?

When we argued for quantifier adjunction to VP, PP, etcetera, we based our arguments on the existence of readings which could not be generated without these adjunction options. Can we come up with analogous positive evidence for adjunction to DP? This is not easy. The fact that (3) represents an attested reading of (1), for instance, shows nothing, because (3) has an equivalent counterpart involving only S adjunction (namely (2)). Is it possible in principle to construct DP adjunction LFs that do not have such equivalent IP adjunction alternatives?

The following example provides a possible candidate:26

(5) John met neither a student from every class nor a professor.

We take it that (5) can mean that neither (i) nor (ii) are true (that is, that both are false):

(i) For every class x, John met a student from x.
(ii) John met a professor.

The connective neither ... nor here conjoins two quantificational DPs, so it seems to have a meaning of type \( <<et,t>,<<et,t>,<et,t>>, \), as follows:27

\[
\text{either} \ldots \text{nor} = \lambda f \in D_{<et,t>} . \lambda g \in D_{<et,t>} . \lambda h \in D_{<et,t>} . f(g) = f(h) = 0.
\]

Indeed, with (6) and (an analog of) (4) at our disposal; we can show that the following LF has precisely the meaning we just described for (5):

\[(5')\]

\[
\begin{array}{c}
\text{IP} \\
\text{DP} \quad \text{1} \quad \text{IP} \\
\text{DP} \quad \text{neither-nor} \quad \text{DP} \\
\text{DP} \quad 2 \quad \text{DP} \\
\text{every class} \quad \text{a student from } t_2 \\
\text{a professor} \\
\text{John met } t_1
\end{array}
\]

(Exercise: Prove this.) But – and this is the point of the example – there is no equivalent alternative in which every class adjoins to an IP-node instead.
Exercise

Show that (i) does not have the same truth-conditions as (5').

(i)

\[
\begin{array}{c}
\text{DP} \\
\text{every class} \\
\text{IP} \\
\end{array}
\]

\[
\begin{array}{c}
\text{DP} \\
\text{neither-nor} \\
\text{IP} \\
\end{array}
\]

\[
\begin{array}{c}
\text{DP} \\
a \text{student from } t_2 \\
\text{IP} \\
\end{array}
\]

\[
\begin{array}{c}
\text{DP} \\
a \text{professor} \\
\end{array}
\]

8.6.2 Indirect evidence for DP adjunction: a problem with free IP adjunction?

The starting point for May's reanalysis of inverse linking readings was a syntactic consideration: We know from the study of movement in general that extraction out of a subject DP leads to ungrammaticality. For instance, \emph{wh}-movement as in (7) is impossible:

(7) *the basket which one apple in \( t \) is rotten . . .

It is \emph{prima facie} implausible, therefore, that the essentially parallel configuration in (2) above should be syntactically well-formed. The assumption that it is leads to unattractive complications in the syntactic theory of movement, and this motivates us to look for an alternative representation of the meaning that (2) expresses. (3) is a more attractive candidate in this regard, because it does not involve extraction \emph{out} of the subject DP, as do (2) and (7).

This consideration by itself is not decisive in favor of allowing (3) and excluding (2). All it does is alert us to a trade-off: if we want the simplest syntax, we must complicate the semantics, and vice versa. But May did not stop here. He, and later Larson, argued that there is empirical evidence for locality constraints on QR which allow a DP-internal quantifier to move only as far as the edge of the containing DP (hence licensing (3) but not (2)).

Consider a sentence with three quantifiers such as Larson's (8).
Two politicians spy on someone from every city.

(8) Two politicians spy on someone from every city.

(8) has a number of possible readings, but Larson reports that it cannot mean "for every city x, there are two politicians who spy on someone from x". Larson's generalization is that sentences of this form do not allow readings where a third, separate DP takes scope in between the two nested DPs. Here is a summary of possible and impossible scope orders that Larson observes:

(a) ok every city x [someone from x]y [two politicians z [z spies on y]]
(b) ok two politicians z [every city x [someone from x]y [z spies on y]]
(c) every city x [two politicians z [[someone from x]y [z spies on y]]]

To get an intuitive grasp of what the missing reading (c) amounts to, imagine a situation in which it would be true, but (a) and (b) false. Let there be two cities, LA and NY. Each has two natives, l1 and l2 for LA, n1 and n2 for NY. There are four politicians p1, . . . , p4, each of whom spies on exactly one person: p1 on l1, p2 on l2, p3 on n1, and p4 on n2. (c) is true here. (a) is false, because not every city (in fact, neither) has a native on whom two politicians spy. (b) is false, because there aren't two politicians (in fact, not even one) who spy on natives of every city.

The point about the example now is this: If a quantifier embedded inside another DP were allowed to move out of the containing DP and adjoin to the IP above it, then all three of (a)–(c) could be generated. If, on the other hand, such a quantifier can move at most as far as the edge of its containing DP, then (a) and (b) can be generated, but (c) cannot be. (Draw the relevant structures and calculate their meanings to convince yourself of this claim.)

So the unavailability of reading (c) for (8) yields an indirect argument for the existence of quantifier adjunction to DP. Once we make adjunction to IP unavailable for DP-contained quantifiers – which is a good move if we want to exclude (c) – then adjunction to DP seems to be the only way to represent those inverse-linking readings that are grammatical.

8.6.3 Summary

The argument we have just presented is sound and important, but the conclusion to which it has led May and Larson is nevertheless rather unattractive. Not only does it force us to admit type flexibility after all; it also raises – as May and Larson were well aware – a serious problem with bound pronouns. Consider an example from May's book:

(9) Someone from every city despises it.
(9) shows that the embedded quantifier in an inverse-linking sentence can be the antecedent of a bound-variable pronoun which is outside the containing DP. The intended reading of (9) is straightforwardly predicted if every city is allowed to adjoin to S:

\[(9')\]

\[
\begin{array}{c}
\text{IP} \\
\text{DP} \\
\phantom{\text{DP}} 1 \\
\text{IP} \\
\phantom{\text{IP}} \text{every city} \\
\end{array}
\]

\[
\begin{array}{c}
\text{DP} \\
\phantom{\text{DP}} \text{someone from } t_1 \\
\text{VP} \\
\phantom{\text{VP}} \text{despises } i_t \\
\end{array}
\]

But in a DP adjunction structure like (9") the pronoun is free:

\[(9'')\]

\[
\begin{array}{c}
\text{IP} \\
\text{DP} \\
\phantom{\text{DP}} 1 \\
\text{DP} \\
\phantom{\text{DP}} \text{every city} \\
\text{VP} \\
\phantom{\text{VP}} \text{despises } i_t \\
\end{array}
\]

\[
\begin{array}{c}
\phantom{\text{IP}} \text{someone from } t_1 \\
\end{array}
\]

In the response to this problem, May and Larson propose far-reaching revisions of the theory of pronominal anaphora. We are not prepared to discuss these at this point, and therefore we must withhold judgment as to whether their overall proposals can ultimately justify the price we have to pay for them.³¹

Notes

4 Ibid., p. 18.
6 May called it the “Predication Condition” (“Grammar of Quantification,” p. 22).
7 The type <e,t> option is due to the existence of Predicate Modification and not directly relevant to the present discussion. Both e and <<et>,t> are possible because we have Functional Application.
8 It may well be that there are syntactic reasons (e.g., Case Theory) that force the outermost argument of each predicate to move anyway, independently of the Θ-Criterion or Interpretability. See N. Hornstein, *Logical Form. From GB to Minimalism* (Oxford, Basil Blackwell, 1995).
9 May’s explanation relied on his “Condition on Analyzability”, which said that a transformation whose structural description mentions a specifier must affect the minimal DP containing that specifier which isn’t immediately dominated by another DP. Independent motivation for this condition came from facts about pied-piping in wh-movement. See May, “Grammar of Quantification,” for details.
12 We are *not* adopting the revision that was tentatively entertained in the exercise in section 8.2.
13 More accurately, this restriction holds for the final landing site of a moved quantifier. Intermediate landing sites are not so constrained, for the same reason for which movement of DPs of type e is not. *Exercise*: Show that Interpretability considerations place no constraints at all on the type of the adjunction site for a moved name or an intermediate trace of a moved quantifier.
14 One exception (but perhaps not an interesting one) results from our decision to treat the complementizers (in relative clauses) as semantically vacuous. This implies that CN constituents have meanings of type t as well.
15 See R. Montague, “The Proper Treatment of Quantification in Ordinary English” (1971), in R. Thomason (ed.), *Formal Philosophy* (New York, Academic Press, 1974), pp. 247–70. We say “the equivalent of”, because Montague worked in a framework of syntax which did not have transformational rules in a literal sense. His syntactic category labels were also different: our VP corresponds to his IV (“intransitive verb”), and our NP to his CN (“common noun”).
16 That QR needs to be able to adjoin to VP was argued by (among others) E. Williams, “Discourse and Logical Form,” *Linguistic Inquiry*, 8 (1997), pp. 101–39;
Constraints

Quantifier Movement


Adjunction to DP was defended by R. May, Logical Form: Its Structure and Derivation (Cambridge, Mass., MIT Press, 1985), as well as by M. Rooth, "Association with Focus" (Ph.D. dissertation, University of Massachusetts, Amherst, 1985); and R. Larson, "Quantifying into NP" (unpubl. MS, 1987). (For these authors, "NP" corresponds to our "DP".)

17 The locution "quantifying into" comes from the philosophical literature, and came into linguistics through Montague's papers. Montague employed something more like a Quantifier Lowering rule in his syntactic derivations. But we use his terminology here in an extended sense. For us, "quantifying into category X" means adjoining a quantifier to X.

18 In (2) and below, we disregard the VP-internal subject position and the movement of the subject from there. This simplification does not affect any points in this section.

19 Whether sentence (1) also has the reading represented by (2) is not relevant to our discussion here. If that reading is altogether absent (and not just less prominent), that is problematic for our overall theory. But let us set this potential problem aside here. What we are concerned with in this section is how to predict the other reading we have described, which is no doubt available.

20 See chapter 5, where we introduced these conditions.


22 See P. Jacobson's "i-within-i Effects in a Variable-Free Semantics and a Categorial Syntax," in P. Dekker and M. Stokhof (eds), Proceedings of the Ninth Amsterdam Colloquium (Amsterdam, University of Amsterdam Institute for Logic, Language, and Computation, 1994), pp. 349–68, for an interesting study of pronoun binding within DP. Jacobson's evidence actually supports the view that NPs (in contrast to APs and PPs) do not have subjects, and therefore there might be a problem with our analysis of (9).

23 May, Logical Form, p. 69.


25 Once we have (4), it is a routine exercise (which we leave to the reader) to define an appropriate meaning for the determiner every itself which combines with the standard meaning of basket to yield (4).

26 This argument does not appear in the literature, as far as we know. Only Rooth, "Association with Focus," and Larson, "Quantifying into NP," have tried to show that there are possible readings of English sentences which can only be generated by quantifying into DP, and their examples involve DPs in the object position of intensional verbs like need, want, and look for. Their argument relies on Montague's analysis for these verbs, which is beyond the scope of this book.

27 We are distorting the syntax of neither ... nor here, by treating it as a unit which appears between the two phrases it coordinates (see LF (5')). A better analysis would probably treat it as negation affixed to either ... or. These refinements shouldn't affect our present point.

We are disregarding any readings on which every city takes scope internal to the NP person from every city. We are also disregarding any readings on which from every city is an argument or modifier of the verb spy rather than a part of the some DP (as in I spied on him from the window).

Huang, "Logical Relations in Chinese," pp. 228–35, drew from examples similar to (8) a very different conclusion. He hypothesized that relative scope always had to be disambiguated at S structure. More precisely, he claimed that if one DP c-commanded another at SS, it also had to c-command it at LF. At first sight, this principle appears to be violated by the existence of wide-scope object readings in sentences like (i).

(i) Two politicians spy on every citizen.

(i) can mean that for every citizen there are two politicians who spy on him/her. In the LF of this reading, every citizen is not c-commanded by two politicians, but at SS, it apparently is. Huang explained away this apparent counterexample to his principle by maintaining that (i) had (at least) two different SSs, including one in which the DP every citizen was extraposed, i.e., right-adjoined to S. He proposed that this SS was the one that went with to the wide-scope object reading.

We cannot discuss the merits of this general proposal here, but Huang's application to examples like (8) is interesting. If nothing extraposes at SS in (8), Huang predicts the reading in (a). (Note that the some DP and the every DP do not stand in a c-command relation one way or the other, so their relative scopes with respect to each other are unaffected by Huang's principle.) If the object of spy on extraposes (as in (i)), we get an SS in which the some DP c-commands the subject, so this forces reading (b). For (c) to be generated, we would have to extrapose the DP every city outside the c-command domain of the subject, without at the same time extraposing the containing some DP so high that it c-commands the subject.
9 Bound and Referential Pronouns and Ellipsis

We have had a semantics for bound-variable pronouns since chapter 5, but up to now we have ignored the referring uses of pronouns. In this chapter, we introduce a treatment of referring pronouns and attend to some consequences of the fact that every occurrence of a pronoun gives rise to a potential bound-free ambiguity. In particular, we will look at the interpretation of ellipsis constructions and show that the so-called strict-sloppy ambiguity in an elided phrase reduces to the ambiguity of pronouns in its antecedent. To set the stage for this argument, we will have to include a rudimentary sketch of a theory of ellipsis.

9.1 Referential pronouns as free variables

9.1.1 Deictic versus anaphoric, referential versus bound-variable pronouns

Traditional grammar distinguishes between “deictic” and “anaphoric” uses of personal pronouns.\(^1\) The terms are typically explained as follows: A pronoun is used \emph{deictically} when it receives its reference from the extralinguistic utterance context, and it is used \emph{anaphorically} when it “picks up its reference” from another phrase in the surrounding text. Paradigm cases of deictic uses are the demonstrative uses: that is, those accompanied by a pointing gesture which bears the main burden of fixing the reference. But the term also applies to cases where the intended referent is sufficiently salient in the utterance context even without the help of pointing. For instance, (1) might be uttered immediately after a certain man has left the room.\(^2\)

(1) I am glad he is gone.
Under the right circumstances, this utterance is felicitous, and he unambiguously refers to the man who just left. No previous reference to the same person need have been made; nor need the referent be physically present and available to be pointed at. This too is classified as a “deictic” use.

It is not clear, however, whether the traditional division into deictic and anaphoric uses has any role to play in linguistic theory. In the more recent tradition of formal linguistics, both generative syntacticians and philosophers of language have frequently advocated the view that a certain subset of the anaphoric uses does not differ in any theoretically relevant way from the deictic uses. The idea is that anaphora may often be viewed as reference to a contextually salient individual as well. It seems to differ from deixis only insofar as the cause of the referent’s salience is concerned. For instance, when the he in an utterance of (2) refers to Smith, this may be attributed to the fact that Smith has just been referred to in the previous sentence and that this has made him salient to the audience.

(2) I don’t think anybody here is interested in Smith’s work. He should not be invited.

The mechanism by which this anaphoric pronoun acquires its reference is not really different from the mechanism by which a typical deictic pronoun (as in (1)) does. Anaphoric and deictic uses seem to be special cases of the same phenomenon: the pronoun refers to an individual which, for whatever reason, is highly salient at the moment when the pronoun is processed. There are many possible reasons why a particular individual might have become salient. Sometimes it is due to an act of pointing by the speaker; at other times it is because the individual in question has just been mentioned (by name or description) by the (current or a previous) speaker; and yet other times its salience is caused by circumstances which were not created by the speaker’s linguistic or non-linguistic behavior at all.

Let us assume, then, that all deictic pronouns and also many anaphoric ones are interpreted by the same general strategy. In disambiguating the pronoun’s reference, listeners assign it to the most salient individual that allows them to make sense of the utterance. This may not be the most salient individual in absolute terms, if that would lead to a reading for the utterance which conflicts with basic common sense and with the speaker’s manifest beliefs and intentions. In such a case, the most salient individual will be passed over in favor of the next most salient one until a plausible overall reading is found. How this works exactly, and how salience interacts with various other factors, is not easy to say, and there is much room here for psycholinguistic research. As semanticists, we abstract away from the strategies of reference resolution and the conditions they require to succeed. We take for granted that where they do apply successfully,
the pronoun denotes a unique individual. In this respect, there is no difference between anaphora and deixis.

Can all examples of anaphoric pronouns be subsumed under this characterization? No. We already know that this is not feasible, for the simple reason that some pronouns don't refer to an individual at all (hence, \textit{a fortiori}, not to the most salient individual that contributes to a sensible overall reading). We have already seen many instances of such \textit{nonreferring} pronouns: for instance, in \textit{such that} clauses (chapter 5) and with quantifying DPs as antecedents (chapter 7). An example of the latter sort is (3), on the reading that this sentence receives most readily when uttered in isolation.

(3) Every man put a screen in front of him.

The pronoun \textit{him} here doesn't refer to an individual any more than its antecedent \textit{every man} does. Rather, it is best treated as a \textit{bound variable}. In the theory we have developed in the last few chapters, this means that it is co-indexed with the QR trace of its antecedent and interpreted by the following rule (repeated from chapter 5).

(4) \textbf{Pronouns and Traces Rule}

\[
\text{If } \alpha \text{ is a pronoun or trace, } i \text{ is an index, and } g \text{ is a variable assignment whose domain includes } i, \text{ then } \llbracket \alpha_i \rrbracket^g = g(i).
\]

Cases like (3) conclusively show that not all anaphoric pronouns can be treated as referential. At least some of them are bound variables.

Would it be possible, on the other hand, to attempt a unification in the opposite direction and treat all anaphoric pronouns as bound variables? Should we, for example, reconsider our analysis of (2) above and attempt, instead, to analyze the \textit{he} in (2) as a bound variable as well? That would require us to posit an LF for (2) in which the DP \textit{Smith} has been raised high enough to c-command the entire two-sentence text. We assume that such LF representations cannot be generated, due to general constraints on movement. If so, then cases of intersentential anaphora must always involve co-reference rather than variable binding. A similar argument can be made about certain cases of \textit{intrasentential} anaphora as well: namely, those where the antecedent is very deeply embedded. Take (5), on the reading where \textit{her} is anaphoric to \textit{Mary}.

(5) Most accidents that Mary reported were caused by her cat.

In order to apply a bound variable analysis to this pronoun, we would have to raise \textit{Mary} all the way from inside the relative clause to the edge of the matrix clause. Assuming that such movement is blocked by island constraints, we have
here another example in which the relevant anaphoric relation must be co-reference rather than variable binding. We conclude, therefore, that at least some “anaphoric” pronouns are best analyzed as referring pronouns, just like the “deictic” pronouns.

We have used the term “co-reference” in opposition to “variable binding” here, and we want to draw your attention to our terminological policy in this regard. Much of the syntactic literature uses “co-reference” in an informal sense that covers bound-variable anaphora along with other semantic relations. “Co-reference” there is used much like the traditional term “anaphora”. A broad descriptive term like this can be useful in order to have short labels for the different readings of a given sample sentence and to indicate quickly which one of them we are talking about. We will use “anaphora” (and also “antecedent”) in this way below. But when we say “co-reference”, we always mean it in a narrow literal sense: two expressions (or occurrences of expressions) co-refer iff they refer to the same individual. It follows that if two expressions co-refer, then each of them refers to something. A bound-variable pronoun therefore cannot possibly co-refer with any expression. Co-reference implies reference. We cannot legislate this terminology to the general linguistic community, but for the purposes of working with this textbook, it is important that you adopt it too and maintain it consistently. Many points we will be making may otherwise be lost on you.

In summary, the descriptive category of “anaphoric” uses of pronouns appears to fall into two semantically rather different groups: bound-variable uses and (co-)referring uses. The traditional taxonomy “anaphoric” versus “deictic” disregards this important semantic distinction. Instead, it focuses on a subdivision within the class of referring uses, which seems to be more relevant to the theory of language use (processing) than to the theory of grammar (semantics and syntax).

9.1.2 Utterance contexts and variable assignments

We have yet to say how referring pronouns are represented in our LFs and treated by our semantic rules. The simplest assumption we can make at this point is that all pronouns have the same internal syntax and semantics. They must all bear an index (numerical subscript) at LF to be interpretable, and they are all interpreted by the same rule, Traces and Pronouns. The only thing that distinguishes referring pronouns from bound-variable pronouns is that they happen to be free variables. In other words, the difference between referential and bound-variable pronouns resides in the larger surrounding LF structure, not in the pronouns themselves.
Treating referring pronouns as free variables implies a new way of looking at the role of variable assignments. Until now we have assumed that an LF whose truth-value varied from one assignment to the next could ipso facto not represent a felicitous, complete utterance. We will no longer make this assumption. Instead, let us think of assignments as representing the contribution of the utterance situation. The physical and psychological circumstances that prevail when an LF is processed will (if the utterance is felicitous) determine an assignment to all the free variables occurring in this LF. Let's implement this formally.

If you utter a sentence like

(6) She is taller than she

then your utterance is felicitous only if the utterance situation provides values for the two occurrences of the pronoun “she”. Given that referring pronouns bear indices at LF, (6) has some representation such as (7),

(7) She$_1$ is taller than she$_2$

and we can think of an utterance situation as fixing a certain partial function from indices to individuals. An appropriate utterance situation for LF (7) is one that fixes values for the indices 1 and 2. That is, it is appropriate for (7) only if the variable assignment it determines includes 1 and 2 in its domain.

Let “c” stand for an utterance situation or “(utterance) context” (we use these terms interchangeably), and let “gc” stand for the variable assignment determined by c (if any). We can thus formulate the following appropriateness and truth-conditions for LFs with free pronouns.

(8) **Appropriateness Condition**
A context c is *appropriate* for an LF $\phi$ only if c determines a variable assignment $g_c$ whose domain includes every index which has a free occurrence$^6$ in $\phi$.

(9) **Truth and Falsity Conditions for Utterances**
If $\phi$ is uttered in c and c is appropriate for $\phi$, then the utterance of $\phi$ in c is *true* if $[\phi]^{gc} = 1$ and *false* if $[\phi]^{gc} = 0$.

For instance, suppose that (6) with LF (7) is uttered in a situation $c_i$ which furnishes the assignment $g_{c_i}$:

$$g_{c_i} = \begin{bmatrix} 1 \rightarrow & \text{Kim} \\ 2 \rightarrow & \text{Sandy} \end{bmatrix}$$
c₁ is appropriate for (7), and thus this is a true utterance if Kim is taller than Sandy, and a false utterance if she isn't.

What about the role which gender, person, and number features of referential pronouns play in their interpretation? For instance, if Kim or Sandy is male, then the context c₁ that we just defined is not intuitively appropriate for the use of LF (7). Since both pronouns in (7) are third person feminine singular, they both must refer to female individuals distinct from the speaker and the addressee. Can we capture this intuition as well?

As it turns out, we already make the correct predictions, at least as regards the gender features. In section 5.5, when we briefly addressed the role of features in constraining the possible indexings for bound pronouns, we proposed in passing a presuppositional account of gender features. Here is a more explicit version. Suppose that features are nodes of their own, adjoined to the DP (at least, that the semantic composition principles treat them this way). For instance, we may have structures like:

```
DP
  /\[third person\]
DP
  /\[feminine\]
DP
  /\[singular\]
she₁
```

The lowest DP-node is interpreted by the Pronouns and Traces rule, the higher ones by Functional Application. Each feature has a suitable lexical entry, for example:

(10) \[
[feminine] = \lambda x : x \text{ is female} . x
\]

So a feature denotes a partial identity function. If the DP-node above such a feature gets a denotation at all, it will be the same as the one of the next lower DP. But if the lower DP's denotation fails to have the appropriate property (for example, femaleness), the one above gets no value.

Consider now what happens if, for example, the LF (7) is uttered in a context c₂ which maps the index 1 to a man. (That is, g_{c₂}(1) is male.) Although this context may qualify as “appropriate” for (7) in the technical sense of our Appropriateness Condition (8) (provided that g_{c₂} also has index 2 in its domain),
it will not provide a referent for the complete DP that corresponds to the
pronoun "she_1". Instead, this DP will fail to be in the domain of \( [[ \]]^{\text{she}} \), due to (10).
Therefore, the whole sentence (7) will also not be in the domain of \( [[ \]]^{\text{she}} \), and
the utterance of (7) in \( c_2 \) is neither true nor false. In short, the result is presup-
position failure.

On this account of gender features, it is not strictly speaking impossible to use
a feminine pronoun to refer to a man. But if one does so, one thereby expresses
the presupposition that this man is female. This is intuitively right. If the dis-
course participants mistakenly believe a male referent to be female, or if they are
willing to pretend that they do, then indeed an occurrence of she can refer to
a man, without any violation of principles of grammar.

Person and number features might be treated in an analogous fashion. You
may add appropriate lexical entries analogous to (10).

9.2 Co-reference or binding?

We have seen some clear cases of referring pronouns and also some cases that
compellingly called for a bound-variable analysis. Whenever an anaphoric pro-
noun has a quantifier as its antecedent, for instance, it cannot be referential,
at least not on the reading which we intend when we say the pronoun is
"anaphoric". Anaphoric pronouns with quantifier antecedents are the paradigm
cases of bound-variable pronouns, but they are by no means the only instances
of bound-variable pronouns that we find in natural language. Bound-variable
readings for pronouns are not confined to sentences involving quantifiers. It is
easy to see that our theory predicts them to have a much wider distribution.
And we will argue that this prediction is welcome and supported by empirical
evidence.

Exercise

The following sentence contains a pronoun (his), but no quantifier. (Assume
our Fregean analysis of the, on which definite DPs have type e.)

(i) The dog that greeted his master was fed.

(a) Give a precise paraphrase of the presupposition and assertion that
(i) expresses on its most salient reading. (Imagine an out-of-the-blue
utterance, with no contextually salient candidate for the reference of *his*.)

(b) Show that this reading can be correctly predicted by means of a bound variable construal of *his*.

(c) Argue that no referential analysis of *his* can adequately capture this same reading.

As we noted in section 9.1.1, many anaphoric pronouns with referring antecedents such as proper names are best analyzed as referring pronouns that co-refer with their antecedents. When the antecedent was in a separate sentence, or deeply embedded in an island, this co-reference analysis was virtually forced on us. But that still leaves many cases where anaphora to a referring antecedent *can* be analyzed as variable binding. Consider (1).

(1) John hates his father.\(^7\)

Disregard any reading where his refers to someone other than John and concentrate on the reading which is true iff John hates his own (John’s) father.

There is nothing that prevents us from generating the following LF for (1).

(2)

```
(2)  S
    /\   S
   /   /
  DP  1 DP VP
   |    |
  John DP VP
       /\   /
      |
      t1 V DP
      |
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free indices), and then we apply the semantic rules and the Truth and Falsity Conditions for utterances to derive that any utterance of (2) is true if John hates John's father and false if he doesn't.

So the relevant ("anaphoric") reading of (1) is adequately captured by a bound-variable analysis of the pronoun his.

It is just as easy, of course, to generate LFs in which the pronoun is free, either by not QR'ing John at all (see (3)), or by giving its trace an index different from the pronoun's (see (4)).

(3)
```
S
  ||
DP  VP
  ||
John V DP
  ||
hates the NP
  ||
  DP N
  he₁ father
```

(4)
```
S
  ||
DP  2 S
  ||
John DP VP
  ||
t₂ V DP
  ||
hates the NP
  ||
  DP N
  he₁ father
```

In (3) and (4), he₁ is free. By the Appropriateness Condition, these LFs thus require an utterance context which assigns a referent to the index 1. Among
the candidates are contexts like $c_1$ such that $g_{c_1} = [1 \to \text{Fred}]$, or $c_2$ such that $g_{c_2} = [1 \to \text{Bill}]$, or $c_3$ such that $g_{c_3} = [1 \to \text{John}]$, etcetera. If it is $c_3$ that prevails when (3) or (4) is uttered — and why shouldn’t it be, given that John has been brought to the listener’s attention by the mention of his name? — then the utterance will be true iff John hates John’s father (as we can easily calculate).

So the pertinent reading of (1) can be just as adequately captured by a referential analysis of the pronoun. We have shown that LF (3) (or (4)) uttered in context $c_3$ has precisely the same truth and falsity conditions as LF (2) (uttered in any context).

As this example illustrates, our current theory predicts a lot of “invisible” ambiguity. We have what intuitively is just one reading, with certain definite truth-conditions, but it is generated in two distinct ways: with distinct LFs, distinct denotations for many of the subsentential constituents, and distinct appropriateness conditions for the utterance context. Both analyses are consistent with our current assumptions about syntax, semantics, and pragmatics. We have no motivation for any particular constraint whose addition to the theory would render one of them unavailable. This being so, we hypothesize that both analyses are indeed grammatical. But it seems to be impossible in principle to obtain empirical confirmation for this hypothesis. All the evidence we have to go on are our intuitions about the truth-conditions of utterances of sentences like (1). And, as we have just seen, these intuitions are equally consistent with (1) being ambiguous at LF between (2) and (3) and with it unambiguously having structure (3).8

Or so it appears. In fact, the conclusion we have just stated has been challenged. A number of authors (including Partee, Keenan, Lasnik, Sag, Williams, and Reinhart9) have argued that there is truth-conditional evidence for this “invisible” ambiguity, after all. But we must look beyond the truth-conditions of sentence (1) by itself to find it. When we consider the truth-conditions of larger units (complex sentences or multisentential texts) in which (1) is followed by an elliptical continuation, we will see that the ambiguity becomes manifest, and we will get evidence that both a bound-variable reading and a co-referential reading are generated by the grammar. We reproduce this line of argument in the remainder of the chapter.

9.3 Pronouns in the theory of ellipsis

9.3.1 Background: the LF Identity Condition on ellipsis

The most discussed construction in the semantic literature on ellipsis is “VP ellipsis”, exemplified in (1) and (2).
(1) He smokes. He shouldn't.

(2) Laura took a nap, and Lena did too.

The second sentence in each text is missing a VP on the surface, but it is understood just as if there was one present ([\textit{v}_\text{p}, \text{smoke}] in (1), [\textit{v}_\text{p}, \text{take a nap}] in (2)). Here we will talk about VP ellipsis as well as a somewhat different construction known as "stripping" or "bare argument ellipsis". The latter is illustrated by (3)-(5).

(3) Some people smoke, but not many.

(4) Laura left Texas, and Lena as well.

(5) Laura drank the milk last night, or perhaps the juice.

The main difference between the two constructions is that in bare argument ellipsis the auxiliary is absent as well. Apart from negation and adverbs like "as well", merely a bare noun phrase remains. Note also that (unlike in VP ellipsis) the remnant phrase need not be the subject, as seen in (5).

We assume that in the derivation of all these sentences, some constituent is deleted on the way from SS to PF. In VP ellipsis, it's a VP, and in bare argument ellipsis, an S. If bare argument ellipsis always deletes an S constituent, we must assume that the "remnant" argument always has been topicalized (adjoined to S) before the deletion. These assumptions, of course, would need to be justified if we were to attempt a serious syntactic analysis of the construction.

Elliptical sentences are thus incomplete as sentences on the surface, but they nevertheless have the semantic interpretation of complete sentences. How can this be? A straightforward answer is that they are complete sentences at the level of LF. The deletion on the PF branch does not take place in the derivation from SS to LF. This way, the LFs can be interpreted by the familiar semantic rules and will wind up with run-of-the-mill sentence meanings. For instance, if Lena as well has essentially the same LF as the complete sentence Lena left Texas as well, then we don't have to say anything further about how it receives its interpretation.

The deletion operation that yields elliptical sentences is evidently not allowed to delete arbitrary material in arbitrary environments. If it were, then the text in (4) would be predicted to have many readings in addition to the one we observe. For instance, it should be able to mean that Laura left Texas and Lena also \textit{moved}. This reading could result from deleting \textit{moved} instead of \textit{left Texas}.

But (4) cannot be understood in this way. We clearly judge (4) false, for instance, if Lena moved only from San Antonio to Dallas. Apparently, the material
that is deleted in the derivation of an elliptical sentence must be identical to material that is present overtly in the antecedent discourse. Roughly, the explanation for why the covert portion of the LF of Lena as well in (4) can be left Texas, but cannot be moved, has to be that the former expression, but not the latter, matches the overt material in the antecedent sentence Laura left Texas.

How exactly should this identity condition be stated, and at what level does it apply? A main insight of work on ellipsis in the 1970s was that the relevant level had to be one which was fully disambiguated with respect to interpretation. In our framework, this is LF. Mere identity of representations at some pre-LF stage of the derivation is not sufficient to license ellipsis. To appreciate this point, we must look at examples in which non-trivial changes occur in the derivation towards LF. (The ones above are too simple.) Take a sentence in which LF movement has an essential disambiguating role to play, such as (6).

\[(6) \quad \text{Laura showed a drawing to every teacher.}\]

(6) displays a scope ambiguity and can mean either that there was a drawing that Laura showed to every teacher, or else that for every teacher there was a drawing that Laura showed her. The two readings have distinct LFs, but the same structure at SS. Now examine your intuitions about possible meanings for the text in (7), in which (6) is followed by an elliptical sentence.

\[(7) \quad \text{Laura showed a drawing to every teacher, but Lena didn't.}\]

The elliptical conjunct “Lena didn’t” here can be understood to say that there wasn’t any drawing that Lena showed to every teacher – but only if the preceding sentence is understood analogously: namely, as saying that there was a drawing that Laura showed to every teacher. Alternatively, “Lena didn’t” can mean that not for every teacher was there a drawing that Lena showed her – but only when the preceding sentence means that for every teacher there was a drawing that Laura showed her. It is not possible to read the first conjunct with one scope order, and the elliptical sentence with the reverse one. This is readily explained if the elided material has to have the same LF as its antecedent, but it would not automatically follow if they merely needed to look alike at some pre-LF stage in the derivation, say SS.

We thus adopt the following general condition on ellipsis:

\[(8) \quad \text{\textit{LF Identity Condition on Ellipsis}}\]

A constituent may be deleted at PF only if it is a copy of another constituent at LF.

Before we return to the topic of pronouns, here are a few sample derivations. The first one involves VP ellipsis, the second one bare argument ellipsis with an
object remnant. The constituent to be deleted at PF is always the one that is boxed.

(9) SS:

In the derivation to LF, everything can stay the same in this case. The VP in the first conjunct is a copy of the deleted VP in the second conjunct, and thus it serves to fulfill the LF Identity Condition.

The case of bare argument ellipsis is slightly more complicated. In the SS, the remnant has already been fronted.

(10) SS for (5):

In deriving the LF, we may now QR the milk, leaving a trace also indexed 1.

(11)

Now we have an LF-identical antecedent for the deleted S-node.
9.3.2 Referential pronouns and ellipsis

What happens when the antecedents of elided phrases contain pronouns? Let us begin with examples of pronouns which are clearly referential. For instance, (12) has a natural reading on which it means that Philipp went to Roman's office.

(12) (On Roman's birthday), Philipp went to his office.

Imagine that an utterance of (12), with the reading just described, is continued as follows.

(13) (On Roman's birthday), Philipp went to his office. Marcel didn't.

What can the continuation mean? The fact is that, given that we have understood the first sentence to mean that Philipp went to Roman’s office, we must understand the second one to mean that Marcel didn’t go to Roman’s office. It cannot then mean that Marcel didn’t go to Philipp’s office, or didn’t go to Felix’s office, or didn’t go to Marcel’s office. However salient one of these other individuals may be, the only choice for the referent of the elided copy of his is Roman, the referent of the overt copy of his.

We would like our theory of ellipsis to make this prediction. Let us see what it predicts as it stands. On the reading we have specified for the first half of (13) (=(12)), its LF must contain a free pronoun. Let's say it looks as in (14), and the utterance context c is such that g_c(1) = Roman.

(14)

```
S
 /    \\
|     |
Philipp I

I
 /    \\
|     |
PAST go to his_1 office
```

Given this first half, what options do we have for the LF of the whole text? The LF Identity Condition requires that the deleted VP in the second part be a copy of the antecedent VP. So this has to be go to his_1 office as well, and not, for instance, go to his_2 office. We predict that (15) is a grammatical LF for (13), but (16) an ungrammatical one.

(15) an ungrammatical one.

(16) an ungrammatical one.
In the grammatical LF (15), the two occurrences of \text{his}_1 cannot fail to co-refer. Given our context \(c\), the second S in (15) is true iff Marcel didn’t go to Roman’s office. The ungrammatical LF (16), on the other hand, could express one of the undesirable readings we seek to exclude. For example, if \(c\) happened to be such that \(g_c(2) = \text{Felix}\), then the second S in (16) as uttered in \(c\) would be true iff Marcel didn’t go to Felix’s office. It is a good thing, then, that the LF Identity Condition excludes (16).

Unfortunately, however, there is another possible LF which we do not yet exclude and which also expresses one of the unavailable readings. Suppose we had chosen to QR Marcel and to give its trace the index 1, as in (17).

\begin{align*}
(15) & \\
S & S \\
\text{Philipp} & \text{Marcel} \\
\text{I} & \text{I} \\
\text{I} & \text{I} \\
\text{VP} & \text{VP} \\
\text{PAST go to his}_1 \text{ office} & \text{didn’t go to his}_1 \text{ office} \\

(16)* & \\
S & S \\
\text{Philipp} & \text{Marcel} \\
\text{I} & \text{I} \\
\text{I} & \text{I} \\
\text{VP} & \text{VP} \\
\text{PAST go to his}_1 \text{ office} & \text{didn’t go to his}_{2} \text{ office} \\

(17) & \\
S & S \\
\text{Philipp} & \text{Marcel} \\
\text{I} & \text{I} \\
\text{I} & \text{I} \\
\text{VP} & \text{VP} \\
\text{PAST go to his}_1 \text{ office} & \text{t}_1 \text{ didn’t go to his}_1 \text{ office} \\
\end{align*}
No principle that we know of excludes this LF, and in particular, it does not violate our LF Identity Condition. But look at its interpretation. Given our context $c$, the first part still says that Philipp went to Roman’s office, but the second part says that Marcel didn’t go to Marcel’s office. If this LF is generated for our text in (13), we are in trouble.

To close this “loophole,” we need to tighten up the theory somehow. How exactly this should be done has been a matter of considerable debate in the literature on ellipsis. For our purposes in this book, a very preliminary remedy will have to do. What appears to be causing the trouble in (17) is that we have a free pronoun in the first sentence but a bound-variable pronoun in the analogous place in the second sentence. Perhaps we should rewrite the LF Identity Condition in such a way that it is sensitive to this difference, even though it is not a difference that can be detected by merely looking inside the LF representations of the deleted phrase and its putative antecedent. Another option is to leave the LF Identity Condition as it stands, but add a general prohibition against LFs in which a given index has both bound and free occurrences. We might as well choose the second option here.

(18) No LF representation (for a sentence or multisentential text) must contain both bound occurrences and free occurrences of the same index.

Given (18), LFs like (17) are simply not generated, whether the derivation involves PF deletion or not: (Note that the adoption of (18) makes no difference to anything other than our predictions about ellipsis. The reason for this is that every LF prohibited by (18) has “notational variants” that are semantically equivalent with it, as well as indistinguishable from it in all respects relevant to syntactic principles other than the identity condition on ellipsis.) With (17) as a source of unwanted readings out of the way, our LF Identity Condition predicts just what we were aiming for. Whenever the first sentence in an utterance of (13) means that Philipp went to Roman’s office, the elliptical second sentence can only mean that Marcel didn’t go to Roman’s office. More generally, we predict that whenever a pronoun in the antecedent phrase refers to an individual $x$, then that pronoun’s counterpart in the deleted phrase must refer to $x$ as well. Referential pronouns keep their reference under ellipsis. This law follows from our current theory.

9.3.3 The “sloppy identity” puzzle and its solution

Against the background of the law we just deduced, the phenomenon of so-called sloppy identity readings can look very puzzling. Consider (13) once more.
(13) (On Roman’s birthday), Philipp went to his office. Marcel didn’t.

In the previous section, you were instructed to read the first sentence as “Philipp went to Roman’s office”, and we did not consider other interpretations. But, of course, it is also possible to read this first sentence as claiming that Philipp went to his own (Philipp’s) office, for instance. In this case, what are the possible readings for the continuation “Marcel didn’t”? Clearly, it is possible then to understand “Marcel didn’t” as claiming that Marcel didn’t go to his own (Marcel’s) office. (This is even the most salient reading.) The text (13) as a whole can mean that Philipp went to Philipp’s office and Marcel didn’t go to Marcel’s office (the so-called sloppy reading). Doesn’t this constitute a blatant counterexample to the law we have just seen? If the antecedent VP means “go to Philipp’s office”, how can the deleted VP mean “go to Marcel’s office” and yet have an identical LF? When the antecedent VP meant “go to Roman’s office”, the deleted VP had to mean “go to Roman’s office”. No “sloppy” alternative was available then. But when the antecedent VP means “go to Philipp’s office”, the deleted VP doesn’t necessarily mean “go to Philipp’s office”! Why this difference between Roman and Philipp? 15

The puzzle persists only as long as we take it for granted that the overt pronoun has to be referential. 16 It is certainly tempting to think that whenever Philipp went to his office has the truth-conditions of “Philipp went to Philipp’s office”, it must be because his refers to Philipp. But we ought to know better. As we saw earlier in this chapter, the truth-conditions of “Philipp went to Philipp’s office” can also come about in an entirely different way: namely, by interpreting his as a bound-variable pronoun. Let us examine this alternative possibility and how it affects the predictions of the LF Identity Condition.

Here is an LF for the text (13) in which the overt as well as the deleted occurrence of the pronoun have been construed as bound variables. In both sentences, the subjects were QR’d, and the same indices were chosen for their traces.

(19)
The LF in (19) expresses precisely the so-called sloppy reading of the text in (13). As we can easily calculate by our semantic rules, both trees in (19) are true iff Philipp went to Philipp's office and Marcel didn't go to Marcel's office. (19) also meets the LF Identity Condition on ellipsis, since the two VPs are identical. And it is not in violation of stipulation (18), since all occurrences of index 1 are bound in the overall representation. (19), then, shows that the "sloppy" identity reading is not, after all, a counterexample to our theory. It only seemed to be a counterexample when we reasoned on the basis of the prejudice that the pronouns here were all referring.

We have solved the sloppy identity puzzle and vindicated our assumptions about ellipsis. In doing so, we have found an indirect, yet compelling, argument for the coexistence of bound-variable readings and referential readings in anaphoric pronouns with referring antecedents. This is the argument which we announced at the end of section 9.2, where it had seemed at first as if the bound-variable reading could not possibly be told apart from a co-referential reading.

There is more to be done in order to establish these conclusions firmly. Ideally, we would like to see that our assumptions about pronouns and our assumptions about ellipsis conspire to yield exactly the right predictions about the distribution of strict and sloppy readings in all sorts of examples. As a first small step in this direction, let's make sure that our simple text (13) receives all and only the readings that it intuitively allows. We have just seen how the "sloppy" reading is accounted for, and we already saw in section 9.3.2 how to generate a "strict" reading. By a "strict" reading, we mean any reading with truth-conditions of the form "Philipp went to x's office, and Marcel didn't go to y's office", for some given person x. x may be Philipp, or Marcel, or a third person such as Roman. The LF in (15) above can represent any one of these readings, given an appropriate utterance context which maps the free index shared by both pronouns to the relevant person x. This is as it should be, since all such strict readings are available. The harder job is to show that we generate no unwanted sloppy readings: that is, no truth-conditions of the form "Philipp went to x's office and Marcel didn't go to y's office", where x ≠ y, except in the one case where x = Philipp and y = Marcel.

We reason as follows. Assume there are x ≠ y such that a given utterance of (13) has the truth-conditions of "Philipp went to x's office and Marcel didn't go to y's office". Now suppose, first, that x ≠ Philipp. Then the meaning "Philipp went to x's office" for the first conjunct cannot have come about through a bound-variable construal of his, but must be due to a referential construal. Hence the overt pronoun in the antecedent must be a free variable, and, by (8) and (18), its counterpart in the deleted predicate must be another free variable with the same index. Then, both pronouns refer to x, and the second conjunct means that Marcel didn't go to x's office. Given that x ≠ y, this contradicts our
initial assumption. Second, suppose that \( y \neq \text{Marcel} \). By analogous reasoning, it follows that the elided pronoun must be a free variable, and therefore that its overt counterpart must be equally free and co-indexed. Hence the first conjunct must mean that Philipp went to \( y \)'s office – again contrary to assumption. In sum, if either \( x \neq \text{Philipp} \) or \( y \neq \text{Marcel} \), we derive a contradiction with our initial assumption that the text has "sloppy" truth conditions (that is, \( x \neq y \)). This proves that the case where \( x = \text{Philipp} \) and \( y = \text{Marcel} \) constitutes the only option for a grammatical "sloppy" reading.

In this section, we have analyzed a simple example of pronoun interpretation in an ellipsis construction. Our main aim has been to argue that the distinction between bound-variable anaphora and co-reference, which was forced upon us by our general analysis of pronouns, has an important role to play in explaining the emergence of "strict" and "sloppy" readings under ellipsis. If the approach we have outlined is correct, it also follows that ellipsis data can give us valuable indirect evidence about the constraints that syntax places on the LF configurations in which pronouns are bound. We will return to this point in the following chapter.

Exercise

Disregarding any readings where the italicized pronoun refers to a person not mentioned in the sentence, each of the following sentences is ambiguous.

(a) Only Bill can get his car started.
(b) Only Mary knows why she is crying.
(c) Only Ed understood the objections to his theory.
(d) Only he was asked a question that he understood.

Your first task is to describe what the ambiguity consists in and how the two readings differ in truth-conditions. Then choose one example and explain in detail how the two readings are accounted for under our current theory. This will involve specifying the relevant properties of the LFs for each of the readings and showing how our semantic rules apply to these LFs.

To do all this, you will have to assume a semantic value for the word only. Assume that this word combines with DPs of type \( e \) (at least in the present examples), but that the DPs of the form \([o\_\text{only} \_D]\) that it forms with them are of type \( \langle e,t\rangle,t\rangle \). Assume the following lexical entry:

\[
\text{[only]} = \lambda x \in D_e . \ \lambda f \in D_{\langle e,t\rangle} . \ \text{for all } y \in D_e \text{ such that } y \neq x, f(y) = 0.
\]
You are permitted — in fact, encouraged — to abstract away from irrelevant complexity by taking the semantic values of certain complex LF chunks for granted instead of deriving them compositionally. For example, you may make assumptions like the following without argument:

$$[[s_t, \text{knows why she}_j \text{is crying}]]^g = 1 \text{ iff } g(i) \text{ knows why } g(j) \text{ is crying.}$$

Notes

5. The assumptions we are making in this paragraph about the limitations on LF movement will be scrutinized more thoroughly in chapters 10 and 11 below. But, to anticipate the results, they will hold up well enough.
6. Recall the definition of “free occurrence” of an index in a tree from chapter 5.
7. As before when we have needed an *ad hoc* account of possessive constructions, we assume that the structure of the DP is something like \[ [\_D] \text{ the } [\_NP \text{ he father}] \], with a covert definite article. Father is a transitive noun (type \(<e, <e,t>\) ), and he saturates its internal argument.
8. The possibility that only (2) and not (3) is grammatical is excluded by the fact that his *can refer* to someone other than John. Such a reading clearly requires structure (3).
11. Even further unavailable readings would be generated if we disregarded the “as well”, which places some constraints of its own on the missing material. Very roughly, as well (and its synonyms also and too) give rise to a presupposition which is only fulfilled if the predicate they co-occur with denotes a superset of some salient predicate in the previous discourse. For instance, Laura left Texas, and Lena moved
too is okay, because \{x : x moved\} is a superset of \{x : x left Texas\}; whereas Laura left Texas, and Lena drank milk too is weird, because \{x : x drank milk\} cannot be taken for granted to be a superset of \{x : x left Texas\}. We will not undertake a serious analysis of these particles here. We would like to stress, however, that not all examples of ellipsis include such a particle. See, e.g., (1), (3), (5). Therefore, the problem of overgeneration from an unconstrained deletion operation is actually even worse than it looks for (4). For instance, why shouldn't the elliptical disjunct in (5) be able to mean “... or perhaps Lena spilled the juice”?

12 An excellent summary of the arguments for this conclusion is found in Sag, *Deletion*, ch. 2. Sag also provides a useful overview of the literature up to 1976.

13 See Williams, “Discourse,” and especially Sag, *Deletion*, for a broader survey of problem cases of this kind and for more general versions of our stipulation and alternatives to it. Sag, in fact, argued at length for a version of the LF Identity Condition which was much more restrictive than the combination of our (8) and (18). It was so restrictive, in fact, that it ruled out many of the derivations we present in this chapter, including (11) and (19). This, in turn, led Sag to adopt a different LF syntax for VPs. A serious introduction to the theory of ellipsis would have to include a thorough discussion of Sag's proposal and the complex collection of data that has been brought to bear on it. This would go far beyond the scope of an introductory semantics text. See R. Fiengo and R. May, *Indices and Identity* (Cambridge, Mass., MIT Press, 1994), for a useful recent overview.

14 For instance, (17) has variants like (i).

(i)

```
\begin{center}
\begin{tikzpicture}
  \node {S} child {node {I} child {node {PAST} child {node {Philipp}} child {node {VP} edge from parent node [left] {go to his\textsubscript{1} office}}} edge from parent node [left] {IVP}} child {node {I} child {node {t\textsubscript{2}} child {node {Marcel}} child {node {VP} edge from parent node [left] {didn't go to his\textsubscript{2} office}}} edge from parent node [left] {IVP}};
\end{tikzpicture}
\end{center}
```

(17) violates (18), (i) abides by it. There is no difference in interpretation, and no difference with respect to the syntactic mechanisms which generate and license LFs – except, of course, for the LF Identity Condition, which doesn't license deletion in (i).

15 Presumably, it was this way of looking at the phenomenon that originally gave rise to the label “sloppy” (due to Ross – see J. R. Ross, “Constraints on Variables in Syntax” (Ph.D. dissertation, MIT, 1967), distributed by Indiana Linguistics Club). If the LF Identity Condition were *strictly* applied, it seems, such readings would be ruled out. But, as we will see right away, there is a different and more insightful way of looking at the phenomenon, which reveals that the “sloppy” reading does not really deserve to be called “sloppy” at all. It actually obeys the LF Identity Condition just as strictly as the “strict” reading. Nevertheless, the descriptive labels “sloppy” and “strict” have stuck, despite the overwhelming success of the theory that contradicted them, and so we are using them too.

16 Among the first authors to see this clearly were Partee (“Opacity”) and Keenan (“Names, Quantifiers”).
In the sections to come, we will first look at syntactic co-indexing procedures and conditions on binding. This will be short, since the details belong in a syntax text. We will then distinguish syntactic and semantic notions of "binding" and state a principle connecting the two. The outcome will be essentially Reinhart's theory of pronoun interpretation, and we will review its predictions about "sloppy identity" readings and the so-called Weak Crossover effect, as well as its implications for the syntax of co-reference.

10.1 Indexing and Surface Structure binding

We have argued that the input for the semantic interpretation component should be LF representations. LF representations are transformationally derived from SS representations by Quantifier Raising (QR). We have adopted a very unconstrained view of QR: it is optional; it can apply to any kind of DP; and it can adjoin to any node. Its application is indirectly constrained by syntactic principles on the one hand and by interpretability on the other.

QR, like all movement, leaves a co-indexed trace. Where do the indices come from? One possibility is that they are added in the application of movement itself, so that even an initially index-less phrase would be co-indexed with its trace after movement. Another possibility is that the phrase to be moved had an index to begin with, which then gets duplicated in the movement. What we have done so far in this book was compatible with either option, but in this chapter, we take the second view: all indices are already present by SS.

Suppose that indices are optionally inserted with any DP at DS and may be chosen freely at this point. When a DP αi moves in the course of the derivation, a trace t_i is left in its original site, and the index of the moved DP is also adjoined to its sister node at the landing site (to form a predicate abstract). So the configuration after movement looks like this:
Though indices are assigned freely, certain choices will later be filtered out. One obvious constraint is that every *pronoun* must be generated with an index, or else it will not be interpretable. For non-pronominal DPs, indices are in principle optional, but will, in effect, be required if the DP is going to move. Otherwise, there will be no index on its trace, and the trace will be uninterpretable. Another way in which the semantics indirectly constrains indexing possibilities has to do with which features like gender, person, and number. Recall from chapters 5 and 9 that these induce certain presuppositions, which will systematically fail if we try to co-index DPs with conflicting features.

Among syntactic constraints on indexing possibilities, we have mentioned (in chapter 5) the Prohibition against Vacuous Binding and, briefly, the so-called Binding Theory. It is the latter’s effects which we want to look at more closely now.

To formulate the Binding Theory, we need a definition of *syntactic binding*, which in turn presupposes a definition of *c-command*. Here are two more or less standard definitions.

(2) **C-command**
- A node \( \alpha \) c-commands a node \( \beta \) iff
  - (i) neither node dominates the other, and
  - (ii) the first branching node dominating \( \alpha \) dominates \( \beta \).

(3) **Syntactic binding**
- A node \( \alpha \) syntactically binds a node \( \beta \) iff
  - (i) \( \alpha \) and \( \beta \) are co-indexed,
  - (ii) \( \alpha \) c-commands \( \beta \),
  - (iii) \( \alpha \) is in an A-position, and
  - (iv) \( \alpha \) does not c-command any other node which also is co-indexed with \( \beta \), c-commands \( \beta \), and is in an A-position.

"A-positions" are the positions of subjects and objects: "non-A (A-bar) positions" are adjoined and complementizer positions. The Binding Theory only applies to binders in A-positions. The following version is from Chomsky.
(4) Binding Theory
(A) Reflexive pronouns and reciprocals must be bound in their Minimal Governing Category.
(B) All other pronouns must be free in their Minimal Governing Category.
(C) All non-pronominal DPs must be free.

"Bound" here is to be understood as "syntactically bound" and "free" as "not syntactically bound". Clauses (A) and (B) give conditions for binding possibilities of different kinds of pronouns. The conditions crucially rely on the notion "Minimal Governing Category". There is much discussion of this in the syntactic literature. For our purposes, the exact definition of the local domain in which reciprocals and reflexives must be bound and other pronouns must be free is not essential. What is essential, however, is that reflexives must be bound, ordinary pronouns can be bound or free, and non-pronominal DPs must be free.

10.2 Syntactic binding, semantic binding, and the Binding Principle

Syntactic binding is not the same as binding in the semantic sense. We discussed the semantic notion in section 5.4. Our discussion yielded the following definition.

(1) Let $\alpha^m$ be a variable occurrence in a tree $\gamma$ which is (semantically) bound in $\gamma$, and let $\beta^n$ be a variable binder occurrence in $\gamma$. Then $\beta^n$ (semantically) binds $\alpha^m$ iff the sister of $\beta^n$ is the largest subtree of $\gamma$ in which $\alpha^m$ is (semantically) free.

Recall also that the only variable binders in our system are the adjoined indices that trigger the Predicate Abstraction Rule.

Compare the following two structures:

(2) $S$
    $\quad$ every diver$_1$
    $\quad$ VP
    $\quad$ defended himself$_1$
The pronoun “himself₁” is syntactically bound in both (2) and (3), but semantically bound only in (3). Accordingly, our semantic rules give very different results when applied to these two structures. In (2), “himself₁” is interpreted like a free variable; it receives its value from the variable assignment supplied by the utterance context (if any). The index on “every diver” receives no interpretation at all. It is not seen by any of our semantic rules, and we assume it is just skipped over. (2) can only mean that every diver defended a certain contextually salient individual. In (3), on the other hand, the pronoun “himself₁” is interpreted like a bound variable. It is bound in the interpretation of the predicate abstract, and the predicted meaning for (3) is that every diver defended himself.

(3) correctly captures the truth-conditions of the English sentence “Every diver defended himself”; (2) does not. We have to find out, then, what excludes (2) as a possible LF for this sentence. In other words, we have to give an explanation for why QR is obligatory in structures like this, even though it is an optional rule and there is no type mismatch forcing us to move the quantifier phrase.

We propose that what is at play here is a principle which enforces a tight connection between syntactic binding at SS and semantic binding at LF. Roughly speaking, every syntactic binding relation must correspond to a semantic binding relation, and vice versa. To give a precise statement, we must first introduce a derivative notion of “semantic binding”, which relates two DPs. (On the literal notion, only variable binders in the semantic sense can bind anything.)

\[ (4) \text{ A DP } \alpha \text{ semantically binds a DP } \beta \text{ (in the derivative sense) iff } \beta \text{ and the trace of } \alpha \text{ are (semantically) bound by the same variable binder.} \]

In this derivative sense, we can say that “every diver₁” in (3) “semantically binds” “himself₁”, even though the real binder is the adjoined index right below it.

We can now state our principle.
(5) **Binding Principle**

Let \( \alpha \) and \( \beta \) be DPs, where \( \beta \) is not phonetically empty. Then \( \alpha \) binds \( \beta \) syntactically at SS iff \( \alpha \) binds \( \beta \) semantically at LF.

(The qualification "not empty" serves to exempt traces.)

The Binding Principle in interaction with the Binding Theory excludes (2) as a possible LF, even though it is a possible SS, of course. The Binding Theory requires syntactic binding between quantifier phrase and reflexive at SS. Because of the Binding Principle, this must be matched by semantic binding at LF. This means that QR *must* apply in this case, hence the correct interpretation of the sentence "Every diver defended himself" is guaranteed.

Let us look at another example.

(6) She\(_1\) said that she\(_1\) saw a shark next to her\(_1\).

In (6), we have two syntactic binding relations: the first pronoun binds the second, and the second binds the third. The Binding Principle requires that each of these syntactic binding relations corresponds to a semantic binding relation. This means that all but the lowest pronoun must QR, and the LF looks as follows:

\[
(6') \quad \begin{array}{c}
S \\
\downarrow \\
\text{she}_1 \\
\downarrow \\
1 \\
\downarrow \\
said \\
\downarrow \\
that \\
\downarrow \\
she_1 \\
\downarrow \\
1 \\
\downarrow \\
S \\
\downarrow \\
t_1 \text{ saw a shark next to her}_1 \\
\end{array}
\]

The Binding Principle imposes a direct correspondence between syntactic binding at Surface Structure and variable binding at Logical Form. Whenever you find syntactic binding of a pronoun at SS, you have a bound-variable interpretation at LF. And whenever you have a bound variable interpretation at LF, you have syntactic binding at SS. The second subclaim is a version of the so-called Weak Crossover condition. The first subclaim is surprising and controversial. Is the bound-variable interpretation really the only way to interpret syntactic binding of pronouns? We will see that there are reasons for a positive answer. They are laid out in detail in Reinhart's book, and we will review them below.
10.3 Weak Crossover

Here is an example of what is known as "Weak Crossover".

(1) The shark next to him attacked every diver.

In (1), the pronoun him cannot receive a bound-variable interpretation. But now consider the following indexing.

(2) The shark next to him\textsubscript{1} attacked [every diver]\textsubscript{1}.

In (2), the pronoun him\textsubscript{1} and the quantifier phrase every diver\textsubscript{1} are co-indexed. Yet neither DP binds the other syntactically, since neither c-commands the other. According to the Binding Principle, this means that neither is allowed to bind the other semantically at LF; hence the following structure is not a possible LF for (2).

(3) 

```
(3) S
    ____________
   /     \
 every diver\textsubscript{1} 1 S
    \     /
       DP     VP
        the NP     attacked t\textsubscript{1}
             shark PP
                     next to him\textsubscript{1}
```

In (3), him\textsubscript{1} is semantically bound, but it was not syntactically bound at SS, in violation of the Binding Principle. We have no choice as far as QR goes, however. The object is of semantic type $\langle e,t\rangle,t\rangle$, hence the type mismatch forces QR. The conclusion, then, is that the Binding Principle indirectly rules out the co-indexing between object and pronoun in (2).\textsuperscript{6} A legitimate indexing for the same sentence is (4).

(4) The shark next to him\textsubscript{2} attacked [every diver]\textsubscript{1}.
The LF representation derivable from (4) presents no problem for the Binding Principle.

(5)

\[
\begin{array}{c}
S \\
\text{every diver}_1 \\
1 \\
\text{S} \\
\text{DP} \\
\text{the} \\
\text{NP} \\
\text{shark} \\
\text{PP} \\
\text{next to} \\
\text{him}_2
\end{array}
\]

In (5), the pronoun him\textsubscript{2} is not semantically bound. If the utterance context doesn't specify a value for it, the sentence winds up without a truth-value, hence is inappropriate. The Binding Principle, then, correctly predicts that the pronoun in (1) must be referential, and cannot be understood as a bound variable.

The Weak Crossover facts confirm one part of the Binding Principle: namely, that bound-variable interpretations of pronouns at LF require syntactic binding at SS. The following section provides more evidence for the same claim as well as for the other part: namely, that syntactic binding forces a bound-variable interpretation.

10.4 The Binding Principle and strict and sloppy identity

Our discussion of ellipsis (section 9.3) gave us evidence that certain pronouns anaphorically related to proper names are truly ambiguous between a co-referential and a bound-variable interpretation. We saw this in examples like (1).

(1) First Philipp cleaned his room, and then Felix (did).

If the pronoun "his" is a free variable, (1) receives a "strict" reading; if it is a bound variable, we obtain the "sloppy" reading.
The Binding Principle which we have posited in this chapter now states that the bound-variable interpretation is grammatical in exactly those cases where we also have syntactic binding at Surface Structure. Given the connection between bound-variable interpretation and sloppy identity, this implies further predictions about the distribution of strict and sloppy readings. In the configuration of example (1), the syntactic Binding Theory permits, but does not force, co-indexing between “Philipp” and “his” at SS. The emergence of both sloppy and strict readings is thus precisely what we expect in this case. Let us look at the predictions for some other cases.7

Since syntactic binding always requires c-command, a pronoun which is anaphorically related to a non-c-commanding antecedent can only be an instance of co-reference anaphora, but never of bound-variable anaphora. We have already seen that therefore the antecedent can never be a quantifier (see section 10.3). It also follows that when the antecedent is a proper name and there is an elliptical continuation, a sloppy reading is systematically unavailable.

(2) His father spoils Roman, but not Felix.

(2) cannot mean that Roman’s father spoils Roman, but Felix’s father doesn’t spoil Felix.

Further test cases for this prediction are the following examples:

(3) (a) Zelda bought Siegfried a present on his birthday, and Felix too (that is, “... and she bought Felix ... too”).
(b) Zelda thought about Siegfried on his birthday, and about Felix too.

(4) (a) You can keep Rosa in her room for the whole afternoon, but not Zelda.
(b) Felix is kissing Rosa in her favorite picture but not Zelda (that is, “he is not kissing Zelda”).
(c) Rosa is wearing a pink dress in her favorite picture, but not Zelda.

Lack of c-command is one possible obstacle to syntactic binding, but not the only one. The Binding Theory (clause (B)) also implies that a non-reflexive pronoun may not be bound when it is “too close” to a potential antecedent, as in *“Bill1 likes him1”. In this case, too, our theory predicts that only co-reference anaphora may obtain, never bound-variable anaphora. For reasons we will turn to in the last section, even co-reference at first seems impossible here. But there are special discourse conditions under which it is not entirely excluded, and these provide us with another kind of test case for the Binding Principle.9 Consider (5).10
I know what Bill and Mary have in common. Mary likes Bill and Bill likes him too.

(5) illustrates that there is a possible context in which “Bill likes him” has the truth-conditions of “Bill likes Bill”. Assuming that the Binding Theory is correct, this configuration precludes co-indexing, and if the Binding Principle is right, the observed truth-conditions can therefore not be the result of construing “him” as a bound variable. This can only be a case of co-reference. It follows, then, that elliptical continuations will unambiguously have a strict reading. The prediction is borne out:

(6) Ann: “Nobody likes Bill.”
Mary: “No, that’s not quite right. Bill likes him, of course, just like you do.”

Mary’s utterance in (6) cannot have a sloppy reading. It can’t mean that Bill likes Bill, just like you like you.

Thus far, we have looked at implications from one direction of the Binding Principle. If there is a bound-variable interpretation, there has to be syntactic binding at Surface Structure. The other part of the biconditional says that whenever there is syntactic binding at SS, there must be semantic binding at LF. This part, too, implies predictions about the distribution of strict and sloppy readings. Let’s look at cases where the syntax doesn’t just allow, but requires, binding. Reflexive pronouns must be syntactically bound according to clause (A) of the Binding Theory. The Binding Principle therefore predicts them to be bound variables in all cases. Co-reference is never an option, and we expect the absence of a “strict” reading in examples like (7).

(7) Ann hurt herself, and Mary (did) too.

Let us finally note (again with Reinhart) that parallel examples with only can be constructed for all of the ellipsis examples in this section. In these examples, capitalization represents focal stress.

(8) Only Philipp cleaned his room.

(9) I only claimed that his father spoils ROMAN.

(10) (a) Zelda only bought SIEGFRIED a present on his birthday.
    (b) Zelda only though about SIEGFRIED on his birthday.
(11) (a) You can only keep ROSA in her room for the whole afternoon.
    (b) Felix is only kissing ROSA in her favorite picture.
    (c) Only Rosa is wearing a pink dress in her favorite picture.

(12) Despite the big fuss about Felix’s candidacy, when we counted the votes, we found out that in fact only Felix himself voted for him.

(13) Only Ann hurt herself.\textsuperscript{13}

We leave it to the reader to deduce (and examine) the relevant predictions about possible readings.

\section*{10.5 Syntactic constraints on co-reference?}

The Binding Principle stipulates a tight correspondence between syntactic binding at SS and bound-variable interpretation at LF, and we have seen ample support for this. Our theory also implies that syntactic binding or the absence thereof has nothing to do with allowing or prohibiting co-reference. If a DP is not syntactically bound, it is not interpreted as a bound variable. But how else it is interpreted in that case depends solely on its lexical content (if it is a non-pronominal DP) or on the assignment contributed by the utterance situation (if it is a pronoun). There is nothing in the theory which predicts that any of the following examples should be ungrammatical or require disjoint reference between the italicized DPs.

(1) (a) \textit{My father} voted for \textit{my father}.
    (b) I hurt \textit{me}.
    (c) \textit{Bill Clinton} overestimates \textit{him}.
    (d) \textit{She} sold \textit{Ann’s car}.

Bound-variable interpretations are precluded, but co-reference is allowed by the theory (provided suitable utterance contexts), and in some cases even necessitated by the DPs’ inherent meanings (including feature-related presuppositions).

For some examples (though admittedly not the ones that come to mind at first), this permissiveness is welcome. For instance, we don’t have to worry about sentences like (2).

(2) She is the boss.
Presumably, clause (C) of the Binding Theory prohibits co-indexation of *she* and the *boss* in (2). But there is clearly no obstacle to co-reference between the two DPs. On our current approach, (2) receives a truth-value only in contexts that fix a reference for *she*. The referent of *she* must have been mentioned earlier, or else must be picked out with a gesture, or be salient otherwise. If the context does indeed specify a referent for *she* in (2), then (2) is true just in case this referent is the boss. And this account seems right.

The point is not limited to identity sentences like (2). For another classic example, imagine yourself at a Halloween party, trying to guess which of your friends is hiding behind which costume. One of the guests just left the party, and somebody utters (3).

(3) He put on John’s coat.

The pronoun “he” refers to the one who just left, and the utterance is just as grammatical and felicitous if this is John as it is if it’s somebody else. Moreover, there is the type of example which we already saw in the section on sloppy identity:

(4) I know what Bill and Mary have in common. Mary likes Bill and Bill likes him too.

In sum, there is some evidence which indicates that it would actually be misguided to search for a theory on which Binding Theory directly constrains reference. Our present approach makes some good predictions.

Still, there are problems. Notwithstanding the existence of examples like (2), (3), and (4), most examples with c-commanded names and locally c-commanded pronouns are clearly judged to disallow co-reference. Except for the identity sentence in (2), we had to set up rather special discourse contexts to bring out the judgment that co-reference was possible. The problem is highlighted by pairs like (5) and (6).

(5) *She* liked the flowers that we bought for *Zelda*.

(6) The flowers that we bought for *her* pleased *Zelda*.

(6), in which the pronoun precedes but does not c-command the name, is readily accepted with a co-referential reading. (5) clearly has a different intuitive status. Once the pronoun c-commands the name, only a much narrower set of conceivable utterance contexts will make co-reference acceptable.

Reinhart offers an explanation for this difference. She proposes a further principle, which establishes a preference ranking between different LFs that
convey the same meaning. In a nutshell, Reinhart’s principle says that if a given message can be conveyed by two minimally different LFs of which one involves variable binding where the other has co-reference, then the variable binding structure is always the preferred one. A precise formulation of the principle is not trivial, and we will content ourselves here with a few elementary and informal illustrations.

Suppose a speaker wants to convey the information that Felix voted for himself. That is, she wants to produce an utterance which is true if Felix voted for Felix and false if he didn’t. There are various different possible utterances which have these truth-conditions. The following pairs of LFs and reference assignments (where needed) represent a few of these choices:

(7) Felix voted for Felix
(8) Felix voted for him1 gc = [1 → Felix]  
(9) Felix 1 [t1 voted for himself1]
(10) He1 voted for Felix gc = [1 → Felix]

In (7), (8), and (10), the object co-refers with the subject, but in (9), it is a bound variable. Reinhart’s principle thus dictates that the speaker choose (9) over (7), (8), or (10). It thus accounts for the fact that speakers will not use the surface sentences corresponding to (7), (8), or (10) if they mean to express these particular truth-conditions.

We have presented the example from the speaker’s perspective, but we can equally well look at it from the hearer’s. Whereas the speaker needed to make a choice between different ways of encoding a given truth-condition, the hearer’s task is to disambiguate a given surface string. Suppose the hearer hears “Felix voted for him” and needs to guess the LF structure and reference assignment that the speaker intends him to recover. Among the candidates might be the following pairs.

(8) Felix voted for him1 gc = [1 → Felix]
(11) Felix voted for him1 gc = [1 → Max]
(12) Felix voted for him1 gc = [1 → Oscar]

Now the hearer is taking into account that the speaker is guided by Reinhart’s principle, and thus he must reason as follows: “If she intended (8), she would
be expressing the information that Felix voted for Felix. But in order to convey *this* message, she could also have chosen a different utterance, the one in (9). And by Reinhart’s principle, she would have preferred that choice. I therefore conclude that she must not have intended (8), but perhaps (11) or (12).” So Reinhart’s principle also explains why a hearer will not assign a co-referential reading to the surface form “Felix voted for him”.

Similar reasoning applies to other examples. For instance, (5) (repeated from above) must not be used with “she” referring to Zelda,

(5) *She* liked the flowers that we bought for *Zelda*

because the meaning thereby expressed would be the same as that of a competing LF which involves variable binding: namely, (13).

(13) *Zelda* \[ t_1 \text{ liked the flowers that we bought for } \text{her}_1 \]

Reinhart’s principle therefore predicts that the binding structure in (13) must be chosen over a co-reference use of (5). It also correctly predicts that (6) is different.

(6) The flowers that we bought for *her* pleased *Zelda*.

(6) on the reading where “she” refers to Zelda does not have any “competitor” that expresses the same meaning through variable binding.\(^{17}\) If (14) or (15) were grammatical, they would qualify, but these LFs cannot be generated without violating the Binding Principle. (See section 10.3.)

(14) *Zelda* \[ \text{the flowers that we bought for } t_1 \text{ pleased her}_1 \]

(15) *Zelda* \[ \text{the flowers that we bought for her}_1 \text{ pleased } t_1 \]

Reinhart’s approach to the connection between syntactic Binding Theory and the interpretation of referential DPs seems roundabout and complicated at first, and it uses concepts and procedures that need to be made precise for it to yield correct predictions. As it has turned out, however, this is not a weakness, but relates directly to its greatest strengths. In particular, the approach leads naturally to an interesting and plausible hypothesis about the so-called “exceptional” cases like (2), (3), and (4) above, in which co-reference is allowed.\(^{18}\) Let us illustrate with example (2).

What we want to explain is why co-reference is okay in (2); that is, why the speaker is permitted to choose the following LF plus reference assignment.

(16) *She*\(_1\) is the boss \[ g_c := [1 \rightarrow \text{the boss}] \]
Why isn’t she required to forgo this choice in favor of the following alternative, which involves variable binding instead of co-reference?

(17) The boss [t₁ is herself₁]  

Reinhart’s answer is that (17) would not convey the same information as (16). This is what distinguishes example (2) from the ones above (for example, “Felix voted for him”). The information that the speaker intended to convey in our story about “Felix voted for him” was adequately expressed by the bound-variable structure “Felix₁ voted for himself₁.” But the information that the speaker wants to convey when she says “She is the boss” is not preserved in the reformulation “The boss is herself”. The latter is completely trivial, and only attributes to the boss the necessary property of self-identity. The original utterance, by contrast, can clearly serve to tell the hearer something new (and possibly false). Reinhart’s principle regulates only choices between possible utterances which convey a given intended message. It thus implies that (17) is not in competition with (16), and therefore leads to the correct prediction that the co-reference in (16) is acceptable.

Incidentally, the account just sketched raises an interesting question about the relevant concept of “same information” or “intended message”. The difference between “She is the boss” (with co-reference) and “the boss is (identical to) herself” is intuitively real enough, but it turns out that it cannot be pinned down as a difference between the truth-conditions of utterances in the technical sense of our theory. As a simple calculation reveals, if the LF she₁ is the boss is uttered in a context c such that gc₁(1) = the boss, then this utterance has precisely the same truth-conditions as an utterance of the LF the boss [t₁ is herself₁] (uttered in any context). Both utterances are true non-contingently, since the boss = the boss, of course. It appears that there is something wrong here with our conception of “truth-condition of an utterance,” or at least that this is not the appropriate theoretical construct to explicate the semantic (or pragmatic) notions that are relevant to the operation of Reinhart’s principle. Unfortunately, we cannot offer a better explication here, since it would require a serious theory of context dependency.

Reinhart’s account of examples (1) and (3) follows the same general strategy. In each of these cases, there is a clear intuition that the speaker would not be conveying exactly her intended message if she replaced the utterance involving co-reference by one involving variable binding. This being so, the preference principle licenses the structure with co-reference. As in the case of (2), it is beyond our current means to analyze precisely what it is that would get lost in the replacement. But for the purposes of the present argument (and Reinhart’s), this doesn’t matter. We have shown that Reinhart’s theory explains not only why Binding Theory appears to constrain co-reference is most cases, but also why it
sometimes doesn’t. This vindicates the view that co-reference, as opposed to variable binding, is not directly represented in linguistic representations, and thus cannot be directly affected by syntactic rules.

10.6 Summary

We started out with a problem and ended up with Reinhart’s theory of anaphora. The problem arose when we wondered what it is that forces anaphoric interpretation in cases like (1).

(1) Ann1 defended herself1

Our semantic rules for proper names and pronouns just don’t allow us to establish a connection between the two noun phrases in (1). The interpretation of nouns doesn’t depend on the index on the noun at all. And the interpretation of pronouns is strictly local. It depends on the index of the pronoun, but it can’t see whether this pronoun is co-indexed with another node or not. We saw that we could guarantee an anaphoric interpretation for (1) only if we had reasons to require that the noun phrase Ann1 must undergo QR. This led to one half of the Binding Principle: namely, that whenever there is syntactic binding, there must be semantic binding. The preceding sections gave support for this generalization, as well as for the other half of the Binding Principle: namely, that whenever there is semantic binding, there must be syntactic binding (a version of the “Weak Crossover Condition”). We also defended the view that neither the Binding Theory nor the Binding Principle make any mention of co-reference.

Notes


2 This looks a little different from the structures we have been drawing so far (see section 7.3). We previously did not include an index in the moved phrase itself. Most of the time, the index in α, will in fact be semantically vacuous, and thus we might as well leave it out. The exceptions are cases in which the moved phrase is itself a variable (e.g., a pronoun).

3 We are not assuming here that the index is part of the pronoun as a lexical item. If pronouns are listed in the lexicon at all, they are listed there without an index and as semantically vacuous items. No semantic rule sees the pronoun itself, and the
lowest node that is interpreted is the indexed DP dominating it. This implies that our definition of “variable” in section 5.4 needs a small correction. As it is written, only terminal nodes can be variables. Instead of “terminal node”, we should have said “node that dominates at most one non-vacuous item”. This allows the indexed DP dominating a pronoun to count as a variable, but still excludes truly complex constituents containing meaningful material in addition to a (single) variable.

4 See N. Chomsky, Lectures on Government and Binding (Dordrecht, Foris, 1981). What we are defining under (3) is actually called “local A-binding” there. But since this is the only syntactic notion of binding we will be talking about, we can afford to omit “local” and “A” from its name.

5 Ibid.

6 This is not quite correct if there are nodes of type t below the subject, which can serve as alternative landing sites for QR (see ch. 8). On that assumption, the SS in (2) will not be ruled out. But note that QR’ing the object to a site below the subject would not result in binding the pronoun him. It would remain a free variable, whose interpretation is unrelated to that of the object’s trace (despite co-indexing), and which can only get a referent from the context. So the empirical prediction remains the same: (1) cannot mean that every diver x was attacked by the shark next to x.

7 This will be a superficial and incomplete survey. See Reinhart’s work for a detailed defense of the predicted generalizations, and much recent work for critical discussion.

8 Taken more or less from Reinhart, Anaphora, p. 153.

9 See ibid., p. 169 for the argument to follow.


11 A wider range of examples with reflexives reveals that the data are more complex. For recent discussion and controversy, see A. Hestvik, “Reflexives and Ellipsis,” Natural Language Semantics, 3/2 (1995), pp. 211–37, and Fiengo and May, Indices and Identity, among others.

12 In many of these examples, only is a VP operator that associates with a focused element in the VP. See M. Rooth, “Association with Focus” (Ph.D. dissertation, University of Massachusetts, Amherst, 1985, distributed by GLSA) for ways of interpreting such constructions.


14 This one is due to Higginbotham (J. Higginbotham, “Anaphora and GB: Some Preliminary Remarks,” Cahiers Linguistiques d’Ottawa, 9 (1980), pp. 223–36). See also Evans, “Pronouns.”

15 Reinhart’s principle is sometimes called a “pragmatic” principle and has been related to general maxims of conversation like “Avoid ambiguity”. (See S. Levinson, Pragmatics (Cambridge, Cambridge University Press, 1983).) It is tempting, indeed, to speculate that the preference for bound-variable anaphora (“Felix voted for himself”) over co-reference (“Felix voted for him”) is due simply to the fact that the reflexive disambiguates the utterance, whereas a plain pronoun creates referential ambiguity. We do not want to endorse a reduction to pragmatics of this simple-minded sort. For one thing, many applications of the principle as stated by Reinhart involve rankings among structures which all contain plain non-reflexive pronouns and thus have equally ambiguous surface forms. Also, the appeal to a general strategy of ambiguity avoidance opens the door to many objections. For example,
why do we allow co-reference in the coordinate sentence “Felix is smart and he is nice” when there is the less ambiguous alternative “Felix is smart and nice”? (Keenan, personal communication). Given that referential ambiguity is easily tolerated in this case, doesn’t this undermine the whole idea behind Reinhart’s proposal?

We do not mean to discourage attempts to reduce Reinhart’s principle to something more general. Clearly, it should not simply be accepted as an isolated stipulation, and perhaps it really does turn out to follow from “pragmatics” in some sense of the word. But in our experience, billing the principle as a “pragmatic” one sometimes misleads people about what it actually says, and leads them to dismiss it without appreciating the considerable empirical support for it.


This example shows that the set of potential “competitors” must be limited to LFs with essentially the same overall structure, in a sense to be made precise. We do not want (6) to have to compete with (i), for example:

(i) Zelda 1[t₁ was pleased by the flowers that we bought for her₁]

(i) is truth-conditionally equivalent with the co-referential reading of (6), and it is a grammatical structure involving variable binding. Nevertheless, its existence does not preempt the co-referential use of (6). The success of Reinhart’s principle thus depends on an appropriate characterization of the set of structures that are being compared in its application. (See references in n. 16.)

Another strong argument for the approach has emerged from the study of language acquisition. Young children typically go through a stage in which they allow co-reference in linguistic structures in which adults do not. An insightful explanation of the children’s performance has been based on the hypothesis that they have adult-like competence in the Binding Theory but an immature ability to apply Reinhart’s preference principle. See Y.-C. Chien and K. Wexler, “Children’s Knowledge of Locality Conditions in Binding as Evidence for the Modularity of Syntax and Pragmatics,” Language Acquisition, 1 (1991), pp. 225–95; J. Grimshaw and S. T. Rosen, “Knowledge and Obedience: The Developmental Status of the Binding Theory,” Linguistic Inquiry, 21 (1990), pp. 187–222; and Grodzinsky and Reinhart, “Innateness,” for details and discussion.

Refer to the definition in section 9.1.2.

11 E-Type Anaphora

We have distinguished referential from bound-variable pronouns and have discussed the relevance of the distinction for the semantics of ellipsis and the interpretation of syntactic Binding Theory. We have argued, in particular, that "anaphoric" relations in the descriptive sense of the term are a mixed bag, and that some cases involve variable binding whereas others involve co-reference. In this chapter, we will show that there are even further kinds of anaphoric relations to be distinguished, and additional interpretations for pronouns as well.

We will begin with a review of predictions from the current theory, specifically its predictions about the distribution of pronouns which are anaphorically related to quantificational antecedents. We will find that the syntactic assumptions that gave rise to these predictions hold up well under scrutiny. Yet, they seem at first to exclude a whole bunch of perfectly grammatical examples. This problem will lead us to reexamine some assumptions: Is it possible, after all, to assign a referential interpretation to a pronoun whose antecedent is a quantifier? What does "antecedent" mean in this case? In the end, we will find conclusive evidence that some pronouns are neither referential nor bound variables. They have a third kind of interpretation, as so-called E-Type pronouns. We will spell out one concrete analysis of E-Type pronouns, on which they are definite descriptions with complex silent predicates, and will briefly mention applications to some famous examples.

11.1 Review of some predictions

QR is subject to some locality constraints. For instance, it is not grammatical to QR the "no" DP in the relative clause of (1) all the way to the matrix S.

(1) Every problem that no man could solve was contributed by Mary.

Such constraints on QR entail constraints on the distribution of bound-variable pronouns. If the predicate abstract created in the movement of no man cannot extend beyond the relative clause, then the binder index that heads this abstract
cannot bind pronouns outside the relative clause. In this particular example, there aren’t any pronouns around in the first place. But look at (2) and (3).

(2) Every problem that no man showed to his mother was easy.

(3) Every problem that no man could solve kept him busy all day.

For the his in (2), an anaphoric reading where its antecedent is no man is entirely natural. This is predicted, since we can generate an LF in which the relative clause looks as follows:

(4)

```
CP
  wh₁
  S
    DP  2
      no man  t₂
    VP
      showed  t₁
to  his₂ mother
```

(3), on the other hand, is strange – unless, of course, we imagine a context in which the reference for him is somehow fixed independently of the information contributed by this sentence. This, too, is predicted. If we QR no man within the relative clause, the pronoun remains free (even if it happens to bear the same index):

(5)

```
S
  DP  3
    every
      NP
        problem
      CP
        kept him₂ busy...
    wh₁
    S
      DP  2
        no man  t₂
could  VP
          solve  t₁
```
This LF requires a context that assigns a referent to the index 2. If we wanted to get the pronoun bound, we would have to QR no man to the matrix S:

\[(6)^*\]

```
S
  \_ 2
  \_/ \\
no  man  DP
    \_ 3
    \_/ \\
      every DP
      /\ \\
problem NP
  \_  t3
  \_/ \\
kept him2 VP
    /\ \\
wh1 S
  \_  t2
  \_/ \\
  \_/ \\
  \_/ \\
  \_/ \\
t2 could VP
    /\ \\
solve t1
```

(6) would mean that there is no man who was kept busy all day by every problem he could solve. No such reading is available for (3), and the prohibition against QR out of relative clauses correctly predicts this.

The binding of the pronoun in (6) would probably be ruled out even if the long-distance application of QR to no man were legitimate. The reason is that the trace of no man fails to c-command him, so the co-indexing in (6) also violates (at least some versions of) the Weak Crossover prohibition (for instance, our Binding Principle from chapter 10). LFs like (6) are thus ruled out redundantly. Still, we have independent evidence – namely, from judgments about the truth-conditions of sentences like (1) (where pronoun binding does not play a role at all!) – that one of the reasons why (6) is ruled out is a constraint on QR. (6) would be impossible even if Weak Crossover did not apply here.

Let us look at a couple more structures in which the scope of a quantifier phrase cannot be wide enough to allow the binding of a certain pronoun. (These cases, too, are redundantly excluded by Weak Crossover.)

The simplest case of this kind arises when the pronoun isn’t even in the same sentence as the quantifier. The highest adjunction site that we could possibly choose when we QR a DP is the root node for the whole sentence. It follows from this that bound-variable anaphora is a sentence-internal relation. Indeed, it is impossible to construe no boy as the antecedent of he in the text in (7).
(7) No boy was invited. He complained.

A similar situation arises in coordinate structures. There is some evidence that QR obeys the Coordinate Structure Constraint, so that a quantifying DP in one conjunct of and, for instance, cannot take scope over the whole conjunction. This implies that it is also impossible for no boy to bind the he in the following variant of (7).

(8) No boy was invited and he complained.

Compare (7) with (9).

(9) John was invited. He complained.

(7) doesn't have an anaphoric reading where no boy is the antecedent of he, but (9) is easily read with he anaphoric to John. This contrast is not mysterious from our present point of view. Due to the sentence-internal nature of QR, neither example permits a bound-variable construal for he. Hence both examples can only be read with he as a referring pronoun. In (9), he may refer to the same individual as John, in which case the two DPs co-refer. In (7), he cannot co-refer with no boy for the simple reason that no boy doesn't refer.

So far, so good. But as we look at a wider range of examples, this simple picture will come to look insufficient.

11.2 Referential pronouns with quantifier antecedents

Consider a variant of an example from Evans:

(1) Only one congressman admires Kennedy. He is very junior.

Only one congressman is not a referring DP. If we wanted to treat it as referring to an individual, we would run into many of the familiar paradoxes afflicting such treatments (see chapter 6). Therefore, only one congressman and he could not possibly co-refer. Nor should it be possible for he to be a bound variable co-indexed with the QR trace of only one congressman. Since the two are in separate sentences, there is no possible landing site for the quantifier phrase that would be high enough to c-command the pronoun. Yet our intuitions tell us that he can be read as anaphorically related to the antecedent only one congressman. How can this be, if neither co-reference nor variable binding are possible?
The first response to this dilemma that comes to mind is that QR is perhaps less constrained than we have thought, after all. Perhaps the constraints that we reviewed above for DPs with the determiner no do not pertain to DPs in general. Could it be that some DPs, among them only one congressman, are able to take scope over whole multisentential texts?

We could implement this possibility by introducing a constituent T (for “text”) that dominates all the Ss in a text. A suitable semantic rule for such T constituents is easy to state:

(2) If $\phi^1$, $\ldots$, $\phi^n$ are Ss, and $\psi$ is $[\phi^1, \ldots, \phi^n]$, then $[\psi] = 1$ iff for all $i$ such that $1 \leq i \leq n$, $[\phi^i] = 1$.

In other words, a text is interpreted as the conjunction of all its sentences. In (1), we could now apply QR in such a way that it adjuncts only one congressman to T. And if Weak Crossover for some reason didn’t apply here either, we could also co-index this DP with the pronoun he. The result would be an LF with the pronoun bound:

(3) is interpretable and receives the following truth-conditions: $[\varphi(3)] = 1$ iff there is only one element of D that is a congressman and admires Kennedy and is very junior.

Does this solve our problem of accounting for the anaphoric reading in (1)? No. As Evans pointed out, the truth-conditions of (3) are not the intuitive truth-conditions of the English text (1). To see why not, imagine a situation in which two congressmen admire Kennedy, one of them very junior and the other one senior. The structure (3) is true in this state of affairs, since there is only one individual who is a congressman, admires Kennedy, and is very junior. Intuitively, however, someone who utters (1) under these circumstances has made a false assertion. (1) contains the assertion that only one congressman admires Kennedy; and this assertion is falsified by our scenario.
We conclude from this, with Evans, that our attempt to account for (1) by liberalizing syntactic constraints on QR and pronoun binding was misguided. Upon closer inspection, (1) actually supports these constraints. The fact that (1) cannot have the meaning of (3) confirms that DPs with the determiner only one are subject to these constraints as well, no less than DPs with no.

It's back, then, to our initial assumption that the scope of only one congressman is confined to the first sentence. We might as well forget again about T-nodes and the rule for their interpretation, and assume that the text in (4) is represented at LF as a sequence of $S$s.

\[
\begin{array}{c}
S \\
\quad DP \\
\quad \\
\quad D \\
\quad N \\
\quad only \ one \ congressman \\
\quad 1 \\
\quad S \\
\quad t_1 \\
\quad VP \\
\quad admires \\
\quad Kennedy \\
\quad he_2 \\
\quad VP \\
\quad is \ very \ junior
\end{array}
\]

In (4), $he_2$ is a free variable. (And it would be just as free if we had insisted on writing $he_1$ instead.) So it must be a referential pronoun. Can we argue that this prediction is correct?

To do so, we have to answer three questions: (i) What does the pronoun refer to (on the intended reading)? (ii) Does this assignment of reference lead to intuitively correct truth-conditions for each sentence in the text? (iii) How does the referent in question become salient enough to be available when the pronoun is processed? Regarding the third question, a somewhat vague and sketchy answer will have to do for us here. But concrete and precise answers to the first two are definitely our responsibility as semanticists.

Here is a quick answer to question (i). The pronoun $he_2$ in (4) refers to the congressman who admires Kennedy.

Here is a quick answer to question (ii). The first sentence in (4) is true iff only one congressman admires Kennedy. Given that $he_2$ denotes the congressman who admires Kennedy, the second $S$ in (4) is true iff the congressman who admires Kennedy is very junior. So for both sentences of (4) to be uttered truly, it must be the case that only one congressman admires Kennedy and that the congressman who admires Kennedy is very junior. This prediction accords with intuition.

And here is a quick answer to question (iii). The congressman who admires Kennedy is a salient referent for the pronoun $he_2$, because a hearer who has just processed the first sentence of text (4) will naturally be thinking about this guy.
These are basically the answers we will stick by, although there are some details to be made more precise. We said that the pronoun refers to the congressman who admires Kennedy. But what if there isn’t one? Or if there are two or more? Taking these possibilities into account, we give the following more precise answer to question (i): When the text (4) is uttered on the “anaphoric” reading that is under consideration here, then the pronoun he_2 refers to the congressman who admires Kennedy, if there is exactly one; otherwise it has no referent. In the latter case, the utterance context is not appropriate for (4) (in the technical sense of our Appropriate Condition from chapter 9), and the second sentence in (4) gets no truth-value. A speaker might, of course, be using (4) in the false belief that there is exactly one congressman who admires Kennedy. But in that case, we are claiming, he acts just like a speaker who uses a pronoun demonstratively and points in a direction where he believes there is something to point at, when in fact there isn’t.

So what we are saying about truth-conditions is, more precisely, the following. A speaker who uses (4) makes two consecutive utterances. The first one is true if only one congressman admires Kennedy, and false otherwise. The second one is neither true nor false if the first one was false. But if the first one was true, then the second one is true if the congressman who admires Kennedy is very junior, and false if he is not.

Let’s also take a second look at question (iii). What causes the referent of he_2 (when it has one, and when this is the congressman who admires Kennedy) to be suitably salient? Clearly, the utterance of the sentence that precedes the pronoun (namely, Only one congressman admires Kennedy) plays a crucial role in this. We said that processing that sentence will make the listener think of this man. Why? Because the sentence Only one congressman admires Kennedy is in some sense “about” him. The appropriate notion of “aboutness” here cannot be defined very precisely. Crucially, we don’t want to say that an utterance is only “about” an entity if that entity is the referent of one of its parts. There is no constituent in Only one congressman admires Kennedy which denotes the congressman who admires Kennedy. The “aboutness” relation is more indirect. The sentence, in effect, makes a claim about the cardinality of a certain set (namely, the intersection of congressmen and admirers of Kennedy). When the claim is true, the set is a singleton and uniquely determines its sole member. Whatever the psychological details, it seems reasonable to assume that a listener who has just interpreted this sentence and imagined it to be true is therefore in a state of mind in which he readily guesses that the intended referent of the subsequent he_2 may be the congressman who admires Kennedy.

In a certain sense, then, quantifiers can serve as “antecedents” for referential pronouns. It is no longer mysterious now that there are perfectly grammatical and interpretable anaphoric pronouns whose quantificational antecedents do not c-command them at SS or at LF. In fact, we may have to wonder at this
point why we were able to find any confirmation at all for our earlier expectation that quantifiers could only antecede pronouns they c-command. Why, for example, did the examples we saw in section 11.1 give us that impression?

(5) Every problem that no man could solve kept him busy all day.

(6) No boy was invited and he complained.

In (5) and (6), there is just no anaphoric reading at all. How come? How do these differ from example (1)?

There seems to be a significant difference between no and only one. We can see this in minimal pairs like (6) versus (7), or (1) versus (8).

(7) Only one boy was invited and he complained.

(8) No congressman admires Kennedy. He is very junior.

When we choose no, anaphoric readings become impossible. The reason for this is actually quite transparent, given the story we have told about (1). If (8) had a reading analogous to the one in (1), the pronoun he in (8) should also refer to the congressman who admires Kennedy. But in this case, it is not compatible with the truth of the first sentence that there is any such congressman. Whenever the first sentence of (8) is uttered truthfully, there is no congressman admiring Kennedy which the pronoun could possibly refer to. Nor is there any other person who is systematically brought to the attention of a listener who processes this negative claim (except, of course, Kennedy, who indeed is the only possible referent for “he”, when (8) is presented out of context).

If this approach is right, it is not the presence of no per se that prevents an anaphoric construal, but the meaning of the antecedent discourse as a whole. If the speaker denies the existence of the relevant potential referent in any other way than by using a sentence with no, the effect should be the same. If, on the other hand, a no DP is embedded in an utterance that doesn’t amount to a denial of existence, anaphoric pronouns should become okay again. There is some evidence that these predictions are on the right track:

(9) I seriously doubt that there is a congressman who admires Kennedy. He is very junior.

(10) I can’t believe that I received no mail at all today. You stole it!

“He” in (9) cannot be taken to refer to the congressman who admires Kennedy, whereas “it” in (10) can refer to the mail I received today.
Let us look at a couple of other quantifiers besides "no" and "only one".

(11) Every woman was invited. She was pleased.

(12) A woman was invited. She was pleased.

(11) does not allow any anaphoric reading. Without a context that independently furnishes a referent, it simply sounds incoherent. Can we explain why? Apparently an utterance of “Every woman was invited” strongly conveys the message that there were several women invited. The semantic analysis of “every” that we have assumed in this book does not predict this as part of the truth-conditions, and so we currently don’t have an explanation for this intuition. But the intuition is clearly there, and this is all we need to know for the present purpose. Given that the first sentence in (11) evokes a scenario with multiple invited women, it is not suited to highlighting any individual woman as the intended referent. From the perspective of our approach to the previous examples, the use of “she” in (11) is infelicitous for essentially the same reason as when you utter “she” while pointing at a crowd. Indeed, this analogy is reinforced by the observation that (11) becomes a lot better when “She was pleased” is changed to “They were pleased”.

Example (12) raises more delicate questions, which have been debated extensively. The basic judgment is that an anaphoric reading in (12) is naturally available. Can we analyze this reading in the same way as we analyzed the anaphoric reading in (1) (with “only one”)? That is, can we say that the pronoun either refers to the unique woman who was invited (provided there is one) or else gets no value? Many concrete uses of such sentences clearly fit this analysis. For example, suppose the topic of conversation is this year’s party convention. Who will be the keynote speaker? In this context, if I utter (12),

(12) A woman was invited. She was pleased.

you will understand me as claiming, first, that a woman was invited (to give the keynote address), and second, that the woman who was invited was pleased. If you suspect that my first claim is false (perhaps because they invited a man, or perhaps because no invitation has been issued yet), then you will accordingly assume (just as the analysis predicts) that my use of “she” may not refer to anyone. The possibility that two or more women might have been invited probably doesn’t even cross your mind, given the scenario we have specified. But it is not strictly inconceivable, of course, and we can examine our intuitions about this case as well.

If two women were invited, then my use of “she” did not refer to an individual who was the unique woman invited – that much is clear as a matter of
logic. But might it still have referred to somebody, perhaps to one of the two invited women? Under appropriate circumstances, this seems to be possible. Imagine, for instance, that you happen to find out not only that two women were invited, but also that one of the invitations occurred in my (the speaker of (12)’s) presence, whereas the other one was issued secretly behind my back. In that case, you will conclude that my use of “she” referred to the woman whose invitation I witnessed. (And accordingly, your judgment about the truth of my second utterance will depend on whether *this* woman was pleased.)

This kind of observation suggests that we should not commit ourselves to any simple-minded recipe for fixing the reference of pronouns which are, in some pre-theoretical sense, “anaphorically related” to quantificational antecedents. A large number of cases of this sort conform to a common pattern: the antecedent sentence is of the form “Det A B”, and the pronoun refers to the unique A that is B, if there is one, or else denotes nothing and causes presupposition failure. But it doesn’t *always* have to work in this way. Reference resolution is a complex cognitive task, and we have been aware all along that we could not provide rules for it, even when we were thinking only about run-of-the-mill deictic uses and cases of co-reference anaphora.

What we have learned in this section is that there are anaphoric relations which are neither variable binding nor co-reference. If the approach we have indicated here is right, then the ways in which previous discourse can furnish a referent for a pronoun have to be varied and sometimes rather indirect. The pronoun does not always simply co-refer with its so-called antecedent. Often, the antecedent contributes to the fact that a referent is made available in a much more roundabout way, in concert with the message conveyed by the surrounding sentence, and with miscellaneous information about the speaker’s grounds for her assertions.

11.3 Pronouns that are neither bound variables nor referential

We have concluded that there are referential pronouns which have quantifiers as “antecedents”, in a certain sense. Although we had not thought of this possibility when we first discussed the distinction between referential and bound-variable pronouns, its existence does not really challenge our theory. The Evans example ((1) in section 11.2) seemed at first to create a problem, but we were eventually able to argue that it could be accommodated without revising any of our syntactic or semantic assumptions. The examples we will look at next are
more recalcitrant. They contain pronouns which, for rather obvious reasons, cannot be considered referential and cannot be treated as bound variables.5

There is, in fact, a general recipe for constructing this problematic type of example. Take a sentence that exhibits the kind of anaphoric dependency we saw in Evans’s congressman sentence:

(1) Only one congressman admires Kennedy, and he is very junior.

Introduce another pronoun into it. For instance, substitute a pronoun for the name “Kennedy”:

(2) Only one congressman admires him and he is very junior.

Finally, embed the whole thing in a construction in which this new pronoun is bound. For example:

(3) Every president thought that only one congressman admired him and he was very junior.

We might have hoped that the analysis we developed for the “he” in the simpler original example would carry over to (3). But there is an obstacle. The “he” in (1), we argued, was a referential pronoun. Can the “he” in (3) be referential as well? What would be its referent? The congressman who admired Kennedy? The congressman who admired Johnson? The congressman who admired Reagan? No. None of these guys is the right choice. Whichever of them we take “he” to denote, the truth-conditions we then predict for (3) as a whole turn out different from the ones we intuitively understand. It seems that the denotation of “he” must be allowed to vary with the assignment somehow. But can “he” be a bound-variable pronoun? Then what would it be bound by? The quantifier “only one congressman”? But we have argued at length that such an analysis makes wrong predictions in sentence (1), and the argumentation carries over entirely to (3).

Here is another pair of examples which illustrates the problem. Take (4) first.

(4) John bought just one bottle of wine and served it with the dessert.

(4) is similar to (1) above and can be accounted for along the same lines: it refers to the bottle of wine John bought. This correctly accounts for the intuition that (4) implies that John served the unique bottle of wine he bought with the dessert, and is not equivalent to “there is just one bottle of wine which John bought and served with the dessert”. The latter could be true if John bought, say, five bottles, four of which he served before the dessert and one of which he served with it. But (4) is judged false in this case.

But now look at (5).
(5) Every host bought just one bottle of wine and served it with the dessert.

Here the same account no longer works. The argument against a bound-variable analysis of it carries over from (4) to (5). (5) does not have the truth-conditions of "for every host, there is just one bottle that he bought and served with dessert". Otherwise it could be true if some of the hosts bought two or more bottles. The English sentence (5) is falsified by such circumstances. So (5) no more involves a bound-variable reading of the pronoun than does (4). But it in (5) cannot be referential either. What could its referent possibly be? The bottle of wine that host John bought? If that were its reference, (5) should mean that every host bought just one bottle of wine and then served the one that John bought for dessert. The actual truth-conditions of (5) are not of the form "every host bought just one bottle of wine and served x for dessert", for any x ∈ D.

So there are some pronouns that are neither free variables nor bound variables. Let's call them "E-Type pronouns". Our current theory interprets all pronouns as (bound or free) variables; so it plainly denies their existence. We will have to make some substantive revision. So far, we only have a label and a negative characterization. What are E-Type pronouns? What is their syntactic representation and their semantic interpretation?

11.4 Paraphrases with definite descriptions

We begin with an observation: E-Type pronouns can always be paraphrased by certain definite descriptions. For instance, (5) in section 11.3 on the relevant reading is truth-conditionally equivalent to (1).

(1) Every host bought just one bottle of wine and served the bottle of wine he had bought with the dessert.

Of course, a paraphrase is not yet a semantic analysis. But it may help us find one. Let us first consider in detail how the paraphrase in (1) is interpreted, and return to the original sentence with the pronoun in the following section.

What is the LF of (1)? Evidently, every host has scope over everything else in the sentence and binds the pronoun he. We have also seen that just one bottle of wine does not take scope over the conjunction and. Since and here conjoins two VPs, and the scope of just one bottle of wine is confined to the left conjunct, there has to be an adjunction site for QR at or below the VP level. As discussed in chapter 8, we can solve this problem by assuming VP-internal subjects. The SS of (1) is then something like (2) (making obvious abbreviations).
In deriving an interpretable LF from this, we need not move every host further (its index already binds the pronoun), but we must QR just one bottle. So the LF is (3) (omitting the internal structure of the definite DP, which stays as in (2)).
We sketch the most important steps in the semantic interpretation of this structure:

\[
\lbrack (3) \rbrack = 1
\]

iff

for every host \( x \), \( \lbrack \lbrack \text{vp} [\text{vp} \ldots] \text{ and } [\text{vp} \ldots] \rbrack \rbrack^{[1 \rightarrow x]} = 1 \)

iff

for every host \( x \),

\( \lbrack \text{just one bottle } 2[\text{tl bought } t_2] \rbrack^{[1 \rightarrow x]} = 1 \)

and \( \lbrack \text{tl served } \ldots \text{ the bottle } \text{he}_1 \text{ bought} \rbrack^{[1 \rightarrow x]} = 1 \)

iff

for every host \( x \),

there is just one bottle \( y \) such that \( \lbrack \text{tl bought } t_2 \rbrack^{[1 \rightarrow x, 2 \rightarrow y]} = 1 \)

and \( x \text{ served } \ldots \lbrack \text{the bottle } \text{he}_1 \text{ bought} \rbrack^{[1 \rightarrow x]} \)

iff

for every host \( x \),

there is just one bottle \( y \) such that \( x \text{ bought } y \)

and \( x \text{ served } \ldots \text{ the bottle } x \text{ bought} \).

11.5 Cooper’s analysis of E-Type pronouns

We have considered the analysis of the definite description paraphrase and have seen an LF that correctly expresses its truth-conditions. Let us now return to our original example with the pronoun, (1).

(1) Every host bought just one bottle of wine and served it with the dessert.

(1), we have noted, has the same truth-conditions we have just calculated. Can we devise an analysis for the E-Type pronoun which will yield this prediction?

One thing is clear. If the pronoun \( \text{it} \) in (1) is to have the same interpretation as the DP \( \text{the bottle } \text{he}_1 \text{ bought} \), then the LF representation of it has to contain an occurrence of the variable \( 1 \). Otherwise, the semantic value of the E-Type pronoun could not possibly vary with the value of that variable. It would not be able to denote different bottles of wine for different choices of host. It is also clear that the variable \( 1 \) cannot be the only meaningful constituent in the LF of it. The semantic values of it as a whole are not hosts, but bottles.

Here is an analysis due to Robin Cooper. Cooper proposes, in effect, that the LF representation of an E-Type pronoun consists of a definite article and a predicate that is made up of two variables. The first variable is of type \(<e,\langle e, t \rangle>\) and remains free in the sentence as a whole. The second variable is of type \( e \), and typically gets bound in the sentence (in our example, that is the variable \( 1 \)). We may think of these variables as unpronounced pro-forms. We will use the
notation $R_i$ for the type $<e,<e,t>$ variable and $pro_i$ for the type $e$ variable (where $i$ in each case is an index). The syntactic representation of an E-Type pronoun then looks like this:

\[(2)\]
\[
\begin{array}{c}
\text{DP} \\
\text{the} \\
\text{NP} \\
\text{N} \\
\text{DP} \\
\text{R}_7 \\
\text{pro}_1
\end{array}
\]

We have little to say about the syntactic or morphological side of the analysis. Let us simply assume that DPs which consist of a definite article followed by nothing but unpronounced items will always be spelled out as pronouns. For instance, (2) may become it.\(^9\)

How are such structures interpreted? The basic idea is that the $R$ variable receives a denotation from the context of utterance. In this respect, it is like a referential pronoun — except that it has a different semantic type, and therefore it requires that the context specify not a salient individual, but a salient 2-place relation. In the example at hand (sentence (1) under the intended reading), the $R$ variable denotes the relation which holds between people and bottles they have bought. (More accurately, it denotes the function $[\lambda x \in D . \lambda y \in D . y \text{ is a bottle that } x \text{ bought}].\(^{10}\)$ This relation (function) will have been made salient to the hearer as a result of his or her processing the first conjunct of the VP in (1), the predicate bought just one bottle. It will therefore be a natural candidate for the value of the free $R$ variable in the second conjunct. Given this choice of value for the node $R_7$, the rest of the interpretation of structure (2) will be straightforward. By Functional Application and the entry for the, the denotation of (2) under an assignment $g$ turns out to be the (unique) bottle bought by $g(1)$. We will spell all this out more precisely right below.

To summarize Cooper’s proposal, an E-Type pronoun is a definite description with an unpronounced predicate. The predicate is complex, and crucially contains a bound variable as well as a free variable. The free variable is essentially a relational noun (like “father”), except that its semantic value is not fixed lexically, but depends entirely on the utterance context. (It also is not restricted to semantic values which happen to be expressible by lexical items or constituents of the language. For instance, in our example the $R$ variable is interpreted as “bottle-bought-by”, but this is presumably not a well-formed constituent in English.) The bound variable serves as the free variable’s argument, and it will typically be co-indexed with other bound variables in the larger structure. As a
result, the E-Type pronoun's value co-varies systematically with the values of other variables in the sentence. (In our example, the bottle assigned to it varies with the person assigned to the QR trace of every host.) The proposal explains why an E-Type pronoun can be paraphrased by an explicit definite description. The explicit description and the E-Type pronoun are interpreted alike, except that the E-Type pronoun relies on context where the explicit description relies on lexical meaning.

To implement this analysis in our semantic theory, we don't need to make any real changes. We just have to make sure that our formulation of certain rules and principles covers higher-type variables as well as those of type e. First, let's assume that an index is not merely a number, but a pair of a number and a semantic type. So the tree we gave in (2) is really an abbreviation of (3) below.

(3)

```
(3) DP
       the NP
           N DP
               R<7,<e<e,t,t>> pro<1,e>
```

(Of course, we will get away with the abbreviated versions for all practical purposes, since the types in interpretable structures are always predictable from the environment.) A variable assignment is then defined as follows:

(4) A partial function g from indices to denotations (of any type) is a (variable) assignment iff it fulfills the following condition:

For any number n and type τ such that <n,τ> ∈ dom(g), g(n,τ) ∈ Dτ.

Now our so-called Pronouns and Traces Rule should really be called "Pro-forms and Traces Rule", and be formulated as follows:

(5) Pro-forms and Traces

If α is pro-form or trace, i is an index, and g is an assignment whose domain includes i, then [[αi]]E = g(i).

The Appropriateness Condition for utterances can stay as it is (cf. section 9.1.2):

(6) Appropriateness Condition

A context c is appropriate for an LF φ only if c determines a variable assignment gc whose domain includes every index which has a free occurrence in φ.
This formulation automatically covers the indices of our free R variables and ensures that they receive values of the appropriate type from the context. This is all we need for a precise analysis of our example. The LF we propose for (1) is (7) below.

\[
\text{(7) } \begin{array}{c}
\text{(7) } \\
\text{S} \\
\text{DP } 1 \text{ S} \\
\text{every host VP and VP} \\
\text{DP } 2 \text{ VP} \\
\text{just one bottle t_1 bought t_2} \\
\text{the NP} \\
\text{N DP} \\
\text{R_7 pro_1}
\end{array}
\]

The utterance context c in which (7) is uttered specifies the following assignment gc:

\[
\text{(8) } gc := [7 \rightarrow \lambda x \in D . \lambda y \in D . y \text{ is a bottle that } x \text{ bought}]
\]

By appropriate calculations (for the most part parallel to those we sketched in section 11.4 above), we determine that the utterance of (7) in the context c is true iff every host bought just one bottle and served the bottle he bought with dessert.

11.6 Some applications

Given the existence of E-Type pronouns, our grammar now predicts a great deal of structural ambiguity for every surface sentence that contains one or more pronouns. Not only are there often multiple choices of well-formed and interpretable indexings, but each pronoun can in principle be either a simple variable or a complex structure. We have not placed any syntactic constraints on the
generation of E-Type pronouns, and thus there should be E-Type readings for pronouns all over the place. Is this a problem?

A little reflection shows that an E-Type interpretation of a given pronoun will generally be indistinguishable from a referential interpretation, unless the pronoun is in the scope of a quantifier. Consider an example which illustrates this point. Before we were persuaded that we needed E-Type pronouns, we looked at examples like (1).

(1) Only one congressman admires Kennedy. He is very junior.

The he in (1), we argued, could be represented as a simple free variable, to which the utterance context assigned the congressman who admires Kennedy. But nothing stops us now from generating another LF in which he is an E-Type pronoun with the structure the [R, pro]. We even have good reason to believe that such an LF would be easy to interpret, since a context in which (1) as a whole is uttered will naturally furnish a salient referent for pro (namely, Kennedy) and a salient referent for R (namely, the relation between people and the congressmen that admire them). Given these choices of values, the structure the [R, pro] will denote the Congressman who admires Kennedy. So the overall interpretation is the same as on our original analysis. In general, when all the variables contained in an E-Type pronoun happen to be free variables, the pronoun as a whole receives a fixed referent (if it is interpretable at all). In that case, we can always give an alternative analysis on which this referent is directly assigned to a simple free variable pronoun. It is not surprising, therefore, that we needed to turn to examples containing a higher quantifier in order to get evidence for the existence of E-Type pronouns in the first place.

We assume that the coexistence of multiple analyses for examples like (1) is harmless. Presumably, when listeners process utterances, they will first consider simpler parses and entertain more complex ones only when forced to. So if there is a highly salient suitable referent for a simple individual variable (as in paradigm examples of deictic uses or of co-reference anaphora), the listener will have no motivation to try out the E-Type analysis. With an example like (1), it is less evident that the single variable analysis is simpler than the E-Type analysis from a processing point of view. We leave this as an empirical question. Our main point here is that the free generation of E-Type pronouns in all sorts of environments does not appear to give rise to unattested readings.

There is even some direct empirical evidence that pronouns which are not in the scope of quantifiers can have E-Type analyses. Consider our earlier example with an elliptical continuation.

(2) John bought just one bottle of wine and served it with the dessert. Bill did too.
(2) allows (in fact, strongly favors) a sort of "sloppy" reading, on which the second sentence means that Bill bought just one bottle and served the bottle that he, Bill, bought with the dessert. Given our assumptions about ellipsis in chapter 9, this reading calls for an E-Type analysis of it. We leave it to the reader to spell out the argument, as well as to explore more systematically what the existence of E-Type pronouns implies for the emergence of strict and sloppy readings under ellipsis.\textsuperscript{12}

Cooper's analysis of E-Type pronouns has been successfully applied to a variety of examples of great notoriety. We will conclude the chapter with a few brief illustrations.

The so-called donkey sentences were brought to the attention of modern linguists through the discussion of Geach.\textsuperscript{13} The typical donkey sentence has an indefinite DP in an "if" clause or relative clause, and a pronoun anaphoric to it in the main clause. Geach's original examples are (3) and (4).

(3) If a man owns a donkey, he beats it.

(4) Every man who owns a donkey beats it.

(5) Nobody who owns just one computer keeps it at home.

Much of our discussion from the previous sections obviously carries over to these examples. It is easy to see, for instance, that the it in (5) cannot be a referential pronoun (there is no contextually given computer it refers to), and that it can also not be a bound-variable pronoun (the sentence does not have the truth-conditions that would result from giving just one computer wide enough scope to bind the pronoun). A definite description paraphrase ("it" = "the computer that he or she owns") captures the intended reading, and it makes transparent that the pronoun's value must vary with the value of the QR trace of the "no" DP. Cooper's analysis exploits the plausible assumption that the phrase "nobody who owns just one computer" is suited to make salient the relation between people and their computers. We leave it to the reader to spell out the LF and interpretation of (5) in more detail.

Another staple of the literature on pronouns are the "paycheck" sentences. Here is Jacobson's edition of Karttunen's classic example.\textsuperscript{14}

(6) A woman who puts her paycheck in a federally insured bank is wiser than one who puts it in the Brown Employees' Credit Union.

The it cannot be referential, since it denotes no particular, contextually salient paycheck. It can also not be a bound variable, since there is no possible binder high enough in the structure. QR'ing the "antecedent" her paycheck all the way
out of its relative clause to a position c-commanding it is definitely not an option, since that would leave the her unbound. Cooper's E-Type analysis applies naturally. Let it be represented as the \([R_i \text{ pro}_j]\), where \(j\) is co-indexed with the relative pronoun (in the second relative clause), and \(i\) is mapped to the relation between people and their paychecks, which has plausibly been brought to salience with the earlier mention of paychecks.\(^\text{15}\)

An interesting extension of Cooper's analysis of paycheck sentences was developed in Jacobson's study of the so-called Bach-Peters paradox.\(^\text{16}\) An example is (7).

(7) Every boy who deserved it got the prize he wanted.

Bach and Peters had observed that there was no possible LF for this sentence in which both pronouns could be bound. Either every boy who deserved it c-commands he, or the prize he wanted c-commands it, but not both. Even if we disregard all syntactic constraints on movement and indexing, it is strictly impossible to derive from (7) an interpretable LF with no free variables. Jacobson argues that in every example of this kind, one of the pronouns is not a bound-variable pronoun, but an E-Type pronoun. In (7), for example, the he is bound by the every DP, but the it can be represented as the \([R_i \text{ pro}_j]\), where the pro is co-indexed with the trace of who, and the R variable refers to the relation "prize-wanted-by" (the relation between boys and the prizes they want). As Jacobson shows, this analysis makes sense of the distribution of possible determiners in these "paradoxical" sentences, as well as of syntactic constraints on the relations between pronouns and antecedents.

Cooper's E-Type analysis also throws light on many examples with pronouns whose antecedents are in the scope of modal verbs and adverbs and of propositional attitude verbs. These including the "fish" sentences discussed since the late 1960s\(^\text{17}\) and the cases of so-called modal subordination.\(^\text{18}\)

(8) John wants to catch a fish and eat it for dinner.

(9) A wolf might come in. It would eat you first.

Like the pronouns we have discussed in this chapter, the occurrences of "it" in (8) and (9) are not in suitable configurations to be bound by their "antecedents"; nor would such an analysis (even if permitted by syntax) represent their intuitive truth-conditions. And they also cannot be simply referential, for which fish or which wolf would be the referent? There may not even be any fish or wolves at all, yet (8) and (9) can be true. The solution, in a nutshell, exploits the fact that want, might, and would are quantifiers (namely, quantifiers over possible worlds). Once this is made explicit, the examples reveal a structure similar to
the ones we have analyzed here. For example, (8) means that every possible world w in which John gets what he wants is such that he catches a fish in w and eats the fish he catches in w for dinner in w. This is just a vague hint at an analysis, but we cannot do much more in the absence of a real analysis of sentence-embedding verbs and modals. Largely, that is beyond the scope of this book, though you will get a first taste in the next chapter.

Notes

1 For recent discussion of the applicability of the Coordinate Structure Constraint to QR, see E. Ruys, The Scope of Indefinites (Ph.D. dissertation, University of Utrecht, 1992; OTS Dissertation Series).

2 Evans’s original example contains “few” and “they” instead of “only one” and “he”. We have changed the example (and others below) in order to avoid complications related to plurality, a topic beyond the scope of this book. See G. Evans, “Pronouns,” Linguistic Inquiry, 11/2 (1980), pp. 337–62. For a fuller elaboration of the most important points in that paper, see idem, “Quantifiers and Relative Clauses I,” Canadian Journal of Philosophy, 7 (1977), pp. 467–536, repr. in G. Evans, Collected Papers (Oxford, Oxford University Press, 1985), pp. 76–152.

3 This holds regardless of whether we favor the “classical” analysis or the “presuppositional” analysis, at least if we stick to the official version of the latter, as it was presented in section 6.8. On that version, “Every woman was invited” presupposes that there is at least one woman, but not that there is more than one. Perhaps we should adopt a stronger version, on which every $\alpha \beta$ always presupposes that $[\alpha]$ has at least two elements. This looks like a reasonable proposal, but we cannot take the space here to examine it seriously.


6 The term was coined by Evans, but its usage in the linguistic literature has broadened, and it is often now applied to pronouns which Evans explicitly distinguished from “E-Type” pronouns. We, too, are applying it indiscriminately to all pronouns which are not covered by the theory we had up to now. See Evans, “Quantifiers”, and Neale, Descriptions, for discussion.

7 A flexible types approach would also be feasible at this point.

8 See Cooper, “Interpretation of Pronouns.” For recent versions and some discussion of alternatives, see also I. Heim, “E-Type Pronouns and Donkey Anaphora,”
One detail that we gloss over here is how the pronoun receives the correct gender, number, and person features. Our presuppositional treatment of features from section 9.1.2 can be applied to E-Type pronouns as well. Suppose that the structure in (2) has features adjoined to it, which determine the shape of the spelled-out pronoun. Then the presuppositions generated by these features will, in effect, put constraints on the denotation of R7. You are invited to elaborate this remark and apply it to an example.

This is the Schönfinkelization of the 2-place relation just described.

See exercise in section 8.2.


Evans calls the pronoun in the paycheck sentence a "pronoun of laziness", which he specifically distinguishes from an "E-Type pronoun". Following Geach's terminology, a pronoun of laziness is equivalent to an (exact) repetition of its antecedent. The paycheck pronoun fits this description, since (6) is indeed accurately paraphrased by substituting "her paycheck" for "it". Notice that our previous examples of E-Type pronouns do not have this property. Example (1) of section 11.5, for example, does not mean the same thing as "Every host bought just one bottle of wine and served just one bottle of wine with the dessert". As we have noted before, we do not adhere to Evans's terminology, but use "E-Type pronoun" in a broader sense. If the analysis we have adopted from Cooper is right, then the indiscriminate terminology is indeed theoretically justified. Be that as it may, it is common. The term "pronoun of laziness" is rarely used these days, but has also been employed in a very broad sense (e.g., in Partee, "Opacity").


12 First Steps Towards an Intensional Semantics

In chapter 2, we made a move that cost us some effort to sell to our readers. We identified the denotations of sentences with their actual truth-values. We put potential protests to rest by demonstrating that the resulting extensional theory as a whole was still able to pair sentences with their truth-conditions. While sticking to an extensional semantics, we managed to develop a theory of meaning.\(^1\) In this chapter, we will have to face the limitations of the extensional framework we have been working in up to now. We will see that it breaks down in certain contexts, a fact that Frege was well aware of.\(^2\) Fortunately, a minor change in the semantic system will repair the problem (at least to a certain extent), and most of what we have learned will stay intact.

12.1 Where the extensional semantics breaks down

One of the basic assumptions underlying our semantics has been that the extension of a complex expression can be computed from the extensions of its parts in a stepwise fashion, using no more than a handful of composition principles. While this assumption was correct for the examples we have discussed so far, it is plain wrong for examples of the following kind.

(1) Mary believes Jan is loyal.

(2) Mary believes Dick is deceitful.
Here is the problem. Our type-driven interpretation system implies that in the two trees above, the denotations of the higher VP-nodes are computed from the denotations of the matrix V-nodes and the denotations of the embedded S-nodes by the same mode of composition. The two V-nodes dominate the same lexical item, hence are assigned the same denotation. The two S-nodes both denote a truth-value. Suppose now that in the actual world, Jan is indeed loyal, and Dick deceitful; that is, \([\text{Jan is loyal}] = [\text{Dick is deceitful}] = 1\). Consequently, the two embedded S-nodes have the same denotation, and we predict that 

\[
[v \text{ believes } [s \text{ Jan is loyal}] = v \text{ believes } [s \text{ Dick is deceitful}]]
\]

and 

\[
[Mary [v \text{ believes } [s \text{ Jan is loyal}]] = Mary [v \text{ believes } [s \text{ Dick is deceitful}]]]
\]

This isn't a welcome prediction, though, since Mary may believe that Jan is loyal without believing that Dick is deceitful (and vice versa).

We want to say that sentences embedded under believe are nonextensional contexts. Nonextensional contexts are also called “oblique”, “opaque”, or “indirect”. Other words that create nonextensional contexts include the verbs hope, fear, look (as in look smart), seem, seek, the adjectives alleged or fake, the preposition about, the connective because, and modal words of various categories like must, may, probably, obviously, provable, and permissible.

**Exercise**

Show that alleged, seem, might, about, and because create nonextensional contexts. Demonstrate this by constructing compositionality arguments of the
sort we have just given for believe. Note that nonextensional contexts do not have to be sentential. The nonextensional context created by alleged is the modified noun phrase: for example, 'son of a sphinx in alleged son of a sphinx.

12.2 What to do: intensions

Frege proposed that in opaque contexts, expressions denote their Sinn (sense). But what is a Fregean Sinn? Recall the quote from Dummett in chapter 2 that addresses a frequently voiced complaint about Frege being rather vague as to what the senses of expressions are:

"It has become a standard complaint that Frege talks a great deal about the senses of expressions, but nowhere gives an account of what constitutes such a sense. This complaint is partly unfair: for Frege the sense of an expression is the manner in which we determine its reference, and he tells us a great deal about the kind of reference possessed by expressions of different types, thereby specifying the form that the senses of such expressions must take. . . . The sense of an expression is the mode of presentation of the referent: in saying what the referent is, we have to choose a particular way of saying this, a particular means of determining something as a referent."³

The Fregean sense of an expression, then, is the mode of presentation of its extension (reference, Bedeutung). It's a particular means of determining the extension. But what kind of formal object is a "means of determining an extension"? It could just be a linguistic expression – an expression of set theory, for example. Or it could be an algorithm computing the values of a function for arbitrary arguments. Different expressions might specify the same set, and different algorithms might compute the same function. Expressions of set theory and algorithms, then, are means of determining extensions. But there are other, more abstract possibilities. One was proposed by Rudolf Carnap, a student of Frege's.⁴ Here is the idea.

The truth of a sentence depends on the circumstances. It's now true that you are in Amsterdam, but in a little while, that's not going to be the case any more. And if circumstances had been different, you might never have left your native Buffalo at all.

The extension of a predicate depends on the circumstances as well. You are a member of the garden club, but you haven't always been, and you might never have joined. Quite generally, then, the extension of an expression depends on
possible circumstances. An \textit{intension} in Carnap's sense is something that determines \textit{how} extensions depend on possible circumstances. David Lewis tells us where to go from there:

What sort of things determine how something depends on something else? \textit{Functions}, of course; functions in the most general set-theoretic sense, in which the domain of arguments and the range of values may consist of entities of any sort whatsoever, and in which it is not required that the function be specifiable by any simple rule. We have now found something to do at least part of what a meaning for a sentence, name, or common noun does: a function which yields as output an appropriate extension when given as input a package of the various factors on which the extension may depend. We will call such an input package of relevant factors an \textit{index}; and we will call any function from indices to appropriate extensions for a sentence, name, or common noun an \textit{intension}.\textsuperscript{5}

A (Carnapian) intension,\textsuperscript{6} then, is a function from indices to appropriate extensions. To simplify matters in this introductory text, let us neglect all index dependence except for dependence on possible worlds. That is, we will neglect temporal dependence, speaker dependence, and what have you. For our limited purposes here, then, (Carnapian) intensions are functions from possible worlds to appropriate extensions. The intension of a sentence is a function from possible worlds to truth-values. The intension of a 1-place predicate is a function that maps possible worlds into characteristic functions of sets of individuals, etcetera. If you wonder about possible worlds, here is what David Lewis says about them:

The world we live in is a very inclusive thing. Every stick and every stone you have ever seen is part of it. And so are you and I. And so are the planet Earth, the solar system, the entire Milky Way, the remote galaxies we see through telescopes, and (if there are such things) all the bits of empty space between the stars and galaxies. There is nothing so far away from us as not to be part of our world. Anything at any distance at all is to be included. Likewise the world is inclusive in time. No long-gone ancient Romans, no long-gone pterodactyls, no long-gone primordial clouds of plasma are too far in the past, nor are the dead dark stars too far in the future, to be part of the same world. . . .

The way things are, at its most inclusive, means the way this entire world is. But things might have been different, in ever so many ways. This book of mine might have been finished on schedule. Or, had I not been such a commonsensical chap, I might be defending not only a plurality of possible worlds, but also a plurality of impossible worlds, whereof you speak truly by contradicting yourself. Or I might not have existed at all
– neither myself, nor any counterpart of me. Or there might never have been any people. Or the physical constants might have had somewhat different values, incompatible with the emergence of life. Or there might have been altogether different laws of nature; and instead of electrons and quarks, there might have been alien particles, without charge or mass or spin but with alien physical properties that nothing in this world shares. There are ever so many ways that a world might be: and one of these many ways is the way that this world is.7

12.3 An intensional semantics

In this section, we will develop a semantic system that allows denotations to be (Carnapian) intensions, and we will show that such a system solves the problem we ran into at the beginning of this chapter.

There are various intensional frameworks in the possible worlds tradition that you find in the literature.8 The framework we chose here for illustration is a conservative extension of the extensional semantics we have been working with all along. We start out with a recursive definition of an intensional system of semantic types (Montague's), which will be followed by a parallel definition of a typed system of semantic domains.

(1) Recursive definition of semantic types
(a) e is a type.
(b) t is a type.
(c) If a and b are types, then <a,b> is a type.
(d) If a is a type, then <s,a> is a type.
(e) Nothing else is a type.

(2) Semantic domains
Let W be the set of all possible worlds. Associated with each possible world w is the domain of all individuals existing in w. Let D be the union of the domains of all possible worlds. That is, D contains all individuals existing in the actual world, but also all individuals existing in any of the merely possible worlds. It is the set of all possible individuals. The set of intensional domains is now defined as follows:
(a) \( D_e = D \)
(b) \( D_t = \{0, 1\} \)
(c) If a and b are semantic types, then \( D_{<a,b>} \) is the set of all functions from \( D_a \) to \( D_b \).
(d) If a is a type, then \( D_{<s,a>} \) is the set of all functions from \( W \) to \( D_a \).
In addition to our familiar extensions, we now have intensions. The domain \( D_{s,t} \) contains all functions from \( W \) to \( \{0, 1\} \), for example: that is, all characteristic functions of subsets of \( W \). Possible world semanticists take such functions to be the formal construals of *propositions*. Construing propositions as characteristic functions of sets of possible worlds is natural, as emphasized in the following statement by Robert Stalnaker:

> The explication of *proposition* given in formal semantics is based on a very homely intuition: when a statement is made, two things go into determining whether it is true or false. First, what did the statement say: what proposition was asserted? Second, what is the world like: does what was said correspond to it? What, we may ask, must a proposition be in order that this simple account be correct? It must be a rule, or a function, taking us from the way the world is into a truth value. But since our ideas about how the world is change, and since we may wish to consider the statement relative to hypothetical and imaginary situations, we want a function taking not just the actual state of the world, but various possible states of the world into truth values. Since there are two truth values, a proposition will be a way – any way – of dividing a set of possible states of the world into two parts: the ones that are ruled out by the truth of the proposition, and the ones that are not.  

Let us look at some examples of lexical entries. Following Montague, we will relativize the interpretation function to a possible world and an assignment function. As before, we can drop reference to an assignment when the choice of assignment doesn’t matter, since we define for any possible world \( w \), and any expression \( \alpha \):

\[
(3) \quad [[\alpha]]^w := [[\alpha]]^{w, \emptyset}.
\]

We now have:

\[
(4) \quad \text{Names} \\
\text{For any possible world } w: \\
[\text{Jan}]^w = \text{Jan.} \\
[\text{Ann}]^w = \text{Ann.} \\
\text{etc.}
\]

Following Saul Kripke, we treat proper names as *rigid designators*.\(^{10}\) Their reference is picked out in the actual world, and they denote the same individual that was so picked out in every possible world.\(^{11}\) Proper names, then, differ from definite descriptions like *the coldest winter*, which may denote different winters in different possible worlds.
The extensions of predicates may vary depending on the circumstances as well:

(5) Easy predicates
For any possible world w:

\[ \text{[smoke]}^w = \lambda x \in D \cdot x \text{ smokes in } w \]
\[ \text{[love]}^w = \lambda x \in D \cdot [\lambda y \in D \cdot y \text{ loves } x \text{ in } w] \]
\[ \text{[cat]}^w = \lambda x \in D \cdot x \text{ is a cat in } w \]

etc.

Nothing exciting happens with quantifiers. Their extension does not depend on the circumstances, but for reasons of generality, we carry the world parameter along, as we did with proper names:

(6) Determiners
For any possible world w:

\[ \text{[every]}^w = \lambda f \in D_{\text{et},t>} \cdot [\lambda g \in D_{\text{et},t>} \cdot \text{ for all } x \text{ such that } f(x) = 1, g(x) = 1] \]

etc.

The fragment we have been building does not yet require any new composition rules. The ones we already have in place will do, except that the interpretation function depends now not just on an assignment, but also on a possible world. Both parameters have to be schlepped along as the interpretation machinery works its way through a given tree. Assignments, too, are what they used to be: partial functions from the set of natural numbers to D (but D has changed, of course).

Consider now the attitude verbs believe, know, hope, and so on. As a starting point, we will pursue an approach (in the spirit of Hintikka\(^{12}\)) that has it that the content of an attitude can be characterized by a set of possible worlds: namely, those that are compatible with the attitude. Here is how David Lewis illustrates this rather simple idea:

The content of someone's knowledge of the world is given by his class of epistemically accessible worlds. These are the worlds that might, for all he knows, be his world; world W is one of them iff he knows nothing, either explicitly or implicitly, to rule out the hypothesis that W is the world where he lives. Likewise the content of someone's system of belief about the world (encompassing both belief that qualifies as knowledge and belief that fails to qualify) is given by his class of doxastically accessible worlds. World W is one of those iff he believes nothing, either explicitly or implicitly, to rule out the hypothesis that W is the world he lives.

Whatever is true at some epistemically or doxastically accessible world is epistemically or doxastically possible for him. It might be true, for all
he knows or for all he believes. He does not know or believe it to be false. Whatever is true throughout the epistemically or doxastically accessible worlds is epistemically or doxastically necessary; which is to say that he knows or believes it, perhaps explicitly or perhaps only implicitly.

Since only truths can be known, the knower's own world always must be among his epistemically accessible worlds. Not so for doxastic accessibility. If he is mistaken about anything, that is enough to prevent his own world from conforming perfectly to his system of belief.

Lexical entries for attitude verbs will accordingly look as follows:

(7) **Attitude verbs**

For any possible world \( w \):

\[
\langle \text{believe} \rangle^w = \lambda p \in D_{\text{es,do}}. [\lambda x \in D . p(w') = 1, \text{for all } w' \in W \text{ that are compatible with what } x \text{ believes in } w]
\]

\[
\langle \text{know} \rangle^w = \lambda p \in D_{\text{es,do}}. [\lambda x \in D . p(w') = 1, \text{for all } w' \in W \text{ that are compatible with what } x \text{ knows in } w]
\]

\[
\langle \text{hope} \rangle^w = \lambda p \in D_{\text{es,do}}. [\lambda x \in D . p(w') = 1, \text{for all } w' \in W \text{ that are compatible with what } x \text{ hopes in } w]
\]

What a person believes, knows, or hopes can vary from one possible world to another. The world parameter in the lexical entries for attitude verbs matters, then, as it does for most other predicates. For example, we can't simply write "\( w \) is a world where Mary's hopes come true" unless we mean: "\( w \) is a world where those hopes come true that Mary has in the actual world". Otherwise, we must make explicit which world we are talking about; that is, we must write things of the form: "\( w' \) is a world where Mary's hopes in \( w \) come true".

For Mary to believe that \( p \) it is **not** sufficient that \( p \) be true in some world that is compatible with what she believes. If her belief worlds include worlds where \( p \) is true as well as worlds where \( p \) is false, then she is agnostic as to whether or not \( p \). Likewise for the other attitude verbs.

We can say that a possible world having such and such properties is compatible or incompatible with what Mary believes, or that \( w \) is a world where Mary's hopes come true. But never say things like "Mary hopes to be in \( w \)", "Mary believes that she is in \( w \)", "Mary knows that she is in \( w \)". Taken literally, these make sense only if Mary is so choosy that there is only one possible world that satisfies her desires, or so opinionated that just one possible world conforms to her beliefs, or omniscient. The same objection applies to locutions like "the world according to Mary", "the world of her dreams", "the world she thinks she lives in", etcetera. We talk this way when we are not semanticists, but it doesn't make literal sense within the present theoretical framework. If we don't want to commit ourselves to such unrealistic assumptions, we have
to express ourselves in a way that makes it clear that actual desires, beliefs, etcetera, pick out sets of worlds. For instance, we might write: "w is a world where Mary's hopes come true", "w is a world that conforms to everything Mary believes", "what Mary knows is true in w".

Attitude verbs denote functions that apply to propositions. But so far, our interpretation procedure doesn't yet deliver any propositions, nor intensions of any other kind. Look at what happens when we try to interpret sentence (8):

(8) Mary believes Jan is loyal.

\[
\begin{align*}
(8') & \quad S \\
& \quad DP \quad VP \\
& \quad \text{Mary} \quad V \quad S \\
& \quad \text{believes} \quad DP \quad VP \\
& \quad \text{Jan} \quad V \quad AP \\
& \quad \text{is} \quad A \\
& \quad \text{loyal}
\end{align*}
\]

We have for any possible world w:

(a) \([\text{Mary} [\text{believes} [\text{Jan is loyal}]]]^w = (\text{FA})\)
(b) \([\text{believes} [\text{Jan is loyal}]]^w ([\text{Mary}]^w) = (\text{lexical entry Mary})\)
(c) \([\text{believes} [\text{Jan is loyal}]]^w (\text{Mary})\)

We are stuck. The denotation of the verb believes applies to a proposition: that is, an intension of type <s,t>. The interpretation system, however, provides only \([\text{Jan is loyal}]^w\), which is a denotation of type t, hence a mere truth-value: 1 if Jan is loyal in w, and 0 otherwise. Here is the proof:

\begin{align*}
[\text{Jan is loyal}]^w & = (\text{FA}) \\
[\text{is loyal}]^w ([\text{Jan}]^w) & = (\text{lexical entry Jan and emptiness of is}) \\
[\text{loyal}]^w (\text{Jan}) & = (\text{lexical entry loyal}) \\
[\lambda x \in D . x \text{ is loyal in } w] (\text{Jan}).
\end{align*}
By the definition of the $\lambda$-notation, $[\lambda x \in D \cdot x \text{ is loyal in } w] (\text{Jan}) = 1$ iff Jan is loyal in $w$.

To make sure that the lexical requirements of the verb believes can be met, we introduce an additional composition principle: Intensional Functional Application (IFA):\(^{14}\)

(9) **Intensional Functional Application (IFA)**

If $\alpha$ is a branching node and $\{\beta, \gamma\}$ the set of its daughters, then, for any possible world $w$ and any assignment $\alpha$, if $[\beta]^w_{\wedge}$ is a function whose domain contains $\lambda w'. [\gamma]^w_{\wedge}$, then $[\alpha]^w_{\wedge} = [\beta]^w_{\wedge}(\lambda w'. [\gamma]^w_{\wedge})$.

We can now continue our computation at the point where we got stuck:

(c) $[\text{believes} [\text{Jan is loyal}]^w (\text{Mary}) = \text{(by } \text{IFA})$

(d) $[\text{believes}]^w (\lambda w'. [\text{Jan is loyal}]^w) (\text{Mary}) = \text{(by previous computation)}$

(e) $[\text{believes}]^w (\lambda w'. \text{Jan is loyal in } w') (\text{Mary}) = \text{(by lexical entry believe)}$

(f) $[\lambda p \in D_{\leq Ap} \cdot [\lambda x \in D \cdot p(w') = 1, \text{for all } w' \in W \text{ that are compatible with what } x \text{ believes in } w]] (\lambda w'. \text{Jan is loyal in } w') (\text{Mary}) = \text{(by definition of } \lambda\text{-notation)}$

(g) $[\lambda x \in D \cdot [\lambda w'. \text{Jan is loyal in } w'] (w') = 1, \text{for all } w' \in W \text{ that are compatible with what } x \text{ believes in } w] (\text{Mary}) = \text{(by definition of } \lambda\text{-notation)}$

(h) $[\lambda x \in D \cdot \text{Jan is loyal in } w', \text{for all } w' \in W \text{ that are compatible with what } x \text{ believes in } w] (\text{Mary})$.

Finally (again by definition of the $\lambda$-notation) we have:

$[\lambda x \in D \cdot \text{Jan is loyal in } w', \text{for all } w' \in W \text{ that are compatible with what } x \text{ believes in } w] (\text{Mary}) = 1$ iff Jan is loyal in $w'$, for all $w' \in W$ that are compatible with what Mary believes in $w$.

We have arrived at the correct truth-conditions for sentence (8). (8) is true in the actual world iff Mary's actual beliefs exclude all possible worlds in which Jan is not loyal. In an analogous way, we obtain the right truth-conditions for (10):

(10) Mary believes Dick is deceitful.
(10) comes out true in the actual world iff there is no world that is compatible with Mary's actual beliefs in which Dick is not deceitful. Given an intensional semantics, we do not run into any difficulties any more when we assume that all of the sentences 11(a)–(d) might actually be true together:

(11) (a) Jan is loyal.
    (b) Dick is deceitful.
    (c) Mary believes that Jan is loyal.
    (d) Mary does not believe that Dick is deceitful.

The problem we started out with in this chapter is now gone. The solution is very much in the spirit of Frege. The usual denotations are extensions. But for nonextensional contexts, Intensional Functional Application allows a switch to intensions. The switch is triggered by particular lexical items — those that create nonextensional contexts. Whether a lexical item does or does not create a non-extensional context, then, is part of the information conveyed by its denotation, like any other information about selectional restrictions.

**Exercise 1**

Look at sentence (i):

(i) Mary hopes that a plumber is available.

(i) may mean that there is a particular plumber, say Mr French, whom Mary hopes will be available (de re interpretation of a plumber). Or else it may mean that she hopes that some plumber or other is available (de dicto interpretation of a plumber). It is customary to treat the ambiguity as a scope ambiguity. On the de re interpretation, a plumber is moved into the matrix clause. On the de dicto interpretation, a plumber can be interpreted in situ.

(a) Draw appropriate LFs for the two readings of (i).
(b) Add lexical entries to our fragment as needed. Treat that and is as semantically empty, but treat a as a quantifying determinant.
(c) Compute the truth-conditions for both LFs.
Exercise 2

Look at sentence (ii):

(ii) Lee is an alleged drug dealer from Springfield.

Draw an appropriate LF for (ii) and compute its truth-conditions. Add lexical entries to our fragment as needed. Treat is and an as semantically empty, and drug dealer as an unanalyzed common noun. The truth-conditions you want to end up with should imply that (ii) is true in the actual world just in case Lee is a drug dealer from Springfield in all those possible worlds that are compatible with the allegations in the actual world. Let's not fuss about the fact that allegations may vary from time to time and place to place, even in a single world. We chose to ignore that kind of dependence here to make things easier for all of us.

12.4 Limitations and prospects

Carnap\textsuperscript{15} insists that sentences embedded under attitude verbs are neither extensional nor intensional contexts. And he is right. Take two sentences that are true in the same possible worlds but do not have to be believed together. Here is an example due to John Bigelow:\textsuperscript{16}

(1) (a) Robin will win.
(b) Everyone who does not compete, or loses, will have done something Robin will not have done.

(2) (a) Marian believes that Robin will win.
(b) Marian believes that everyone who does not compete, or loses, will have done something Robin will not have done.

It may take some time to figure this out, but (1a) and (1b) are true in exactly the same possible worlds. They express the same proposition, then. But if what we believe are propositions (as our analysis of attitude verbs assumes), anybody who believes (1a) should also believe (1b). This is not right. (2a) and (2b) can have different truth-values. We are in trouble again. Propositions are still not good enough as objects of beliefs and other attitudes.
Carnap proposed the concept of “intensional isomorphism” or “intensional structure” as a remedy.\textsuperscript{17} David Lewis follows up on this idea and identifies “meanings” with “semantically interpreted phrase markers minus their terminal nodes: finite ordered trees having at each node a category and an appropriate intension.”\textsuperscript{18} A slightly simpler construct, “structured propositions”, is proposed by Cresswell and von Stechow.\textsuperscript{19} There is no agreement on the issue yet. The good news is, however, that the uncertainty in the area of propositional attitudes does not seem to have a lot of repercussions on the way linguists do semantics every day. A slight change led us from an extensional system to an intensional one. The switch to a hyperintensional system should not be much more eventful. What we have learned about particular extensional or intensional phenomena should be adaptable to a new foundation without too much ado. Barbara Partee observes:

Many of the most fundamental foundational issues in formal semantics (and in semantics as a whole) remain open questions, and there may be even less work going on on them now than there was in the seventies; perhaps this is because there is more work by linguists and less by philosophers, so the empirical linguistic questions get most of the attention now.\textsuperscript{20}

The empirical linguistic questions have certainly been our main concern in this book.

Notes

1 At least of those ingredients of meaning that have to do with truth-conditions. There are other components of meaning that are not usually dealt with in semantics, but in pragmatics.
6 The term “Carnapian intension” is used by Lewis (ibid.).


Scholars differ with respect to their views on trans-world identity of individuals. David Lewis has argued, for example, that individuals should not be assumed to exist in more than one possible world. Individuals may be related to very similar individuals in other possible worlds via a counterpart relationship. See D. K. Lewis, "Counterpart Theory and Quantified Modal Logic," *Journal of Philosophy*, 65 (1968), pp. 113–26; idem, *Plurality of Worlds*.

Hintikka, *Knowledge and Belief*.


Carnap, *Meaning and Necessity*.


Carnap, *Meaning and Necessity*.

Lewis, *Plurality of Worlds*, p. 27.

Lewis, "General Semantics," p. 182.


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