## Morita-type equivalences for operator algebras

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This talk will describe work of George Eleftherakis (Athens).

### 1 TRO equivalence

1.1 A TRO  $\mathcal{M}$  is a linear subspace of some  $B(H_1, H_2)$  such that

$$\mathcal{M}\mathcal{M}^*\mathcal{M} \subseteq \mathcal{M}$$
.

TRO's arise as:

- Corners of C\*-algebras,
- injective operator spaces,
- normalizing spaces of (possibly nonselfadjoint) algebras (w. Todorov).
- A TRO  $\mathcal{M}$  is a Hilbert (bi)-module over the C\*-algebras  $\mathcal{M}^*\mathcal{M}$  and  $\mathcal{M}\mathcal{M}^*$ . Conversely, Hilbert modules can be represented as TRO's.

w\*-closed TRO's are generated by their partial isometries; they need not contain any unitaries.

1.2 Normalisers: A normaliser of  $\mathcal{A} \subseteq B(H_1)$  into  $\mathcal{B} \subseteq B(H_2)$  is  $T \in B(H_1, H_2)$  such that  $T^*\mathcal{B}T \subseteq \mathcal{A}$  and  $T^*\mathcal{A}T \subseteq \mathcal{B}$ .

TRO equivalence: existence of 'sufficiently many' normalisers.

**Definition 1** Two (w\*-closed) algebras  $\mathcal{A}$ ,  $\mathcal{B}$  are called TRO equivalent if there exists a TRO  $\mathcal{M}$  such that  $\mathcal{A} = [\mathcal{M}^*\mathcal{B}\mathcal{M}]^{-w^*}$  and  $\mathcal{B} = [\mathcal{M}\mathcal{A}\mathcal{M}^*]^{-w^*}$ .

Examples: unitarily equivalence; Morita equivalence of  $W^*$ -algebras.

Example:  $\mathcal{A} \sim_{TRO} M_n(\mathcal{A})$  with  $\mathcal{M} = C_n(\mathcal{A})$ :

$$\binom{*}{*} \cdot (*) \cdot (**) = \binom{**}{**} \qquad \binom{*}{*} \cdot \binom{**}{**} \cdot (**) = (*)$$

- 1.3 TRO equivalence is an equivalence relation.
- 1.4 **Theorem 2** The reflexive algebras A, B are TRO equivalent if and only if there exists a \*-isomorphism

$$\theta: \Delta(\mathcal{A})' \to \Delta(\mathcal{B})'$$

such that

$$\theta(\text{Lat}(\mathcal{A})) = \text{Lat}(\mathcal{B}).$$

Reflexive algebras: Given a (unital, w\*-closed) algebra  $\mathcal{A} \subseteq B(H)$  and a (strongly closed) lattice of projections  $\mathcal{L}$ , let

$$\operatorname{Lat}(\mathcal{A}) = \{ P \in \operatorname{Proj}(H) : P^{\perp} \mathcal{A} P = \{0\} \}$$
$$\operatorname{Alg}(\mathcal{L}) = \{ A \in B(H) : P^{\perp} A P = 0 \ \forall P \in \mathcal{L} \}$$
$$\mathcal{A} \text{ reflexive } : \mathcal{A} = \operatorname{Alg} \mathcal{L}.$$

Examples: W\* algebras, nest algebras (here  $\mathcal{L}$  is totally ordered - a nest) CSL algebras (elements of  $\mathcal{L}$  commute).

# 2 TRO equivalence vs spacial Morita equivalence.

**Example 3** There exist nests  $\mathcal{N}_1, \mathcal{N}_2$  which are isomorphic but the algebras  $\mathcal{N}_1'', \mathcal{N}_2''$  are not isomorphic. Thus isomorphism of the lattices does not guarantee TRO equivalence, even for multiplicity free nest algebras. See item 3.1 below.

On  $H_1 = \ell^2(\mathbb{Q} \cap [0,1])$  define, for each  $t \in [0,1]$ ,

$$Q_t^+ = \{f : \operatorname{supp} f \subseteq [0, t]\}$$

$$Q_t^- = \{f : \operatorname{supp} f \subseteq [0, t)\}.$$

On  $L^2([0,1],\lambda)$  define, for each  $t \in [0,1]$ ,

$$N_t = \{f : \operatorname{supp} f \subseteq [0, t]\}.$$

Let

$$\mathcal{N}_1 = \{Q_t^{\pm} : t \in [0, 1]\} \text{ on } H_1 = \ell^2(\mathbb{Q} \cap [0, 1])$$
  
 $\mathcal{N}_2 = \{Q_t^{\pm} \oplus N_t : t \in [0, 1]\} \text{ on } H_2 \equiv H_1 \oplus L^2([0, 1], \lambda).$ 

2.2 **Proposition** Two CSL algebras A, B have isomorphic lattices iff they are "spacially Morita equivalent":

**Definition** Let  $\mathcal{A} \subset B(H_1)$ ,  $\mathcal{B} \subset B(H_2)$  be  $w^*$ -closed algebras. If there exist spaces  $\mathcal{U} \subset B(H_1, H_2)$ ,  $\mathcal{V} \subset B(H_2, H_1)$  such that  $\mathcal{BUA} \subset \mathcal{U}$ ,  $\mathcal{AVB} \subset \mathcal{V}$ ,  $[\mathcal{VU}]^{-w^*} = \mathcal{A}$ ,  $[\mathcal{UV}]^{-w^*} = \mathcal{B}$  then we say that the algebras  $\mathcal{A}$ ,  $\mathcal{B}$  are spacially Morita equivalent and the system

$$\begin{pmatrix} A & V \\ U & B \end{pmatrix}$$

is called a spacial Morita context.

- 2.3 In general, TRO equivalence  $\Rightarrow$  spacial Morita equivalence.
- 2.4 Definition: Given a commutative subspace lattice (CSL)  $\mathcal{L}$ , the class of all w\*-closed algebras  $\mathcal{A}$  containing a masa and s.t. Lat  $\mathcal{A} = \mathcal{L}$  has a minimal element  $\mathcal{A}_{min}(\mathcal{L})$  [Arveson]. In case  $\mathcal{A}_{min}(\mathcal{L}) = \text{Alg } \mathcal{L}$  we say  $\mathcal{L}$  is synthetic.

Using these methods, one shows:

**Proposition 4** Synthesis is preserved under lattice isomorphisms.

Qu: Under epimorphisms?

### 3 TRO equivalence and CSL algebras.

**Definition 5** Let  $\phi : \mathcal{L}_1 \to \mathcal{L}_2$  be an isomorphism between CSL's. Then  $\phi$  restricts to an isomorphism between the 'atomic parts', but not necessarily between the 'continuous parts' (see example above). If it does, we say  $\phi$  respects continuity.

**Proposition 6** Two CSL algebras are TRO equivalent iff there exists an isomorphism between their lattices which respects continuity.

#### Consequences of TRO equivalence for CSL algebras:

Let  $\mathcal{A}, \mathcal{B}$  be CSL algebras which are TRO equivalent via  $\mathcal{M}$ . There is a bijective correspondence between the  $w^*$ -closed ideals of the algebras  $\mathcal{A}$  and  $\mathcal{B}$ . It follows that

$$\mathcal{A}$$
 strongly reflexive  $\Leftrightarrow \mathcal{B}$  strongly reflexive  $R_1(\mathcal{A}) = 0 \Leftrightarrow R_1(\mathcal{B}) = 0$ 

$$\mathcal{A} \text{ semisimple } \Leftrightarrow \mathcal{B} \text{ semisimple}$$

$$\mathcal{A} = \operatorname{Rad}(\mathcal{A})^{-w^*} \Leftrightarrow \mathcal{B} = \operatorname{Rad}(\mathcal{B})^{-w^*}$$

# 4 Morita equivalence for W\*-algebras (Rieffel)

The category  $_{\mathcal{A}}\mathfrak{M}$  of Hilbert  $\mathcal{A}$ -modules for a W\* algebra  $\mathcal{A}$ : Objects:  $(H, \pi)$  where  $\pi$  is a normal (: w\*-continuous) \*-representation. Morphisms:

$$\operatorname{Hom}_{\mathcal{A}}(H_1, H_2) = \{ T \in B(H_1, H_2) : T\pi_1(a) = \pi_2(a)T \ \forall a \in \mathcal{A} \}.$$

A functor 
$$\mathcal{F}: \mathcal{A}\mathfrak{M} \to \mathcal{B}\mathfrak{M}$$
 is called a \*-functor if  $T \in \operatorname{Hom}_{\mathcal{A}}(H_1, H_2) \Rightarrow \mathcal{F}(T)^* = \mathcal{F}(T^*)$ 

Example: If  $\phi: \mathcal{B} \to \mathcal{A}$  is a \*-isomorphism, then  $\mathcal{F}(H, \alpha) = (H, \alpha \circ \phi)$  on objects and  $\mathcal{F}(T) = T$  on morphisms defines an equivalence \*-functor  $\mathcal{F}: {}_{\mathcal{A}}\mathfrak{M} \to {}_{\mathcal{B}}\mathfrak{M}$ .

**Theorem 7 (Rieffel)**  $_{\mathcal{A}}\mathfrak{M}$  and  $_{\mathcal{B}}\mathfrak{M}$  are equivalent via \*-functors (write  $A \sim_R B$ ) iff there is an 'abstract Morita context'  $\begin{pmatrix} A & \mathcal{X} \\ \mathcal{X}^* & \mathcal{B} \end{pmatrix}$  for  $W^*$ -algebras.

equivalently [Connes]

$$A \sim_R B \iff \exists \text{ faithful normal reprs. s.t. } \alpha(\mathcal{A})' \simeq \beta(\mathcal{B})'.$$

In view of Theorem 2 this remark has the following equivalent version:

**Theorem 8** The W\*-algebras  $\mathcal{A}$ ,  $\mathcal{B}$  are Morita equivalent if and only if they have faithful normal representations  $\alpha$ ,  $\beta$  on Hilbert spaces such that the algebras  $\alpha(\mathcal{A})$ ,  $\beta(\mathcal{B})$  are TRO equivalent.

$$A \sim_R B \iff \exists faithful \ normal * reps: \quad \alpha(A) \sim_{TRO} \beta(B)$$

### 5 The category $\mathfrak{M}$ for dual operator algebras.

- 5.1 Dual operator algebras: (unital) w\*-closed subalgebras  $(A, w^*)$  of some B(H). There is an abstract characterisation (Le Merdy): they are the operator algebras that have an *operator space* predual.
- 5.2 The categories  $_{\mathcal{A}}\mathfrak{M}$  and  $_{\mathcal{A}}\mathfrak{D}\mathfrak{M}$  of Hilbert  $\mathcal{A}$ -modules.

Objects  $Ob(_{\mathcal{A}}\mathfrak{M})$ :

Completely contractive unital w\*-contrs. reps (normal reprs.)  $(H, \alpha)$ .

$$\operatorname{Hom}_{\mathcal{A}}(H_1, H_2) = \{ T \in B(H_1, H_2) : T\alpha_1(a) = \alpha_2(a)T \ \forall a \in \mathcal{A} \}.$$

Objects:  $Ob(_{\mathcal{A}}\mathfrak{DM}) = Ob(_{\mathcal{A}}\mathfrak{M}).$ 

$$\operatorname{Hom}_{\mathcal{A}}^{\mathfrak{D}}(H_1, H_2) = \{ T \in B(H_1, H_2) : T\alpha_1(a) = \alpha_2(a)T \ \forall a \in \Delta(\mathcal{A}) \}.$$

Observe that  $_{\mathcal{A}}\mathfrak{M}\subseteq _{\mathcal{A}}\mathfrak{DM};$  If  $\mathcal{A}$  is a  $W^*$  algebra then  $_{\mathcal{A}}\mathfrak{M}= _{\mathcal{A}}\mathfrak{DM}.$ 

5.3 A  $\Delta$ -extension of a functor  $\mathcal{F}$  is a functor

$$\mathcal{F}^{\delta}: \mathcal{A}\mathfrak{DM} \longrightarrow \mathcal{B}\mathfrak{DM}$$

such that

$$\mathcal{A}\mathfrak{M} \hookrightarrow \mathcal{A}\mathfrak{D}\mathfrak{M}$$
 $\mathcal{F}\downarrow \mathcal{F}^{\delta}\downarrow$ 
 $\mathcal{B}\mathfrak{M} \hookrightarrow \mathcal{B}\mathfrak{D}\mathfrak{M}$ 

commutes.

#### 6 The main theorem.

6.1 **Definition** We say that the unital dual operator algebras  $\mathcal{A}, \mathcal{B}$  are  $\Delta$ -equivalent if there exists an equivalence functor

$$\mathcal{F}: A\mathfrak{M} \leftrightarrow B\mathfrak{M}$$

with a  $\Delta$ -extension to an equivalence \*-functor

$$\mathcal{F}^{\delta}: \mathcal{ADM} \leftrightarrow \mathcal{BDM}.$$

We write  $\mathcal{A} \sim_{\Delta} \mathcal{B}$ .

Observe that for W\*-algebras,  $\mathcal{A} \sim_{\Delta} \mathcal{B} \iff \mathcal{A} \sim_{R} \mathcal{B}$ .

6.2 **Theorem** Let  $\mathcal{A}, \mathcal{B}$  be unital dual operator algebras.

$$\mathcal{A} \sim_{\Delta} \mathcal{B} \iff \exists \text{ compl. isom. normal reps: } \alpha(\mathcal{A}) \sim_{TRO} \beta(\mathcal{B}).$$
 (cf. theorem 8)

### 7 Properties of the equivalence functors

Suppose that  $\mathcal{A}, \mathcal{B}$  are unital dual operator algebras and  $\mathcal{A} \sim_{\Delta} \mathcal{B}$  via  $\mathcal{F}$ .

- 7.1  $\mathcal{F}$  is equivalent to a functor  $\mathcal{F}_{\mathcal{U}}$  of 'tensoring by' a suitable bimodule  $\mathcal{U}$ .
- 7.2  $\mathcal{F}$  maps completely isometric representations to completely isometric representations.
- 7.3  $\mathcal{F}$  'preserves reflexivity': If  $\alpha$  is a compl. isom. repr. and  $\beta = \mathcal{F}(\alpha)$  then  $\alpha(\mathcal{A})$  is reflexive iff  $\beta(\mathcal{B})$  is.

7.4  $\mathcal{F}$  'respects' the lattices: If  $(H, \alpha) \in \mathcal{A}\mathfrak{M}$  with corresponding object  $(\mathcal{F}(H), \beta) \in \mathcal{B}\mathfrak{M}$  then

$$\mathcal{F}^{\delta}(\operatorname{Lat}(\alpha(\mathcal{A}))) = \operatorname{Lat}(\beta(\mathcal{B})).$$

7.5  $\mathcal{F}$  is a normal functor, i.e. 'w\*-continuous'.

### 8 Examples and applications.

- 8.1 If  $\mathcal{A} \sim_{\Delta} \mathcal{B}$  then  $\mathcal{A}$  can be completely isometrically reresented as a CSL algebra iff  $\mathcal{B}$  can be so represented.
- 8.2 If  $\mathcal{L}$  is a non-synthetic CSL then  $\mathcal{A}_{min}(\mathcal{L})$  cannot be isometrically represented as a CSL algebra, although its diagonal,  $\mathcal{L}'$ , is a CSL algebra (a vN algebra with abelian commutant). But note  $\mathcal{A}_{min}(\mathcal{L})$  can be represented as a reflexive algebra: consider  $(\mathcal{A}_{min}(\mathcal{L}))^{(\infty)}$  (infinite ampliation).
- 8.3 Two CSL algebras are  $\Delta$ -equivalent iff they are TRO equivalent.
  - It follows from Theorem 2 that two CSL algebras, both with continuous lattices, or both with totally atomic lattices, are  $\Delta$ -equivalent iff they have isomorphic lattices.
- 8.4 If two nest algebras  $\mathcal{A}, \mathcal{B}$  are similar then there exists an equivalence functor

$$\mathcal{F}: \mathcal{A}\mathfrak{M} \leftrightarrow \mathcal{B}\mathfrak{M}.$$

8.5 But they are not always  $\Delta$ -equivalent: There exists an example of similar nest algebras with unitarily equivalent diagonals which are not  $\Delta$ -equivalent.

**Example 9** Let  $\mathcal{N}_1, \mathcal{N}_2$  be the nests of Example 3. Define the 'ordinal sums' of  $\mathcal{N}_1$  and  $\mathcal{N}_2$ , namely the nests

$$\mathcal{L}_1 = \{ N \oplus 0 : N \in \mathcal{N}_1 \} \cup \{ I_{H_1} \oplus M : M \in \mathcal{N}_2 \} \subset B(H_1 \oplus H_2)$$
  
$$\mathcal{L}_2 = \{ M \oplus 0 : M \in \mathcal{N}_2 \} \cup \{ I_{H_2} \oplus N : N \in \mathcal{N}_1 \} \subset B(H_2 \oplus H_1).$$

Now  $\mathcal{N}_1$  and  $\mathcal{N}_2$  are similar (by the similarity theorem [Davidson]), hence so are  $\mathcal{L}_1$  and  $\mathcal{L}_2$ . So by (8.4), if  $\mathcal{A} = \text{Alg}(\mathcal{L}_1)$  and  $\mathcal{B} = \text{Alg}(\mathcal{L}_2)$  the categories  $_{\mathcal{A}}\mathfrak{M}$  and  $_{\mathcal{B}}\mathfrak{M}$  are equivalent. Observe that  $\Delta(\mathcal{A}) \simeq_{unit} \Delta(\mathcal{B})$  because  $\Delta(\mathcal{A}) = \mathcal{N}_1'' \oplus \mathcal{N}_2''$  and  $\Delta(\mathcal{B}) = \mathcal{N}_2'' \oplus \mathcal{N}_1''$ .

Suppose  $\mathcal{A} \sim_{\Delta} \mathcal{B}$ . Then, by (8.3),  $\mathcal{A} \sim_{TRO} \mathcal{B}$ . So by 2 there exists a \*-isomorphism

$$\theta: \Delta(\mathcal{A}) \to \Delta(\mathcal{B})$$
 such that  $\theta(\mathcal{L}_1) = \mathcal{L}_2$ .

Since the diagonals are mass the map  $\theta$  must be unitarily implemented. Now there exist two possibilities:

$$\theta(I_{H_1} \oplus 0) = \begin{cases} M \oplus 0 & \text{for some } N \in \mathcal{N}_2 \ (a) \\ I_{H_2} \oplus N & \text{for some } N \in \mathcal{N}_1 \ (b) \end{cases}$$

$$(a) \implies \mathcal{N}_{1}'' \simeq_{unit} (\mathcal{N}_{2}''|_{M(H_{2})})$$

$$(b) \implies \mathcal{N}_{1}'' \simeq_{unit} (\mathcal{N}_{2}'' \oplus \mathcal{N}_{1}''|_{N(H_{1})})$$

Both lead to a contradiction because  $\mathcal{N}_1''$  is totally atomic while the others are not.